

# Clone temporal closeness centrality

In static network settings the closeness centrality of a vertex  $a$  describes how close  $a$  is to all other vertices  $c \in V \setminus a$ , in the network specified by the shortest paths. This centrality measure was extended [1, 2] for a graph sequence  $\mathcal{G} = G_1, \dots, G_S$  by proposing the *temporal closeness centrality* (TCC) for a vertex  $a$  as

$$TCC_{1,S}(a) = \sum_{k=1}^{S-1} \sum_{c \in V \setminus a} \frac{1}{|\gamma_{k,m,S}(a,c)|}, \quad (1)$$

where  $|\gamma_{k,m,S}(a,c)|$  denotes the length of the shortest temporal path between vertices  $a$  to  $c$  over the snapshot sequence  $G_k, \dots, G_S$ . TCC captures the dynamics by summing over the graph sequence that gets successively reduced by one snapshot. In analogy to CTBC, we propose the *clone temporal closeness centrality* (CTCC) as an extension of Eq. (1):

$$CTCC_{1,S}(a) = \sum_{k=1}^S \sum_{j_k=1}^{J_k} \sum_{c \in V \setminus a} \frac{1}{|\gamma_{k,m,S}^{j_k}(a,c)|}, \quad (2)$$

where  $|\gamma_{k,m,S}^{j_k}(a,c)|$  denotes the length of a shortest temporal path starting at the  $j_k$ -th clone of the  $k$ -th snapshot. If there is no such temporal path from  $a$  to  $c$  then the distance between  $a$  and  $c$  is defined as infinite and we assume that  $\frac{1}{\infty} = 0$ .

## References

- [1] Tang, J., Musolesi, M., Mascolo, C., Latora, V., Nicosia, V.: Analysing information flows and key mediators through temporal centrality metrics. In: Proceedings of the 3rd Workshop on Social Network Systems. SNS '10, pp. 3–136. ACM, New York, NY, USA (2010)
- [2] Kim, H., Anderson, R.: Temporal node centrality in complex networks. Physical Review E **85**, 026107 (2012)