

Electronic supplementary material

Appendix A: Repeated n -player Prisoner's Dilemma for the smallest possible population size N ($N = n$) for general strategy

Each game involves all the N individuals in the population and consists of one or more rounds. We consider two strategies: ALLD and R. ALLD always defect no matter what other individuals do. R can be any strategy that cooperates at least once during each game when there are only one R and $N - 1$ ALLDs in the population. The payoffs, $V(C|l)$ and $V(D|l)$, have the same properties as above, given by (S1)-(S4).

First, let us consider whether strategy R is ESS_N or not. From (S1) and (S3), we have $V(D|l) > V(C|l)$. This inequality means that a single mutant ALLD in a population of strategy R does not have a lower fitness and thus the first condition for R to be ESS_N is violated. Therefore, we conclude strategy R that is ESS_N does not exist.

Second, we examine whether ALLD is ESS_N or not. Since $V(D|1) > V(C|1)$ and a mutant R cooperates at least once, selection opposes R invading ALLD, that is, the first condition for ALLD to be ESS_N is satisfied. The fixation probability of R is given by

$$\rho_R = 1 / \left(1 + \sum_{k=1}^{N-1} \prod_{i=1}^k \frac{1 - w + wG_i}{1 - w + wF_i} \right),$$

where i denotes the number of R individuals and F_i and G_i are the expected payoffs of R and ALLD, respectively (see main text). Since $F_1 < G_1$ and $F_i \leq G_i$ ($2 \leq i \leq N - 1$) from the definition of R, the second condition for ALLD to be ESS_N is satisfied. Therefore, we can conclude that ALLD is ESS_N .

To sum up, even in repeated n -player Prisoner's Dilemma, cooperation cannot evolve when population size and group size are the same. Note that the argument in this

subsection applies for any payoff functions that satisfy (S1)-(S4): we do not have to assume specific payoff functions such $V(C|k) = bk/n - c$ and $V(D|k) = bk/n$. Also, the argument does not presuppose weak selection.

Appendix B: One-shot n -player Prisoner's Dilemma for general population size

Groups of n individuals are sampled from the population of size N and a game is played within each group. A game consists of only one round, in each of which individuals either cooperate or defect. When some individuals cooperate, all the individuals in the group gain a benefit from it while only the cooperating individuals have to pay a cost.

Let $V(C|l)$ and $V(D|l)$ be the payoffs to individuals choosing cooperation and defection given that l of n individuals in the group choose cooperation. The n -player Prisoner's Dilemma demands that these payoffs have the following properties:

$$V(D|l) > V(C|l + 1), \quad (\text{S1})$$

$$V(D|l + 1) > V(D|l), \quad (\text{S2})$$

$$V(C|l + 1) > V(C|l), \quad (\text{S3})$$

$$(l + 1)V(C|l + 1) + (n - l - 1)V(D|l + 1) > lV(C|l) + (n - l)V(D|l), \quad (\text{S4})$$

where $0 \leq l \leq n - 1$. Using (S1) and (S3), we can conclude that D is always traditional ESS while C never is. Now, let us examine for general N whether either or both C and D can be ESS_N . First, consider the case when a mutant D appears in a population of C. Using (S1) and (S3), we can obtain an inequality, that is,

$$(N - 1)V(D|l) > (N - n)V(C|l + 1) + (n - 1)V(C|l).$$

This inequality implies that the first condition for C to be ESS_N is violated (let $a_k = V(D|k - 1)$, $b_k = V(C|k)$ and $l = n - 1$ in (5) in the main text), that is, a single mutant D in a population of C has a higher fitness. Therefore, we conclude that C cannot be ESS_N .

Second, suppose that a mutant C has appeared in a population of D. From (S1) and (S2), we obtain

$$(N-1)V(C|l+1) < (N-n)V(D|l) + (n-1)V(D|l+1),$$

which suggests that the first condition for D to be ESS_N is satisfied (let $a_k = V(C|n-k+1)$, $b_k = V(D|n-k)$ and $l = 0$ in (5) in the main text). The fixation probability of C is given by

$$\rho_C = 1 / \left(1 + \sum_{k=1}^{N-1} \prod_{i=1}^k \frac{1-w+wG_i}{1-w+wF_i} \right),$$

where i denotes the number of C individuals and F_i and G_i are the expected payoffs of C and D, respectively (see main text). From (S1) and (S2), we have

$$\sum_{k=1}^n \frac{\binom{i-1}{n-k} \binom{N-i}{k-1}}{\binom{N-1}{n-1}} V(C|n-k+1) < \sum_{k=1}^n \frac{\binom{i}{n-k} \binom{N-i-1}{k-1}}{\binom{N-1}{n-1}} V(D|n-k),$$

that is, $F_i < G_i$ ($1 \leq i \leq N-1$). Hence, $\rho_C < 1/N$ and thus the second condition for D to be ESS_N is satisfied. Therefore, we can conclude that D is ESS_N .

To sum up, in nonrepeated n -player Prisoner's Dilemma, C is neither ESS_N nor traditional ESS, and D is not only ESS but also ESS_N . Hence, cooperation cannot emerge in populations consisting of defectors. Note that the argument in this subsection applies for any payoff functions that satisfy (S1)-(S4) and for any selection intensity.

Appendix C: For $n=2, 3$, no payoff matrix exists for which both A and B are traditional ESS but neither is ESS_N .

In two-player games,

$$2a_1 + a_2 < 2b_1 + b_2$$

$$a_1 + 2a_2 > b_1 + 2b_2$$

$$a_1 > b_1$$

$$a_2 < b_2$$

All these inequalities cannot be satisfied simultaneously. Hence, for $n=2$, no payoff matrix exists for which both A and B are traditional ESS but neither is ESS_N .

Similarly, in three player games,

$$3a_1+2a_2+a_3 < 3b_1+2b_2+b_3$$

$$a_1+2a_2+3a_3 > b_1+2b_2+3b_3$$

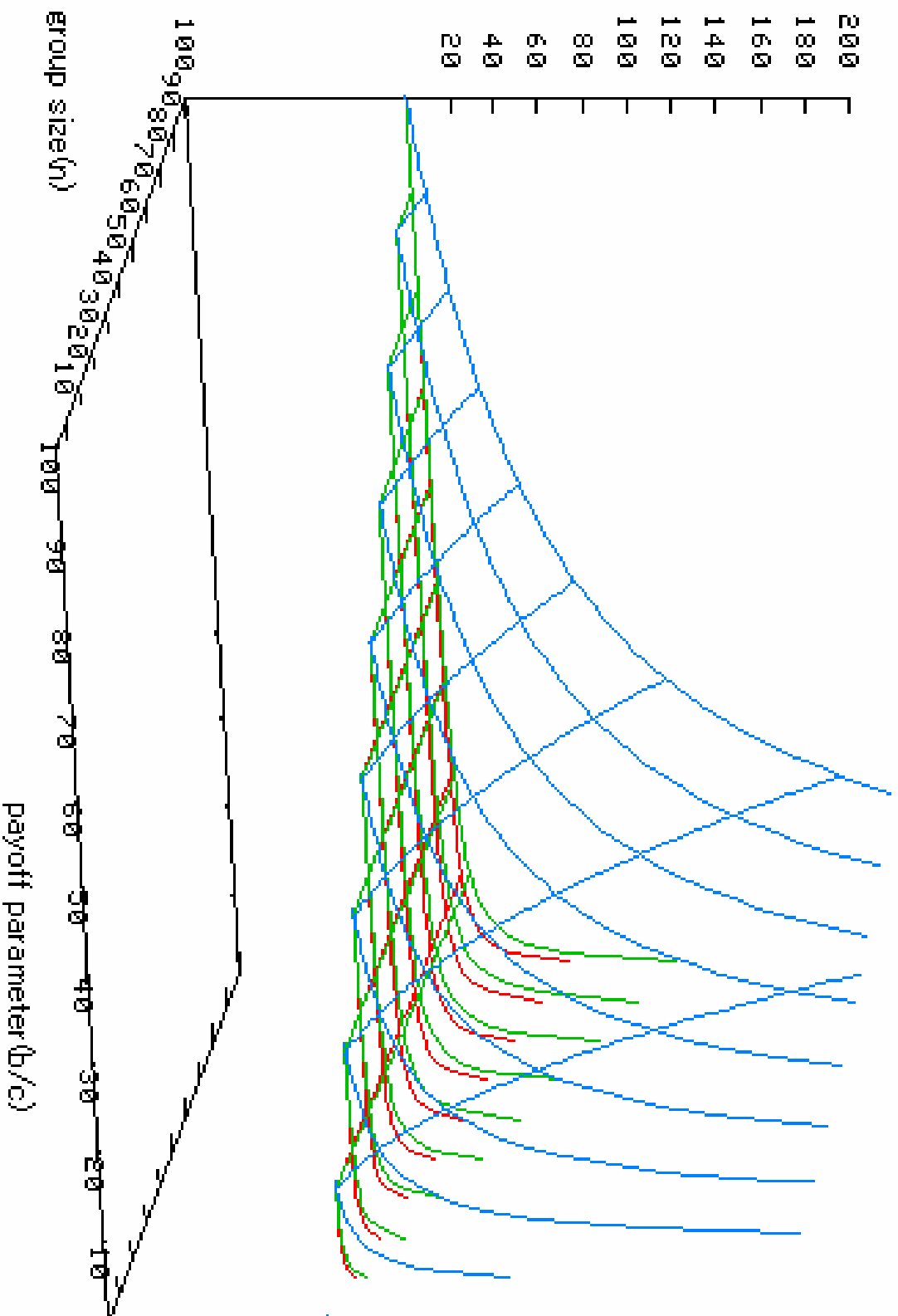
$$a_1 > b_1$$

$$a_3 < b_3$$

All these inequalities cannot be met simultaneously. Hence, for $n=3$, no payoff matrix exists for which both A and B are traditional ESS but neither is ESS_N .

Appendix D: Figure illustrates the relationship between payoff parameter (b/c), group size (n) and the minimum number of rounds (m) required for $\rho_{ALLD} < 1/N$, $\rho_{TFT_{n-1}} > \rho_{ALLD}$ or $\rho_{TFT_{n-1}} > 1/N$ when N is large. The blue represents the number of rounds for which $\rho_{ALLD} = 1/N$, while the green represents that for which $\rho_{TFT_{n-1}} = \rho_{ALLD}$ and the red represents that for which $\rho_{TFT_{n-1}} = 1/N$. The parameter space is divided into the following four regions: (i) TFT_{n-1} is ESS_N and $ALLD$ is not ESS_N ; (ii) both TFT_{n-1} and $ALLD$ are ESS_N and $\rho_{TFT_{n-1}} > \rho_{ALLD}$; (iii) both TFT_{n-1} and $ALLD$ are ESS_N and $\rho_{TFT_{n-1}} < \rho_{ALLD}$; and (iv) TFT_{n-1} is not ESS_N and $ALLD$ is ESS_N . This figure reduces to figure 1a when $b/c = 1.5$ and to figure 1b when $b/(nc) = 0.51$.

the number of rounds(m)



Figure