

# Supporting Information

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## SI Text

**Stochastic Model of Backtrack Recovery.** We describe the stochastic motion of the backtracked polymerase as a 1D continuous time random walk between different states  $i \in [1, 2, \dots, n-1, n, n+1, \dots, \infty)$ . Each state represents the backtracked distance in nucleotides, with zero being the elongation-competent state. Because the backtrack recovery was experimentally probed at low forces (with a mean value of 1.9 pN), we assume that no external forces are exerted on the polymerase; therefore, its stochastic motion can be described by an unbiased diffusion process. The polymerase is represented by a Brownian particle that jumps between adjacent states with rate  $k$  and uses cleavage to return to zero with rate  $k_c$ . For Pol I and Pol II, the cleavage rate is not constant but depends on the backtrack depth (Fig. 2B),  $\lambda$  being the cleavage cutoff at which  $k_c$  drops to zero. The statistics of backtrack recovery for a finite value of  $\lambda$  can only be obtained numerically.

Because in the case of Pol II-TFIIS, we do not observe a threshold depth in backtrack recovery (Fig. 3A), we consider a constant cleavage rate and  $\lambda \rightarrow \infty$ . In the case of Pol I A12.2ΔC and Pol II ΔRpb9, there is no cleavage reaction, and the backtrack recovery can be described by a 1D diffusion process ( $k_c = 0$ ). Therefore, we derive exact analytical expressions for the statistics of backtrack recovery for both Pol II-TFIIS and the mutants as shown below.

For simplicity and because of the fast elongation rate from  $i = 0$  [ $\sim 10$  times faster than the backtrack diffusion rate (10)], we neglect the backward diffusion rate from zero to one in our model. Therefore, we consider that polymerases that reach the elongation-competent state ( $i = 0$ ) will most likely elongate. With  $p_i(t)$ , we denote the probability of the particle to be at state  $i$  at time  $t \geq 0$ . The dynamics of the system can be described by the following set of master equations:

$$\frac{dp_1}{dt} = kp_2 - (2k + k_c)p_1, \quad [\text{S1}]$$

$$\frac{dp_2}{dt} = kp_3 - (2k + k_c)p_2 + kp_1, \quad [\text{S2}]$$

and

$$\frac{dp_i}{dt} = kp_{i+1} - (2k + k_c)p_i + kp_{i-1} \text{ for } i \geq 3. \quad [\text{S3}]$$

Here,  $p_i \equiv p_i(t)$  for convenience. The elongation-competent state at  $i = 0$  is not explicitly considered.

The recovery time  $\tau_{\text{rec}}$  of a polymerase initially backtracked  $n$  nucleotides equals the first-passage time to the elongation state  $i = 0$ . The calculation of  $\tau_{\text{rec}}$  can be simplified considering a continuous model, where the position of the polymerase,  $x$ , is a continuous random variable. We define  $\rho(x, t)dx$  as the probability for a polymerase to be in the interval  $[x, x+dx]$  at time  $t$ . The probability density  $\rho(x, t)$  evolves in time according to the following Fokker-Planck equation:

$$\frac{\partial \rho(x, t)}{\partial t} = k \frac{\partial^2 \rho(x, t)}{\partial x^2} - k_c \rho(x, t). \quad [\text{S4}]$$

We assume that the polymerase is initially located at  $x = n$  [that is, we set the initial condition to  $\rho(x, 0) = \delta(x - n)$  with an absorbing

boundary at the elongation-competent state  $\rho(0, t) = 0$ ]. The solution of Eq. S4 in the half-plane  $x \in [0, \infty)$  can be expressed as

$$\rho(x, t; n) = \frac{e^{-k_c t}}{\sqrt{4\pi kt}} \left[ e^{-\frac{(x-n)^2}{4kt}} - e^{-\frac{(x+n)^2}{4kt}} \right]. \quad [\text{S5}]$$

The distribution  $\rho(\tau_{\text{rec}}; n)$  of the recovery time for a polymerase initially at  $n$  to be within the interval  $[\tau_{\text{rec}}, \tau_{\text{rec}} + d\tau_{\text{rec}}]$  is given by the probability density current into  $x = 0$ , which contains a contribution of the diffusion and a contribution of the cleavage:

$$\begin{aligned} \rho(\tau_{\text{rec}}; n) &= k \frac{\partial \rho(x, t; n)}{\partial x} \Big|_{x=0, t=\tau_{\text{rec}}} + k_c S(\tau_{\text{rec}}; n) \\ &= \rho_{\text{diff}}(\tau_{\text{rec}}; n) + \rho_c(\tau_{\text{rec}}; n), \end{aligned} \quad [\text{S6}]$$

where  $S(\tau_{\text{rec}}; n)$  is the survival probability or the probability of a polymerase initially at  $n$  to be in  $x > 0$  at time  $\tau_{\text{rec}}$ . The probability density current into  $x = 0$  due to diffusion (8) equals to

$$\rho_{\text{diff}}(\tau_{\text{rec}}; n) = k \frac{\partial \rho(x, t; n)}{\partial x} \Big|_{x=0, t=\tau_{\text{rec}}} = e^{-k_c \tau_{\text{rec}}} \frac{n}{\sqrt{4\pi k \tau_{\text{rec}}^3}} e^{-\frac{n^2}{4k \tau_{\text{rec}}}}. \quad [\text{S7}]$$

At time  $\tau_{\text{rec}}$ , the survival probability equals to

$$S(\tau_{\text{rec}}; n) = \int_0^\infty \rho(x, \tau_{\text{rec}}; n) dx = e^{-k_c \tau_{\text{rec}}} \text{erf} \left( \frac{n}{\sqrt{4k \tau_{\text{rec}}}} \right), \quad [\text{S8}]$$

where erf is the error function. The probability density current into  $x = 0$  by cleavage equals to

$$\rho_c(\tau_{\text{rec}}; n) = k_c e^{-k_c \tau_{\text{rec}}} \text{erf} \left( \frac{n}{\sqrt{4k \tau_{\text{rec}}}} \right). \quad [\text{S9}]$$

Similarly, the probability density  $R(\tau_{\text{rec}}; n)$  of recovery from an initial backtrack depth  $n$  in a time  $\tau_{\text{rec}}$  or recovery probability is given by

$$R(\tau_{\text{rec}}; n) = 1 - S(\tau_{\text{rec}}; n) = 1 - e^{-k_c \tau_{\text{rec}}} \text{erf} \left( \frac{n}{\sqrt{4k \tau_{\text{rec}}}} \right). \quad [\text{S10}]$$

For the case  $k = 0$ , the recovery probability simplifies to

$$R(\tau_{\text{rec}}; n) = \text{erfc} \left( \frac{n}{\sqrt{4k \tau_{\text{rec}}}} \right), \quad [\text{S11}]$$

where erfc is the complementary error function. Eq. S11 is used to fit the experimental data of recovery probability in Pol I and Pol II mutants (Pol I A12.2 ΔC and Pol II ΔRpb9) in Figs. 1 E and F and 2 C and D.

The recovery time probability density can be obtained by summing Eqs. S7 and S9:

$$\rho(\tau_{\text{rec}}; n) = e^{-k_c \tau_{\text{rec}}} \frac{n}{\sqrt{4\pi k \tau_{\text{rec}}^3}} e^{-n^2/4k \tau_{\text{rec}}} + k_c e^{-k_c \tau_{\text{rec}}} \text{erf} \left( \frac{n}{\sqrt{4k \tau_{\text{rec}}}} \right). \quad [\text{S12}]$$

In the presence of both diffusion and cleavage, the recovery time averaged over many polymerases initially backtracked  $n$

nucleotides or mean recovery time,  $\langle \tau_{\text{rec}} \rangle_n$ , is given by the following equation:

$$\langle \tau_{\text{rec}} \rangle_n = \frac{1}{k_c} \left[ 1 - \exp\left(-\frac{n}{\sqrt{\frac{k}{k_c}}}\right) \right]. \quad [\text{S13}]$$

Eq. S13 is used to fit the experimental data of mean recovery time of Pol II–TFIIS in Fig. 5 and Fig. S9.

For a cleavage-deficient polymerase, where  $k_c = 0$ , recovery can only proceed by diffusion. In this case, the recovery time distribution has a power law tail  $\sim \tau_{\text{rec}}^{-3/2}$  with a diverging mean recovery time. Note that, even in the absence of cleavage, a finite value of the average is obtained when the average is done only over polymerases that recover in a given time (5 min in our experiments). An alternative statistic that can be considered is the median recovery time  $\tau_{\text{rec},n}^*$ , which can be obtained from the cumulative recovery time distribution  $C(\tau_{\text{rec}}; n)$  that equals the survival probability at time  $\tau_{\text{rec}}$ :

$$C(\tau_{\text{rec}}; n) = \int_0^{\tau_{\text{rec}}} \frac{n}{\sqrt{4\pi k s^3}} e^{-\frac{n^2}{4ks}} ds = \text{erfc}\left(\frac{n}{\sqrt{4k\tau_{\text{rec}}}}\right). \quad [\text{S14}]$$

At the median recovery time, the cumulative distribution equals 1/2, which yields

$$\tau_{\text{rec},n}^* = \left[ 2 \text{erf}^{-1}\left(\frac{1}{2}\right) \right]^{-2} \frac{n^2}{k} \approx \frac{n^2}{k}, \quad [\text{S15}]$$

because the prefactor  $[2 \text{erf}^{-1}(1/2)]^{-2} \simeq 1.099 \approx 1$ , with  $\text{erf}^{-1}$  being the inverse error function. The mode recovery time  $\hat{\tau}_{\text{rec},n}$  is also finite for a purely diffusive recovery and equal to

$$\hat{\tau}_{\text{rec},n} = \frac{n^2}{6k}. \quad [\text{S16}]$$

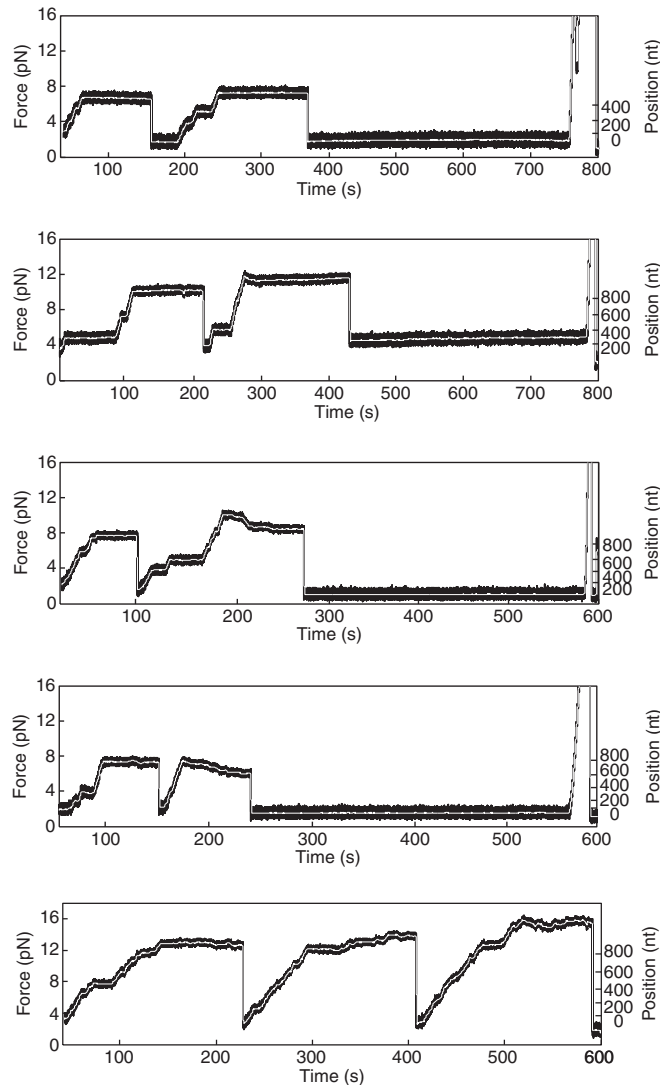
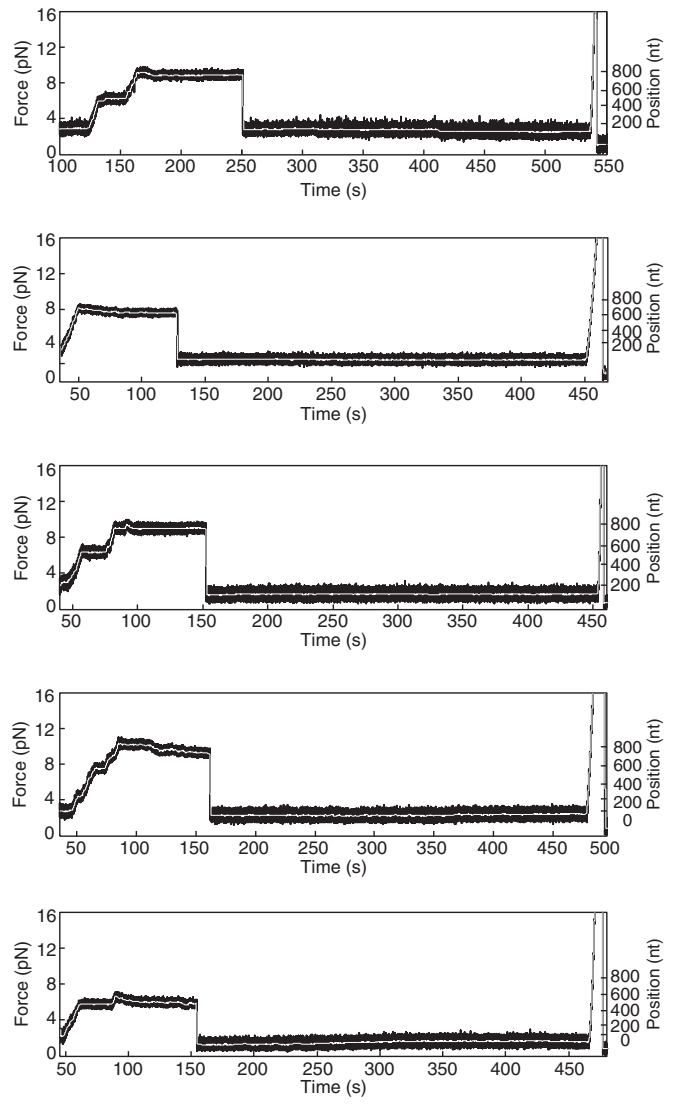


Fig. S11. Additional examples of force reduction experiments that resulted in at least one backtrack recovery.



**Fig. S2.** Additional examples of force reduction experiments that did not result in backtrack recovery.





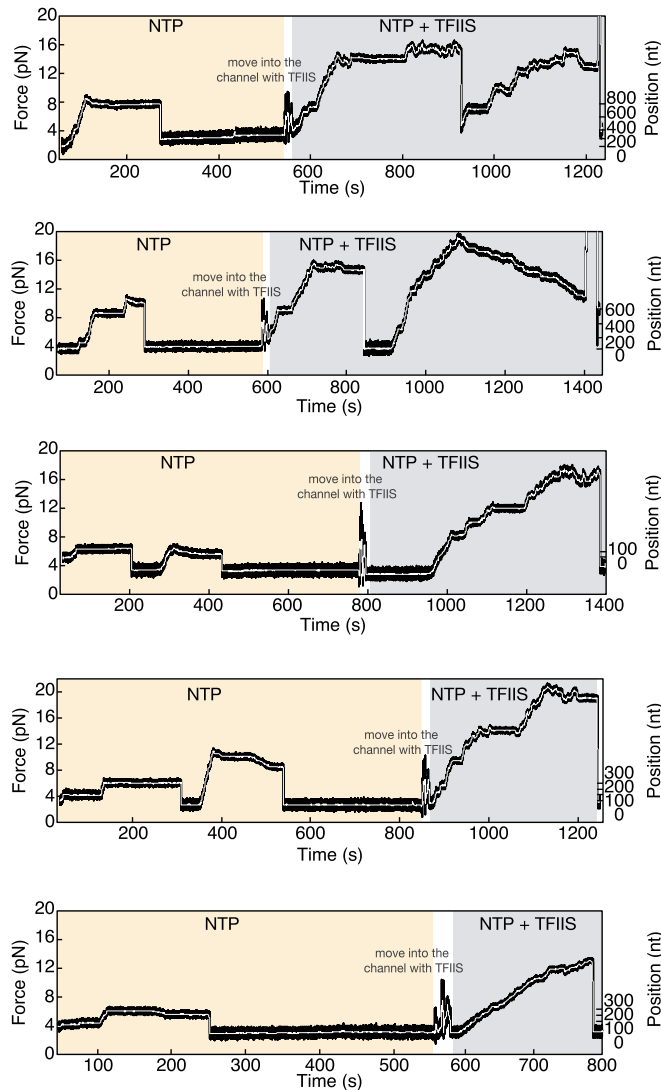


Fig. S7. Additional examples of backtrack recovery experiments of Pol II with adding TFIS after the enzyme has already backtracked.

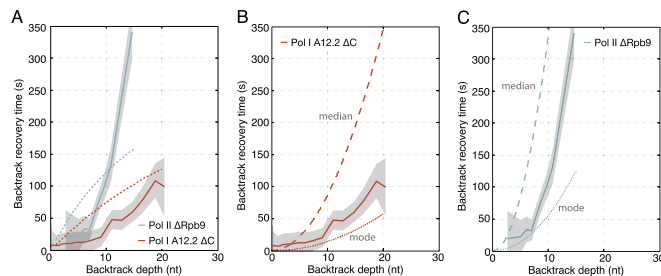


Fig. S8. Backtrack recovery times of Pol I A12.2  $\Delta$ C and Pol II  $\Delta$ Rpb9. Backtrack recovery times (solid lines) as a function of backtrack depth for Pol I A12.2  $\Delta$ C and Pol II  $\Delta$ Rpb9, with SDs obtained by bootstrapping (gray) (*Materials and Methods*). (A) Dashed lines are predictions from the fit of the backtrack recovery probability data (Fig. 2 E and F). Note that the agreement between theoretical prediction and data is not very good here. We speculate that the reason for this is that, specifically for a diffusive first-passage process without cleavage, the recovery time distributions are very broad and not well-represented by an average from a finite and small dataset, like in our experimental data. Please note that the theoretical value of the average recovery time diverges here, because the distribution of the recovery time follows a  $t^{-3/2}$  power law (7) (*SI Text*). (B and C) For this reason, we instead calculate the median (Eq. S15) and the mode (Eq. S16) of the recovery time distribution (*SI Text*) and compare these with the experimental data. Dotted lines represent modes, and dashed lines represent the medians of the backtrack recovery time. Note that the parabolic behavior now is correctly captured in B and C, and the measured values are between the mode and the median of the respective backtrack recovery time.

