

1 Example of a Reduced Gradient Nonlinear Programming 2 Algorithm

3 A general equality-constrained nonlinear programming problem is shown in (S1.1). Here, $f(\mathbf{x})$ is the
4 objective function to be minimized, which may be nonlinear. The vector function $g(\mathbf{x})$ is the system of
5 equality constraints, which may in general be nonlinear. The variables, \mathbf{x} , may additionally be bounded by
6 lower and upper bounds, lb and ub , respectively.

$$7 \quad \begin{aligned} & \text{minimize} && f(\mathbf{x}) \\ & \text{s.t.} && g(\mathbf{x}) = \mathbf{0} \\ & && lb \leq \mathbf{x} \leq ub \end{aligned} \quad (S1.1)$$

8 The Lagrangian (S1.2) may be formed from the model components introduced in (S1.1).

$$9 \quad L(\mathbf{x}, \lambda) = f(\mathbf{x}) - \lambda^T g(\mathbf{x}) \quad (S1.2)$$

10 In equation (S1.2), $L(\mathbf{x}, \lambda)$ is the Lagrangian, and λ is the vector of Lagrange multipliers corresponding to
11 the equality constraints, $g(\mathbf{x})$. Each $g_i(\mathbf{x})$ has an associated Lagrange multiplier λ_i . Optimizing an equality-
12 constrained program involves finding stationary points of the Lagrangian (S1.3). A stationary point is
13 identified as one in which $\nabla L(\mathbf{x}, \lambda) = \mathbf{0}$.

$$14 \quad \nabla L(\mathbf{x}, \lambda) = \nabla f(\mathbf{x}) - [\nabla g(\mathbf{x})]^T \lambda = \mathbf{0} \quad (S1.3)$$

15 In (S1.3), $\nabla L(\mathbf{x}, \lambda)$ is the gradient of the Lagrangian with respect to the model variables, \mathbf{x} . This gradient
16 may be zero at a maximum, minimum, or a saddle point of the Lagrangian. The term $\nabla g(\mathbf{x})$ is a matrix
17 whose rows consist of the gradients of the constraints (S1.4), and is also known as the Jacobian of the
18 constraints, $J(\mathbf{x})$.

$$19 \quad \nabla g(\mathbf{x}) = \begin{pmatrix} -\nabla g_1(\mathbf{x})^T & - \\ -\nabla g_2(\mathbf{x})^T & - \\ \vdots & \\ -\nabla g_m(\mathbf{x})^T & - \end{pmatrix} = J(\mathbf{x}) \quad (S1.4)$$

20 With the notational substitution of (S1.4), (S1.3) becomes (S1.5).

$$21 \quad \nabla L(\mathbf{x}, \lambda) = \nabla f(\mathbf{x}) - J(\mathbf{x})^T \lambda = \mathbf{0} \quad (S1.5)$$

22

23 The form shown in (S1.5) is used in constructing the reduced gradient algorithm.

24 The generalized reduced gradient (GRG) algorithm works well for nonlinear optimization problems in
25 which the objective function and constraints are smooth and differentiable [1,2]. It is particularly well-
26 suited for optimization problems in which the total number of equality constraints is similar to the total
27 number of model variables. This is because the GRG algorithm involves partitioning the variables into
28 sets. One set consists of basic variables, \mathbf{x}_b , and the other of nonbasic variables, \mathbf{x}_{nb} . At each iteration of
29 the GRG algorithm, the nonbasic variables are propagated to a better point (one that improves the
30 objective function), and the basic variables are propagated to the new point such that they satisfy the
31 equality constraints [1,3]. The following list of steps in a general GRG algorithm was adapted from Drud
32 [3]:

- 33 1. Find a feasible starting point, \mathbf{x}_0 . At this point, compute the objective function $f(\mathbf{x}_0)$, gradient of the
34 objective function $\nabla f(\mathbf{x}_0)$, and Jacobian of the constraints $J(\mathbf{x}_0)$.
- 35 2. Select a set of n basic variables, \mathbf{x}_b , such that J_b , the submatrix of the basic columns of J , is
36 nonsingular.

- 37 3. Solve for the multipliers $J_b^T \lambda = \nabla_b f$.
 38 4. Compute the reduced gradient $\mathbf{r} = \nabla f - J^T \lambda$.
 39 5. Terminate if a stopping criterion is met.
 40 6. Find a search direction, \mathbf{d} , for the nonbasic variables based on \mathbf{r} and compute the tangent direction.
 41 7. Perform a search along direction \mathbf{d} . For each step, adjust \mathbf{x}_b to satisfy $g(\mathbf{x}_b, \mathbf{x}_{nb}) = \mathbf{0}$ using a pseudo-
 42 Newton process.
 43 8. Repeat these steps with the current point.

44 This process is next illustrated in a simple example. This example GRG algorithm is simply used for
 45 illustration purposes. Large-scale GRG algorithms, such as CONOPT [2,3], are far more sophisticated,
 46 making them much more effective and efficient. Therefore, the following example is illustrative only, and is
 47 not representative of the CONOPT implementation.

48 The example equality-constrained nonlinear program is shown in (S1.6).

$$\begin{aligned}
 & \text{minimize} && f(\mathbf{x}) = 200 x_1^2 + x_2^2 + x_1 x_2 - 360 x_1 - 240 x_2 + 168 \\
 & \text{s.t.} && g_1(\mathbf{x}) = x_1^2 + x_3^2 - x_1 - x_2 + 0.5 = 0 \\
 49 & && g_2(\mathbf{x}) = x_1 x_2 + x_1 + x_4 - 1 = 0 \\
 & && g_3(\mathbf{x}) = x_5 - x_2 - x_1 = 0 \\
 & && \mathbf{x} \geq 0
 \end{aligned} \tag{S1.6}$$

50 This example was chosen because it consists of a nonlinear objective function with nonlinear constraints.
 51 Specifically, $g_3(\mathbf{x})$ is linear, while $g_1(\mathbf{x})$ and $g_2(\mathbf{x})$ are both nonlinear. Also note that (S1.6) is in the general
 52 format shown in (S1.1). Here, the objective function is only a function of two of the model variables, x_1 and
 53 x_2 . Problems with this structure are very common in least-squares model fitting. This problem has a total
 54 of 5 variables and 3 equality constraints.

55 Operating the GRG algorithm requires repeated use of the gradient of the objective function and the
 56 Jacobian of the constraints. They are listed here for the example problem for reference:

$$\nabla f(\mathbf{x}) = \begin{pmatrix} 400 x_1 + 200 x_2 - 360 \\ 400 x_2 + 200 x_1 - 240 \\ 0 \\ 0 \\ 0 \end{pmatrix}, J(\mathbf{x}) = \begin{pmatrix} 2x_1 - 1 & -1 & 2x_3 & 0 & 0 \\ x_2 + 1 & x_1 & 0 & 1 & 0 \\ -1 & -1 & 0 & 0 & 1 \end{pmatrix} \tag{S1.7}$$

58

59 Step 1: Initialize with a feasible starting point

60 Start with a point that satisfies the constraints $g(\mathbf{x}) = \mathbf{0}$. Calculate the objective function, gradient of the
 61 objective function, and Jacobian at this point. We will begin with the arbitrarily chosen feasible point
 62 $\mathbf{x}_0 = [0 \ 0.5 \ 0 \ 1 \ 0.5]^T$. This starting point is chosen simply because it satisfies the constraints. The
 63 objective function $f(\mathbf{x}_0)$, objective function gradient $\nabla f(\mathbf{x}_0)$, and Jacobian $J(\mathbf{x}_0) = J_0$ at this point are
 64 shown in (S1.8). The subscript, 0, is the index of the current iteration.

$$f(\mathbf{x}_0) = 98.0, \quad \nabla f(\mathbf{x}_0) = \begin{pmatrix} -260 \\ -40 \\ 0 \\ 0 \\ 0 \end{pmatrix}, \quad J_0 = \begin{pmatrix} -1 & -1 & 0 & 0 & 0 \\ 1.5 & 0 & 0 & 1 & 0 \\ -1 & -1 & 0 & 0 & 1 \end{pmatrix} \tag{S1.8}$$

66

67 Step 2: Select a set of basic variables

68 To simplify the notation, the subscripted iteration index is omitted from J for the remaining steps. The best
69 set of basic variables is usually not obvious and is determined by the solver, however it is generally wise
70 to choose a set of basic variables such that the submatrix, J_b , of the Jacobian corresponding to the basic
71 variables is well-conditioned. This is because J_b must be invertible [1]. The remaining nonbasic columns
72 of J are stored in matrix J_{nb} . The variable vector, \mathbf{x} , is partitioned accordingly. It is important to note that
73 for a given choice of basic variables, J_b may be well-conditioned at one point, but ill-conditioned at
74 another point. Therefore, the best set of basic variables may change from one iteration to the next. In this
75 example, x_3 would be a poor choice for a basic variable in the first iteration. This is because the column of
76 the Jacobian corresponding to x_3 is all zeros at the initial point. Choosing x_2 , x_4 , and x_5 to be basic, the
77 vectors \mathbf{x}_b and \mathbf{x}_{nb} , and matrices J_b and J_{nb} are shown in (S1.9) for the first iteration.

$$78 \quad \mathbf{x}_b = \begin{pmatrix} x_2 \\ x_4 \\ x_5 \end{pmatrix}, \quad \mathbf{x}_{nb} = \begin{pmatrix} x_1 \\ x_3 \end{pmatrix}, \quad J_b = \begin{pmatrix} -1 & 0 & 0 \\ 0 & 1 & 0 \\ -1 & 0 & 1 \end{pmatrix}, \quad J_{nb} = \begin{pmatrix} -1 & 0 \\ 1.5 & 0 \\ -1 & 0 \end{pmatrix} \quad (\text{S1.9})$$

79

80 Step 3: Solve for the Lagrange multipliers

81 Candidate optima are found as the stationary points of the Lagrangian, defined in (S1.2). These
82 stationary points, \mathbf{x}^*, λ^* , will satisfy $\nabla L(\mathbf{x}^*, \lambda^*) = \mathbf{0}$. Here we seek points that satisfy (S1.10).

$$83 \quad \nabla L(\mathbf{x}, \lambda) = \nabla f(\mathbf{x}) - J(\mathbf{x})^T \lambda = \mathbf{0} \quad (\text{S1.10})$$

84 Considering only the basic variables, (S1.10) becomes (S1.11), which is an exactly determined system for
85 the multipliers, λ .

$$86 \quad \nabla f_b - J_b^T \lambda = \mathbf{0} \quad (\text{S1.11})$$

87 Solving (S1.11) for λ gives (S1.12) for the first iteration.

$$88 \quad \lambda = J_b^T^{-1} \nabla f_b = \begin{pmatrix} -1 & 0 & -1 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}^{-1} \begin{pmatrix} -40 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 40 \\ 0 \\ 0 \end{pmatrix} \quad (\text{S1.12})$$

89 In solvers using the GRG method, J_b^T is typically not inverted at each iteration of the algorithm, as this
90 would be computationally inefficient. Instead, the inverse is computed at the first iteration and updated at
91 each subsequent iteration using a suitable approximation [1].

92

93 Step 4: Compute the reduced gradient

94 Next, the reduced gradient, \mathbf{r} , is computed using (S1.13). The reduced gradient is calculated as the
95 gradient of the Lagrangian using the values of the multipliers calculated in (S1.12).

$$96 \quad \mathbf{r} = \nabla f - J^T \lambda \quad (\text{S1.13})$$

97 Only the nonbasic variables may have nonzero values in the reduced gradient. This results from the
98 values of the multipliers being calculated using (S1.11). The value of the reduced gradient is shown in
99 (S1.14) for the first iteration.

100

$$\mathbf{r} = \begin{pmatrix} -260 \\ -40 \\ 0 \\ 0 \\ 0 \end{pmatrix} - \begin{pmatrix} -1 & 1.5 & -1 \\ -1 & 0 & -1 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 40 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} -220 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} \quad (\text{S1.14})$$

101

102 Step 5: Check the stopping criterion

103 As the iterations proceed, points are generated which become successively closer to optimality. As these
 104 points approach optimality, the gradient of the Lagrangian (S1.5) moves closer to zero, and the
 105 magnitude of the reduced gradient decreases. At optimality, the magnitude of the reduced gradient is
 106 identically zero. Numerically, it is necessary to specify a threshold at which the algorithm terminates to a
 107 point that is "good enough". This optimality tolerance on the magnitude of the reduced gradient is a user-
 108 specified input, ε . At each iteration of the GRG algorithm, (S1.15) is checked, and once satisfied, the
 109 algorithm terminates.

$$110 \quad \|\mathbf{r}\| \leq \varepsilon \quad (\text{S1.15})$$

111 At the input point, the magnitude of the reduced gradient is $\|\mathbf{r}_0\| = 220.0$.

112

113 Step 6: Find a search direction

114 To minimize the objective function subject to the constraints, one choice for the search direction for the
 115 nonbasic variables is $-\mathbf{r}_{nb}$, where \mathbf{r}_{nb} is the reduced gradient of the nonbasic variables. This gives a
 116 steepest descent direction in the nonbasic variables, which will be used in this example. Designating the
 117 search direction as \mathbf{s} gives (S1.16) for the nonbasic variables.

$$118 \quad \mathbf{s}_{nb} = -\mathbf{r}_{nb} \quad (\text{S1.16})$$

119 The search direction for the basic variables is chosen with the goal of minimizing the violation of the
 120 constraints, $g(\mathbf{x}) = 0$. To find this direction, consider each constraint $g_i(\mathbf{x})$ to be a function of \mathbf{x}_b and \mathbf{x}_{nb} , so
 121 that $g_i(\mathbf{x}) = g_i(\mathbf{x}_b, \mathbf{x}_{nb})$. The total differential of $g_i(\mathbf{x}_b, \mathbf{x}_{nb})$ is shown in (S1.17) with \mathbf{x} partitioned in this way
 122 for a single constraint $g_i(\mathbf{x}) = g_i(\mathbf{x}_b, \mathbf{x}_{nb})$ [4].

$$123 \quad dg_i(\mathbf{x}_b, \mathbf{x}_{nb}) = \left(\frac{\partial g_i}{\partial \mathbf{x}_b} \right)^T d\mathbf{x}_b + \left(\frac{\partial g_i}{\partial \mathbf{x}_{nb}} \right)^T d\mathbf{x}_{nb} = \nabla_b g_i^T d\mathbf{x}_b + \nabla_{nb} g_i^T d\mathbf{x}_{nb} \quad (\text{S1.17})$$

124 Considering the full vector of m equality constraints, $g(\mathbf{x}_b, \mathbf{x}_{nb})$, (S1.17) is generalized to (S1.18).

$$125 \quad dg(\mathbf{x}_b, \mathbf{x}_{nb}) = \begin{pmatrix} -\nabla_b g_1(\mathbf{x})^T & - \\ \vdots & \\ -\nabla_b g_m(\mathbf{x})^T & - \end{pmatrix} d\mathbf{x}_b + \begin{pmatrix} -\nabla_{nb} g_1(\mathbf{x})^T & - \\ \vdots & \\ -\nabla_{nb} g_m(\mathbf{x})^T & - \end{pmatrix} d\mathbf{x}_{nb} = J_b(\mathbf{x})d\mathbf{x}_b + J_{nb}(\mathbf{x})d\mathbf{x}_{nb} \quad (\text{S1.18})$$

126 For a small finite change $\Delta\mathbf{x}$, the constraints $g(\mathbf{x}) = 0$ are satisfied to first order using (S1.19), obtained
 127 from (S1.18) by setting the constraint violation, Δg , equal to zero.

$$128 \quad J_b(\mathbf{x})\Delta\mathbf{x}_b + J_{nb}(\mathbf{x})\Delta\mathbf{x}_{nb} = \mathbf{0} \quad (\text{S1.19})$$

129 Rearranging (S1.19) and substituting $\Delta\mathbf{x}_{nb} = \mathbf{s}_{nb}$ and $\Delta\mathbf{x}_b = \mathbf{s}_b$ provides a way to compute \mathbf{s}_b from \mathbf{s}_{nb} at
 130 each iteration (S1.20).

$$131 \quad \mathbf{s}_b = -J_b^{-1}J_{nb}\mathbf{s}_{nb} \quad (\text{S1.20})$$

132 The search direction for the nonbasic variables, s_{nb} , and basic variables, s_b , is thereby determined for the
 133 current iteration. The full search direction, s , is now defined. For the initial point, the search direction is
 134 $s_0 = \begin{bmatrix} 220 & -220 & 0 & -330 & 0 \end{bmatrix}^T$.

135

136 **Step 7: Perform a line search along the search direction. Use a pseudo-**
 137 **Newton process to correct any constraint violation after the step.**

138 Using the search direction, s , the next step is to perform a line search in the direction of s . The next point
 139 is calculated as shown in (S1.21).

140
$$\mathbf{x}_{i+1} = \mathbf{x}_i + \alpha_i \mathbf{s}_i = \mathbf{x}_i + \alpha_i \|\mathbf{s}_i\| \hat{\mathbf{s}}_i \quad (\text{S1.21})$$

141 Here, $\hat{\mathbf{s}}_i$ is the unit vector in the search direction, \mathbf{s}_i , at iteration i and $\|\mathbf{s}_i\|$ is the magnitude of \mathbf{s}_i . Ideally, α is
 142 chosen to find the global minimizer of the objective along the search direction. However, this is highly
 143 computationally expensive in practice. For practical applications, a step size is accepted if it satisfies
 144 some sufficient decrease condition [5]. For the purpose of illustrating the GRG algorithm in this simple
 145 example, α is held constant at $\alpha_i = 0.001$ for each iteration. The magnitude of the reduced gradient, \mathbf{r} , is
 146 large far from an optimum and decreases as iterations approach an optimum. Because s is determined
 147 from \mathbf{r} , $\|\mathbf{s}\|$ will decrease with subsequent iterations, and the step length will automatically decrease as the
 148 optimum is approached despite α being held constant in the example. From the initial point, the next point
 149 is calculated as shown in (S1.22) using (S1.21).

150
$$\tilde{\mathbf{x}}_1 = \begin{pmatrix} 0 \\ 0.5 \\ 0 \\ 1 \\ 0.5 \end{pmatrix} + 0.001 \begin{pmatrix} 220 \\ -220 \\ 0 \\ -330 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0.5 \\ 0 \\ 1 \\ 0.5 \end{pmatrix} + 0.001 \cdot 453.54 \begin{pmatrix} 0.4851 \\ -0.4851 \\ 0 \\ -0.7276 \\ 0 \end{pmatrix} = \begin{pmatrix} 0.2200 \\ 0.2800 \\ 0 \\ 0.6700 \\ 0.5000 \end{pmatrix} \quad (\text{S1.22})$$

151 The “~” decoration is added to \mathbf{x}_1 in (S1.22) because this point may violate some constraints. After the
 152 step, it may be necessary to correct the constraint violation to within a tolerance. The constraint violation
 153 for $\tilde{\mathbf{x}}_1$ in the example is shown in (S1.23).

154
$$g(\tilde{\mathbf{x}}_1) = \begin{pmatrix} 0.0484 \\ -0.0484 \\ 0.0000 \end{pmatrix} \quad (\text{S1.23})$$

155 In the example, only the first two constraints are violated, and the constraint violation is small. To correct
 156 the constraint violation, pseudo-Newton iterations are performed until the constraint violation is satisfied
 157 within a tolerance [4]. The pseudo-Newton iterations are of the form shown in (S1.24).

158
$$\mathbf{y}_{i+1} = \mathbf{y}_i - J(\mathbf{y}_i)^T [J(\mathbf{y}_i)J(\mathbf{y}_i)^T]^{-1} g(\mathbf{y}_i) \quad (\text{S1.24})$$

159 The pseudo-Newton iterations are repeated until the constraints $g(\mathbf{x})$ are satisfied to within a tolerance.
 160 After performing the pseudo-Newton iterations, the new point, \mathbf{x}_1 , is found. For the example, this point is
 161 shown in (S1.25) along with the objective function $f(\mathbf{x}_1)$.

162

$$\mathbf{x}_1 = \begin{pmatrix} 0.2371 \\ 0.3191 \\ 0 \\ 0.6872 \\ 0.5562 \end{pmatrix}, \quad f(\mathbf{x}_1) = 52.80 \quad (\text{S1.25})$$

163

164 **Step 8: Repeat all steps with the new point**

165 After the new point is calculated, repeat all steps until the stopping criterion is met in **Step 5**. The results
166 of this example are shown in **Table S1.1**. The problem is depicted graphically in **Fig S1.1**.

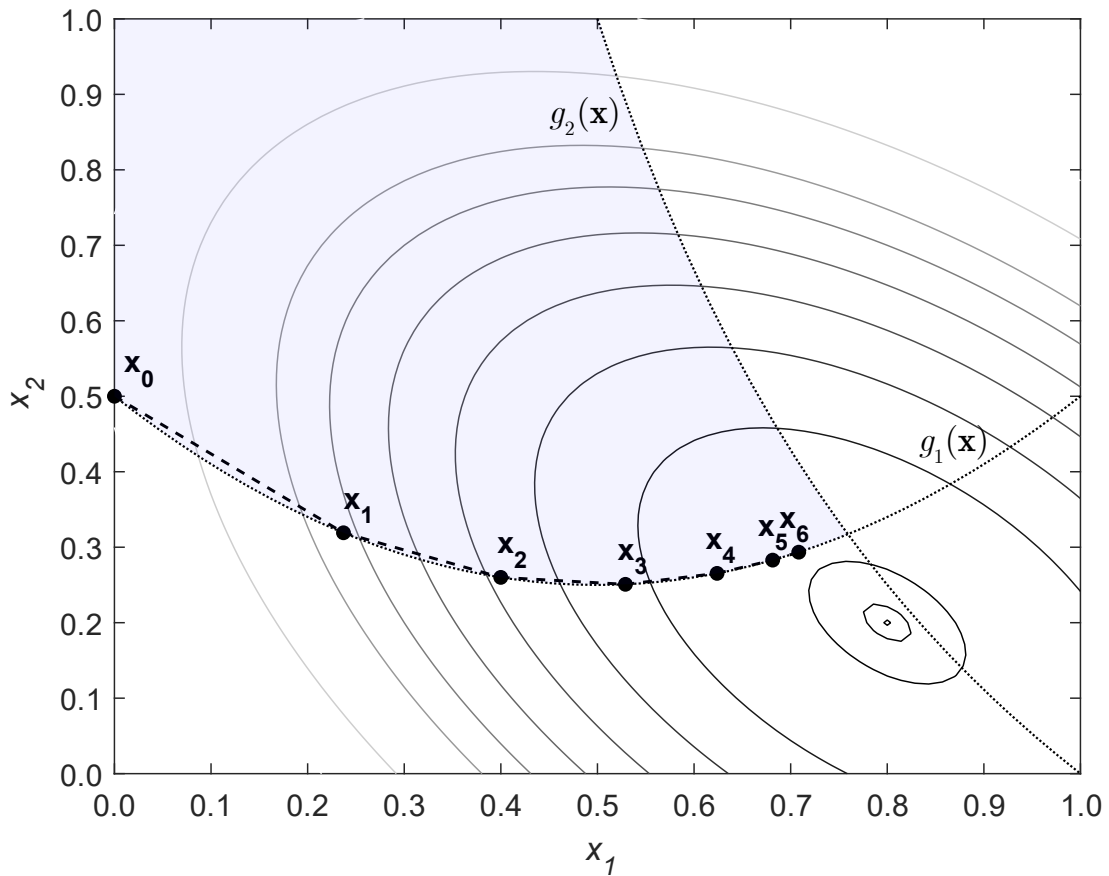
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Table S1.1. The results of the first fifteen iterations of the GRG algorithm are shown for the example problem. The algorithm was operated with α held constant at 0.001. The iterations converge quickly to the optimum of the constrained problem. Iterations 9-15 continue to improve the objective function, however the value was rounded to three decimal places. The magnitude of the reduced gradient approaches zero as the iterations progress, and the algorithm is stopped when this value is within a user-defined tolerance.

Iteration	\mathbf{x}	$f(\mathbf{x})$	$\ \mathbf{r}\ $
0	$[0.0000 \ 0.5000 \ 0.0000 \ 1.0000 \ 0.0000]^T$	98.00	220.0
1	$[0.2371 \ 0.3191 \ 0.0000 \ 0.6872 \ 0.5562]^T$	52.80	167.2
2	$[0.3999 \ 0.2600 \ 0.0000 \ 0.4961 \ 0.6599]^T$	27.93	136.8
3	$[0.5289 \ 0.2508 \ 0.0000 \ 0.3384 \ 0.7798]^T$	12.46	100.2
4	$[0.6238 \ 0.2653 \ 0.0000 \ 0.2107 \ 0.8891]^T$	4.761	59.67
5	$[0.6814 \ 0.2829 \ 0.0000 \ 0.1259 \ 0.9643]^T$	2.222	27.46
6	$[0.7084 \ 0.2934 \ 0.0000 \ 0.0838 \ 1.0018]^T$	1.713	10.05
7	$[0.7183 \ 0.2977 \ 0.0000 \ 0.0678 \ 1.0160]^T$	1.647	3.205
8	$[0.7215 \ 0.2991 \ 0.0000 \ 0.0627 \ 1.0206]^T$	1.640	0.9662
9	$[0.7225 \ 0.2995 \ 0.0000 \ 0.0611 \ 1.0220]^T$	1.639	0.2857
10	$[0.7228 \ 0.2996 \ 0.0000 \ 0.0606 \ 1.0224]^T$	1.639	0.0840
11	$[0.7229 \ 0.2997 \ 0.0000 \ 0.0605 \ 1.0225]^T$	1.639	0.0247
12	$[0.7229 \ 0.2997 \ 0.0000 \ 0.0605 \ 1.0226]^T$	1.639	0.0072
13	$[0.7229 \ 0.2997 \ 0.0000 \ 0.0605 \ 1.0226]^T$	1.639	0.0021
14	$[0.7229 \ 0.2997 \ 0.0000 \ 0.0605 \ 1.0226]^T$	1.639	0.0006
15	$[0.7229 \ 0.2997 \ 0.0000 \ 0.0605 \ 1.0226]^T$	1.639	0.0002

173
174



175

176 **Fig S1.1.** The contours of the objective function are shown for the example problem. The first two
 177 constraints, $g_1(\mathbf{x})$ and $g_2(\mathbf{x})$, are represented in two-dimensional space. In this representation, x_3^2 and x_4
 178 can be thought of as slack variables for the two constraints, $g_1(\mathbf{x})$ and $g_2(\mathbf{x})$, respectively. Because $\mathbf{x} \geq 0$,
 179 the feasible region (shaded in blue) for x_1 and x_2 is bounded below by $g_1(\mathbf{x})$ and above by $g_2(\mathbf{x})$ in this two-
 180 dimensional space. The third constraint, $g_3(\mathbf{x})$ is not plotted because it does not constrain the feasible
 181 space. The variable x_5 may be arbitrarily increased to satisfy $g_3(\mathbf{x})$ for any values of x_1 and x_2 . The first six
 182 iterations of the GRG algorithm are plotted for the example problem. The points progress along the
 183 boundary of $g_1(\mathbf{x})$ and approach the minimum of the objective function as closely as possible without
 184 violating the constraints. In this case, the optimum of the objective function lies outside the feasible region
 185 of the constraints, therefore the optimal solution of the constrained problem is not the same as the
 186 optimum of the unconstrained problem.

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189 **References**

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