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## **Supplemental information**

## Fast and accurate matching of cellular barcodes

### across short-reads and long-reads of single-cell

### **RNA-seq experiments**

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# Supplementary Material

Sample	Brute	Brute match	scTagger =	FLAMES =	scTagger =	% of I Pa	# of I Do
	matched?	unique?	Brute?	Brute?	FLAMES?	70 01 LRS	# OI LRS
N	<ul> <li>✓</li> </ul>	✓	<ul> <li>✓</li> </ul>	<ul> <li>✓</li> </ul>	<ul> <li>✓</li> </ul>	23.5%	1,734,479
	<ul> <li>✓</li> </ul>	✓	<ul> <li>✓</li> </ul>	×	×	26.7%	1,965,586
	<ul> <li>✓</li> </ul>	$\checkmark$	×	<ul> <li>✓</li> </ul>	×	0.1%	3,742
	<ul> <li>✓</li> </ul>	$\checkmark$	×	×	×	1.5%	111,101
	<ul> <li>✓</li> </ul>	×	✓ ✓	×	×	3.2%	235,318
	<ul> <li>✓</li> </ul>	×	×	×	×	20.6%	$1,\!521,\!474$
	×	×	×	×	×	24.4%	1,799,495
NOA1	<ul> <li>✓</li> </ul>	✓	<ul> <li>✓</li> </ul>	✓ <i>✓</i>	<ul> <li>✓</li> </ul>	28.4%	$1,\!606,\!970$
	<ul> <li>✓</li> </ul>	$\checkmark$	<ul> <li>✓</li> </ul>	×	×	28.2%	$1,\!597,\!480$
	<ul> <li>✓</li> </ul>	$\checkmark$	×	✓	×	0.0%	1,586
	✓ ✓	$\checkmark$	×	×	×	0.7%	42,423
	<ul> <li>✓</li> </ul>	×	<ul> <li>✓</li> </ul>	×	×	4.0%	228,219
	<ul> <li>✓</li> </ul>	×	×	×	×	14.2%	803,268
	×	×	×	×	×	24.5%	$1,\!385,\!299$
NOA2	<ul> <li>✓</li> </ul>	$\checkmark$	✓	✓	✓	28.0%	$1,\!321,\!345$
	<ul> <li>✓</li> </ul>	$\checkmark$	$\checkmark$	×	×	28.4%	1,337,904
	<ul> <li>✓</li> </ul>	$\checkmark$	×	$\checkmark$	×	0.1%	4,581
			×	×	×	2.1%	99,454
	<ul> <li>✓</li> </ul>	×	$\checkmark$	×	×	2.0%	93,328
		×	×	×	×	15.3%	720,629
	×	×	×	×	×	24.1%	1,136,768

Table S1: Statistics comparing the matches of FLAMES and scTagger against the Brute-force method, related to Figure 6.

Barcode size $= 4$	$r_1 \cdot \text{ACTTGGT}$	Node #	Read set
		3	$\{r_1, r_2\}$
e = 1	$r_2$ : ACTTGAT	4	$\{r_1, r_2\}$
k-mer size = $4 + 1 = 5$	5	$\{r_1, r_2\}$	
		8	$\{r_1, r_2\}$
1:A 2:C 3:	Γ — 4:T — 5:G	9	$\{r_1, r_2\}$
		10	$\{r_1\}$
	10:G	13	$\{r_1, r_2\}$
Root O:C /:I O:	9:G 16:A	14	$\{r_1\}$
		15	$\{r_1\}$
	14:G 15:T	16	$\{r_2\}$
$11:T \rightarrow 12:T \rightarrow 13:$		17	$\{r_2\}$
		18	$\{r_2\}$

Figure S1: Illustrative example of trie construction in scTagger, related to Figure 2. Assuming barcode size of 4 and maximum allowed error e = 1, the k-mer size is 5. The k-mers of the two LR barcode segments,  $r_1$  and  $r_2$ , are inserted into the trie. The nodes IDs correspond to the order by which they were inserted into the trie. Nodes at layers k - e = to k + e = (i.e. the deepest three layers in the example) contain nodes that barcode alignment can successfully terminate within the allowed error. A map with these nodes as keys is maintained. Each node maps to a set of LR segments from which we extracted k-mers that threaded through the node at insertion time.



Figure S2: Illustrative example of trie querying, related to Figure 2. The different sub-queries of the search at a given node: A) Matching: If the a child node's character matches the next barcode character, we explore this branch while incrementing the barcode index. B) Mismatch: If the a child node's character does not match the next barcode character, we explore this branch while incrementing the barcode index and the total error incurred. C) Deletion: If a character is deleted from the long-read segment, we skip that corresponding barcode index by incrementing the index while stay on the same node in the trie and increment the total error incurred. D) Insertion: If a character is inserted in the long-read segment, we skip that node and to its children without increasing the barcode index while incrementing the total error incurred.

**Algorithm S1** Automatic detection of the ranges of the adapter alignments on the LRs, related to Figure 7.

function $GETRANGES(F)$	$\triangleright$ F[i] is the number adapter alignments on the <i>i</i> -th position
ranges = list()	
while $Q.size > 0$ do	
S = sum(F)	
if $S < 0.01*T$ then	
break	
end if	
Q = Queue()	
Q.enqueue(P)	
first = P	$\triangleright$ first location in the range
last = P	$\triangleright$ last location in the range
while $Q.size > 0$ do	
i = Q.deque()	
if $i \leq P$ and $F[i-1] > 0.0$	01*T then $\triangleright$ Try expanding range to left
Q.enqueue(i-1)	
first=i-1	
end if	
if $i \ge P$ and $F[i+1] > 0$ .	001*T <b>then</b> ▷ Try expanding range to right
Q.enqueue(i+1)	
last=i+1	
end if	
F[i]=0	
end while	
F[first-20:last+20] = 0	$\triangleright$ Set values in range's neighborhood to 0
ranges.append((first,last))	
end while	
return ranges	
end function	

Algorithm S2 Depth-first search in the trie, related to Figure 2.	
function DFS(node, error_budget, index, barcode)	
$\mathbf{if} \text{ index} = \text{barcode.length } \mathbf{then}$	
$edit_distance = MAX_ALLOWED_ERRORS - error_budget$	
$\mathbf{output}$ node.LRs,edit_distance	
end if	
$\mathbf{if} \operatorname{error\_budget} > 0 \mathbf{then}$	
$\mathbf{DFS}(\text{node}, \text{error\_budget} - 1, \text{index} + 1, \text{barcode})$	$\triangleright$ Deletion
end if	
for child in node.children do	
$\mathbf{if} \ barcode[index] = child.char \mathbf{then}$	
$\mathbf{DFS}(\text{child}, \text{error}_\text{budget}, \text{index} +1, \text{barcode})$	$\triangleright$ Match
end if	
$\mathbf{if} \text{ barcode[index]} \mathrel{!=} \mathbf{child.char then}$	
$\mathbf{DFS}(\text{child}, \text{error}_\text{budget} - 1, \text{index} + 1, \text{barcode})$	$\triangleright$ Mismatch
end if	
if $\operatorname{error\_budget} > 0$ then	
$\mathbf{DFS}(\text{child}, \text{error}_\text{budget} - 1, \text{ index}, \text{ barcode})$	$\triangleright$ Insertion
end if	
end for	
end function	
function QUERYTRIE(barcode)	
DFS(root, MAX_ALLOWED_ERRORS, 0, barcode)	
end function	