

## Supplementary Material: Proof of NP-Completeness of FBA-GAP

We specify an instance of the decision problem FBA-GAP as follows:

FBA-GAP[ $R, M, k, l, \mathbf{L}, \mathbf{U}, \mathbf{L}^{src}, \mathbf{U}^{src}, \mathbf{L}^{esc}, \mathbf{U}^{esc}$ ]. Given a set of metabolites  $M$  and reactions  $R$ , does there exist a set of  $k$  source reactions and a set of  $l$  escape reactions that can be added such that there exist fluxes satisfying the lower and upper bounds, including having flux through the biomass reaction more than  $L_{biomass}$ ?

**Theorem.** *FBA-GAP is NP-Complete.*

*Proof.* Follows directly from Lemmas 1 and 2 below.

**Lemma 1.** *FBA-GAP is in NP.*

*Proof.* If a set of source reactions for a set of metabolites  $M_s$  with  $|M_s|=k$  and escape reactions for metabolites  $M_e$  with  $|M_e|=l$  is given, then set  $x_i=1$  for  $i \in M_s$  and  $x_i=0$ , otherwise and set  $y_i=1$  for  $i \in M_e$  and  $y_i=0$ , otherwise. Solve the following linear program

$$\max v_{biomass}$$

Subject to

$$\begin{aligned}
\mathbf{S}\mathbf{v} + \mathbf{b}^{esc} + \mathbf{b}^{esc} &= \mathbf{0} \\
\mathbf{L} &\leq \mathbf{v} \leq \mathbf{U} \\
(\mathbf{L}^{src})^T \mathbf{x} &\leq \mathbf{b}^{src} \leq (\mathbf{U}^{src})^T \mathbf{x} \\
(\mathbf{L}^{esc})^T &\leq \mathbf{b}^{esc} \leq (\mathbf{U}^{esc})^T \mathbf{y}
\end{aligned}$$

If the objective value associated with an optimal solution has  $v_{biomass} \geq L_{biomass}$ , then the sets of source and escape reactions are sufficient. Solving the linear program is polynomial in the size of the inputs, so FBA-GAP is in *NP*.

**Lemma 2.** *FBA-GAP is NP-hard.*

*Proof.* CLOSED HEMISPHERE, known to be strongly NP-complete (Johnson & Preparata, 1978), can be reduced to FBA-GAP. The decision problem for CLOSED HEMISPHERE is stated as follows:

CLOSED HEMISPHERE  $[m, d, \mathbf{A}, n]$ . Given a set of  $m$  linear inequalities  $\mathbf{A}\mathbf{x} \geq \mathbf{0} : \mathbf{A} \in \mathbb{R}^{m \times d}$  does there exist an  $x \in \mathbb{R}^d$  such that  $\|\mathbf{x}\| > 0$  and at least  $n$  of the inequalities are satisfied?

Suppose an instance of CLOSED HEMISPHERE,  $[m, d, \mathbf{A}, n]$  is given. The problem is equivalent to the following intermediate problem: given  $m, d, \mathbf{A}$ , is there an  $x \in \mathbb{R}^d$  with  $\|\mathbf{x}\| > 0$  and an  $\mathbf{s} \in \mathbb{R}_+^m$  such that at least  $n$  of the equalities  $\mathbf{A}\mathbf{x} - \mathbf{s} = \mathbf{0}$  is satisfied?

This intermediate problem can be phrased as a collection of  $d$  instances of FBA-GAP. Each column of  $A$  corresponds to a reaction and each row corresponds to a metabolite. Positive elements in a column of  $A$  are the stoichiometric coefficients for products and negative elements are (the negative of) stoichiometric coefficients for reactants. Each variable  $x_j$  corresponds to the flux through reaction  $j$ . For now, there are no bounds on the flux through each reaction. Each of the  $s_i$  variables corresponds to a one-sided escape reaction for a metabolite.

With the sets of metabolites  $M$  and reactions  $R$  as specified, create an instance of FBA-GAP for each  $x_j, j = 1, \dots, d$  where  $x_j$  corresponds to the flux through the biomass reaction  $v_{biomass}$ . Set  $L_{biomass} = 0$ , for each instance. Each instance of FBA-GAP involves determining if there exists a set of  $k = m - n$  source reactions (and  $l = 0$  escape reactions) such that the flux through the biomass reaction is more than 0.

If the answer to an instance of CLOSED HEMISPHERE is ‘yes’, then there exists a certificate for the instance with  $x_j > 0$  for some  $j$ . This certificate will solve the instance of FBA-GAP where  $x_j$  corresponds to  $v_{biomass}$ . If one of the  $d$  instances of FBA-GAP has a ‘yes’ certificate, then there is a solution to the intermediate problem with the corresponding  $x_j > 0$ . Therefore,  $\|\mathbf{x}\| > 0$ , and the instance of CLOSED HEMISPHERE has a ‘yes’ certificate. The transformation of CLOSED HEMISPHERE to  $d$  instances of FBA-GAP involves only the

introduction of the  $s_i$  variables for each instance, and is therefore a polynomial-time transformation.