

A. Technical algorithmic details

A.1. Identifying conditionally independent parameters

The following algorithm (called `find_conditionally_independent`) is used by MultiBUGS to identify sets of conditionally-independent parameters $W_1, \dots, W_l \subseteq U$:

Input: $G = (E, V)$, a DAG; U , a set of nodes (with identical topological depth)

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 $l \leftarrow 1$ 
 $M \leftarrow \emptyset$ 
while  $|U| > 0$  do
  for  $u$  in  $U$  do
    if  $\text{ch}_G(u) \cap M = \emptyset$  then
       $W_l \leftarrow W_l \cup \{u\}$ 
       $U \leftarrow U \setminus \{u\}$ 
       $M \leftarrow M \cup \text{ch}_G(u)$ 
    end if
  end for
   $M \leftarrow \emptyset$ 
   $l \leftarrow l + 1$ 
end while

```

Output: $\{W_1, \dots, W_l\}$

A.2. Identifying parallelisable likelihoods

Nodes for which the likelihood calculations should be partitioned across cores are identified using the following algorithm, called `find_partial_product_parallel`:

Input: $G = (E, V)$, a DAG; C , a number of cores; h , a topological depth; h^* , the maximum topological depth in G ; T , a computation schedule; r , the current schedule row

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 $\overline{U} \leftarrow D_G^h$ 
 $\overline{\text{ch}} \leftarrow \text{mean}_{v \in S_G} |\text{ch}_G(v)|$ 
for  $u$  in  $U$  do
  if  $|\text{ch}_G(u)| > 2 \times \overline{\text{ch}}$  or  $h^* = 1$  then
     $r \leftarrow r + 1$ 
    for  $c$  in 1 to  $C$  do
       $T_{rc} \leftarrow u$ 
    end for
     $U \leftarrow U \setminus \{u\}$ 
  end if
end for

```

Output: $\{T, U, r\}$

A.3. Building a computation schedule

The overall algorithm for allocating computation to cores is as follows:

Input: $G = (E, V)$, a DAG; C , a number of cores

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Initialise  $T$ , a table with  $C$  columns
 $r \leftarrow 0$ 
 $h^* \leftarrow \max_{v \in S_G} d_G(v)$ 
for  $h$  in  $h^*$  to 1 do
   $\{T, U, r\} \leftarrow \text{find\_partial\_product\_parallel}(G, C, h, h^*, T, r)$ 
   $\{W_1, \dots, W_l\} \leftarrow \text{find\_conditionally\_independent}(G, U)$ 
  for  $i$  in 1 to  $l$  do
     $c \leftarrow 0$ 
    for  $j$  in  $\max_{w \in W_i} |\text{ch}_G(w)|$  to 1 do
      for  $x$  in  $\{w \in W_i : |\text{ch}_G(w)| = j\}$  do
        if  $c \bmod C = 0$  then
           $r \leftarrow r + 1$ 
        end if
      end for
    end for
  end for

```

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         $c \leftarrow 0$ 
    end if
     $T_{r(c+1)} \leftarrow x$ 
     $c \leftarrow c + 1$ 
end for
end for
end for
Output:  $T$ 
```