

Supplementary material

Commands

Approximate Bayesian Branch Support Annotation

```
iqtree2 -s SEQ_ALIGNMENT -te GENE_TREE -m TVM+I+G4 -abayes -pre ANNOTATED_GENE_TREE
```

Note: When inferring support as a post-processing step, the same model used for inferring the tree should be used, a task that requires care when the original trees are inferred using a different tool (e.g., RAxML). TVM+I+G4 is simply an example.

Running wASTRAL

Exact commands when running on gene trees with approximate Bayesian/Bootstrap/SH-like supports.

```
astral-hybrid -x 1 -n 0.333 APPROXIMATE_BAYESIAN_ANNOTATED_GENE_TREE
```

```
astral-hybrid -x 100 -n 0 BOOTSTRAP_ANNOTATED_GENE_TREE
```

```
astral-hybrid -x 1 -n 0 SH_LIKE_ANNOTATED_GENE_TREE
```

Table S1. Counters w^* are defined for each node w in each gene tree, and Q is defined globally. Here, X, Y, Z are distinct colors of A, B , and C . Let u, v be the children of w ; e be the parental edge of w ; p be the parent of w ; $\mathcal{P}_{x,w}$ be the path between x and w ; $s(\mathcal{P}) = 1 - \prod_{\hat{e} \in \mathcal{P}} (1 - s(\hat{e}))$; $m(i, j) = \text{MRCA of } i \text{ and } j$. Counters for leaves are set to zero unless explicitly noted. For each counter, we show a recursive equation on top and the equivalent non-recursive definition on the bottom.

w_X	$(u_X + v_X)e^{-l(e)}$ for internal node w ; $e^{-l(e)}$ for leaf node w colored X
	$\sum_i e^{-l(\mathcal{P}_{i,p})}$ for all leaf nodes i colored X under w
(w_{XX}^+, w_{XY}^+)	$(u_{XX}^+ + v_{XX}^+ + u_X v_X + u_{XY}^+ + v_{XY}^+ + u_X v_Y + u_Y v_X)$
	$\sum_{i,j} e^{-l(\mathcal{P}_{i,j})}$ for all leaf nodes i colored X and j colored X/Y under w
(w_{XX}^-, w_{XY}^-)	$(u_{XX}^- + v_{XX}^- + u_X v_X)(1 - s(e)), (u_{XY}^- + v_{XY}^- + u_X v_Y + u_Y v_X)(1 - s(e))$
	$\sum_{i,j} e^{-l(\mathcal{P}_{i,j})} (1 - s(\mathcal{P}_{m(i,j),p}))$ for all leaf nodes i colored X and j colored X/Y under w
$(w_{XX Y}, w_{XY Z})$	$(u_{XX Y} + v_{XX Y} + (u_{XX}^+ - u_{XX}^-)v_Y + u_Y(v_{XX}^+ - v_{XX}^-))e^{-l(e)},$ $(u_{XY Z} + v_{XY Z} + (u_{XY}^+ - u_{XY}^-)v_Z + u_Z(v_{XY}^+ - v_{XY}^-))e^{-l(e)}$
	$\sum_{i,j,k} e^{-l(\mathcal{P}_{i,j}) - l(\mathcal{P}_{k,p})} s(\mathcal{P}_{m(i,j),m(i,k)})$ for leaf nodes i colored X , j colored X/Y , k colored Z under w , and $m(i, j)$ under $m(i, k)$
$w_{XX YZ}$	$v_X u_Y Z X + u_X v_Y Z X + u_{XX} v_Y + v_{XX} u_Y + u_{XX Y} v_Z + v_{XX Y} u_Z$ $+ (u_{YZ}^+ v_{XX}^- - u_{YZ}^- v_{XX}^+) + (u_{XX}^+ v_{YZ}^+ - u_{XX}^- v_{YZ}^-)$
	$\sum_{h,i,j,k} w_G(hi jk)$ for all leaf nodes h, i colored X , j colored Y , k colored Z , and $w = \text{MRCA } h, i, j, k$
Q	$\sum_{G \in \mathcal{G}} \sum_w (w_{AA BC} + w_{BB AC} + w_{CC AB})$ for internal nodes w in G
	$\sum_{G \in \mathcal{G}} \sum_{h,i,j,k} w_G(hi jk)$ for leaf nodes h, i, j, k in G where h, i have the same color and i, j, k have different colors; when species coloring matches all gene trees, $Q = W[A B C] = \sum_{G \in \mathcal{G}} W(A B C, G)$ (Proposition 5).

Table S2. Running time of species tree inference methods on biological datasets. We use 5.17.3 version of ASTRAL-III if not otherwise clarified.

Dataset	n	k	Method	#Cores	Wall-clock time	CPU time
OneKP	1178	410	wASTRAL-h	16	17.1 min	4.57 hr
			ASTRAL-III (5.0.3)	1	17.2 hr	17.2 hr
Canis	48	449450	wASTRAL-h	1	17.7 hr	17.7 hr
Avian	48	14446	wASTRAL-h	16	1.76 min	28.1 min
			ASTRAL-III	16	20.9 min	5.57 hr
Cetacean	98	3191	wASTRAL-h	16	35.2 sec	9.39 min
			ASTRAL-III	16	1.97 min	31.5 min
Nomiinae	32	853	wASTRAL-h	1	5.93 sec	5.93 sec
			ASTRAL-III	1	8.64 sec	8.64 sec
Lepidoptera	203	1930	wASTRAL-h	16	2.02 min	32.3 min
			ASTRAL-III	16	9.14 min	2.44 hr
Papilionidae	61	6405	wASTRAL-h	16	24.8 sec	6.61 min
			ASTRAL-III	16	1.11 min	17.8 min

Supplementary Figures and Tables

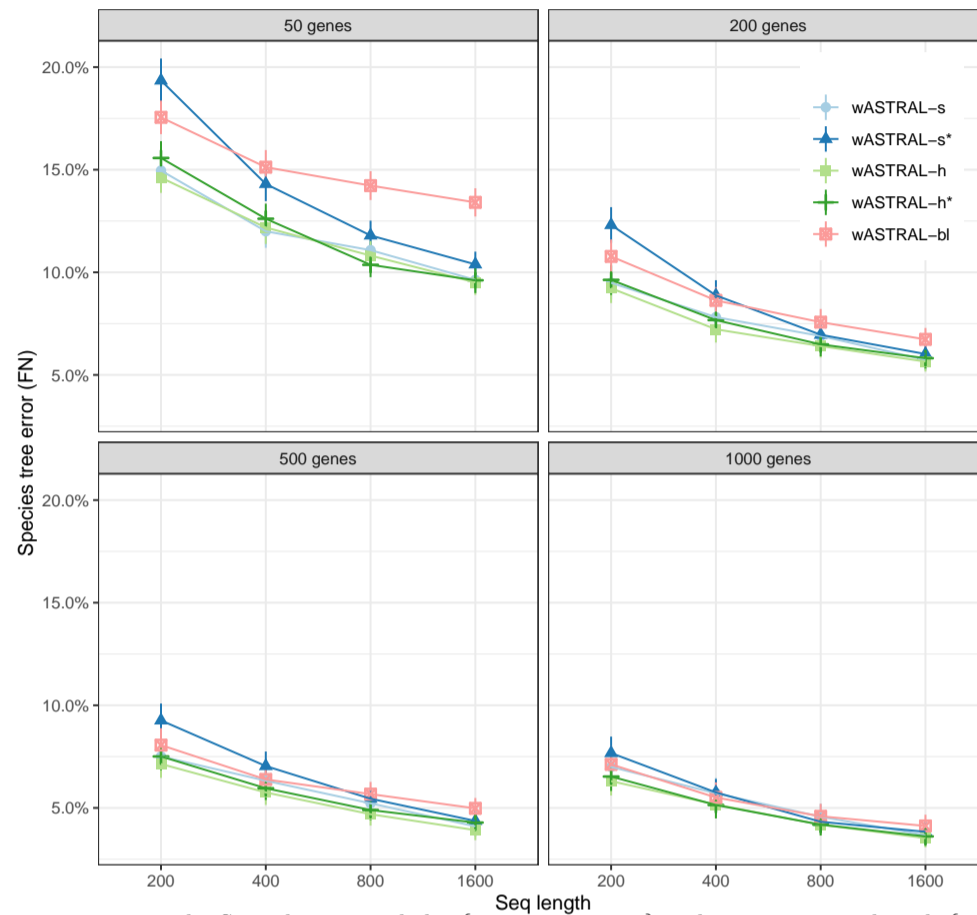


FIG. S1. Species tree error on the S100 dataset with $k = \{50, 200, 500, 1000\}$ and gene sequence length $\{200, 400, 800, 1600\}$. Top: comparison of by weighting scheme. Results with aBayes supports are labeled wASTRAL-s and wASTRAL-h; results with bootstrap support are labeled wASTRAL-s* and wASTRAL-h*. Bottom: comparison with other methods, similar to Figure 2a, but with a y-axis kept fixed.

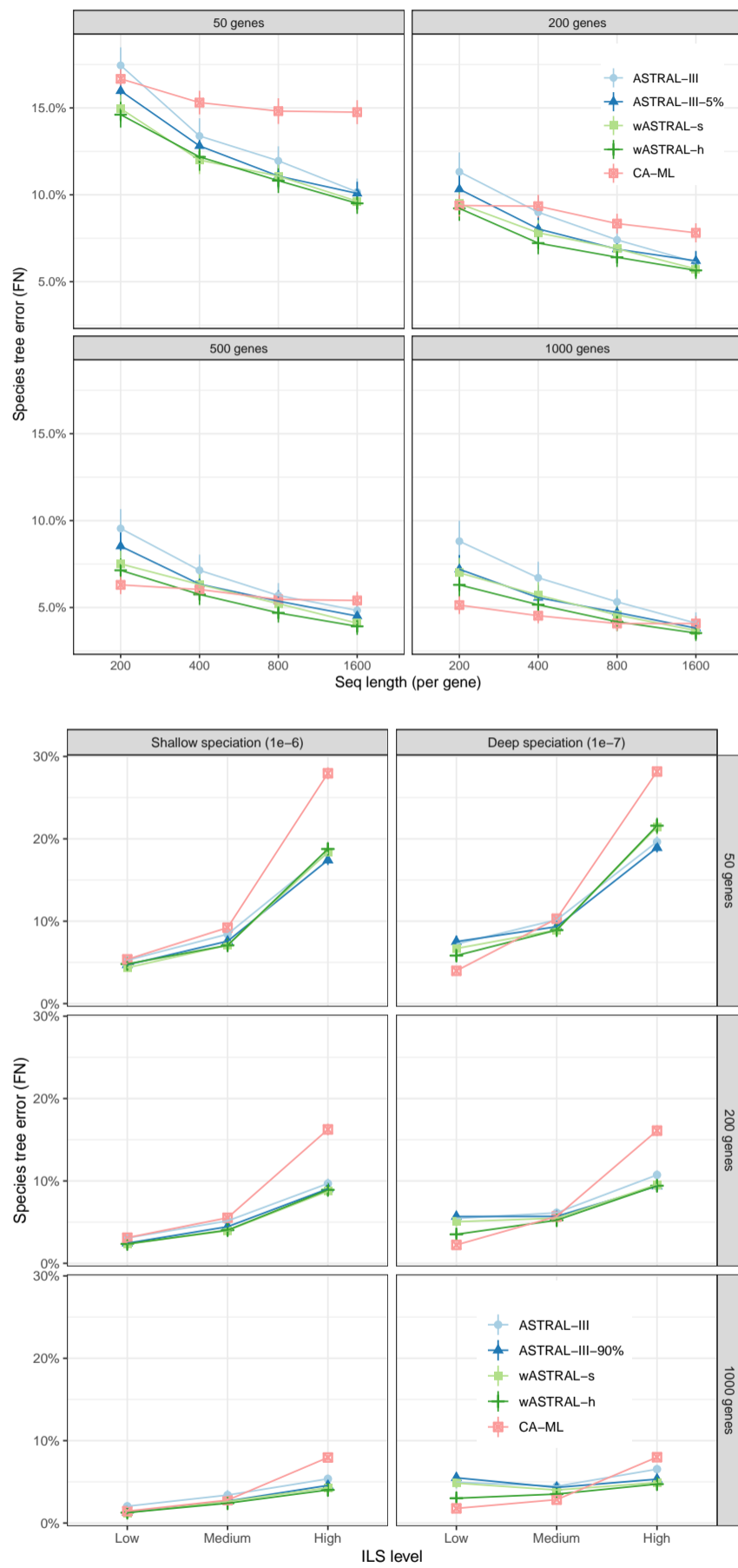


FIG. S2. Species tree error on the S100 dataset (top) and S200 dataset (bottom), similar to Figures 2 and 3 of the main paper, but with y axis kept fixed.

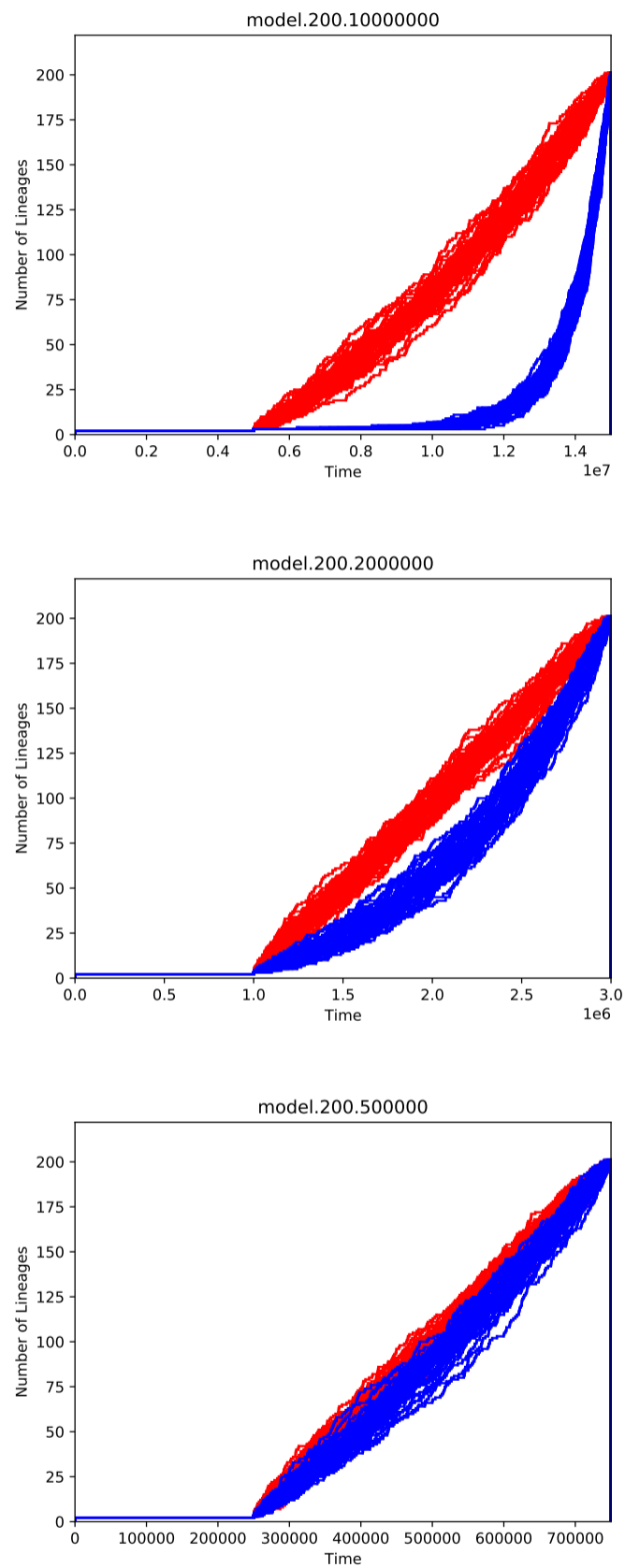


FIG. S3. Lineage Through Time (LTT) plots for three simulated model conditions with 10^{-7} (red) and 10^{-6} (blue) rates tend to lead to deeper and shallower speciation.

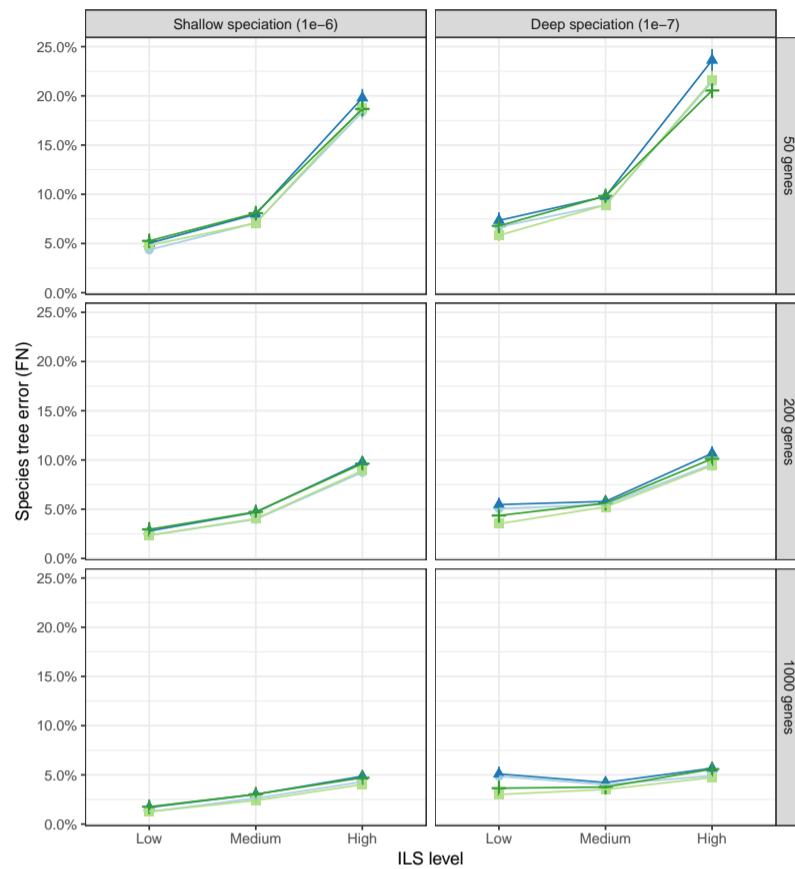


FIG. S4. Species tree error by weighting scheme on the S200 dataset with $k = \{50, 200, 1000\}$ and population size (ILS levels). Species tree shape with parameters E1-6 and E1-7 are used. Results with aBayes supports are labeled wASTRAL-s and wASTRAL-h; results with SH-like support are labeled wASTRAL-s* and wASTRAL-h*.

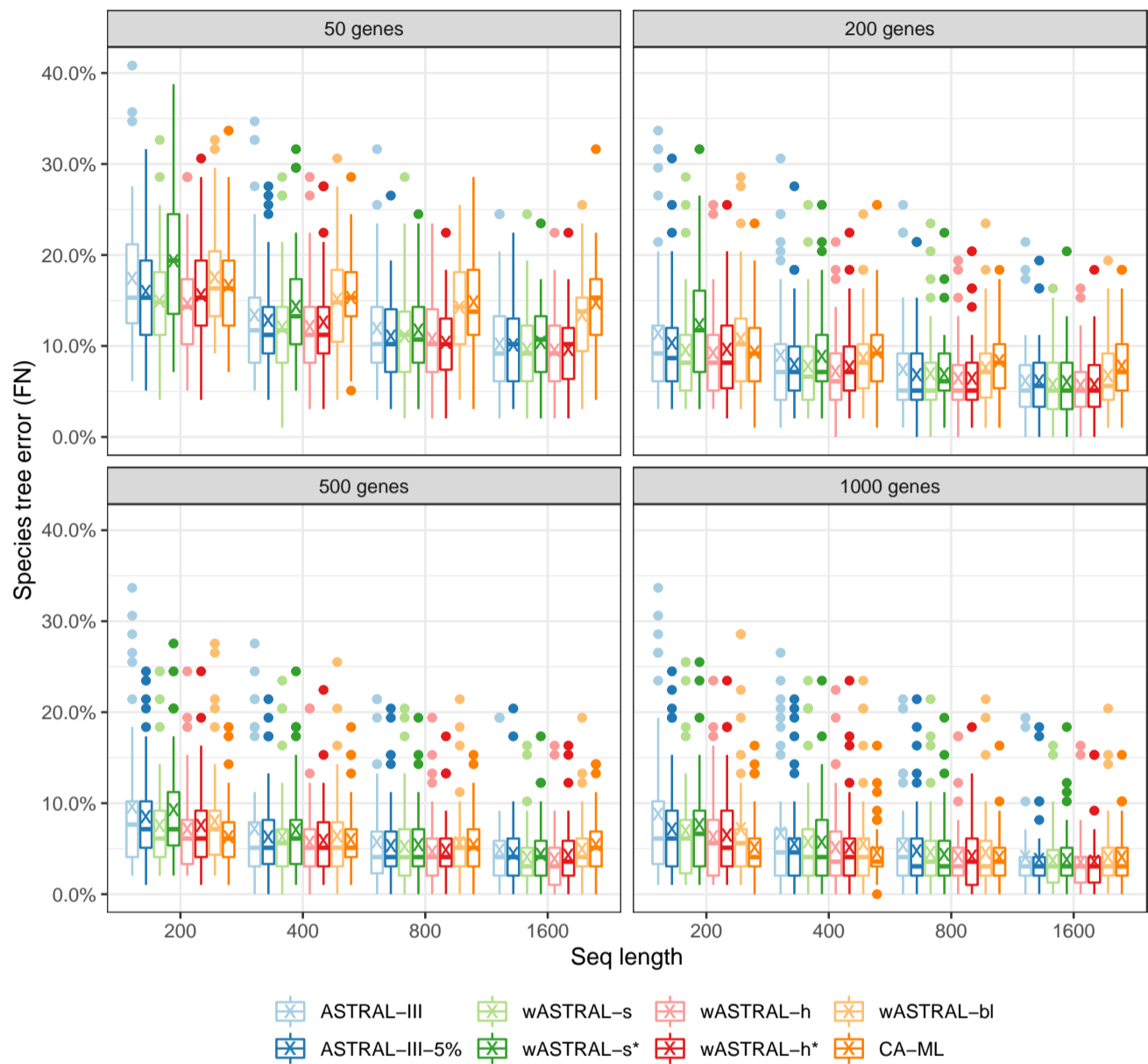


FIG. S5. Species tree error on the S100 dataset with $k = \{50, 200, 500, 1000\}$ and gene sequence length $\{200, 400, 800, 1600\}$. Results with aBayes supports are labelled wASTRAL-s and wASTRAL-h; results with bootstrap support are labelled wASTRAL-s* and wASTRAL-h*.

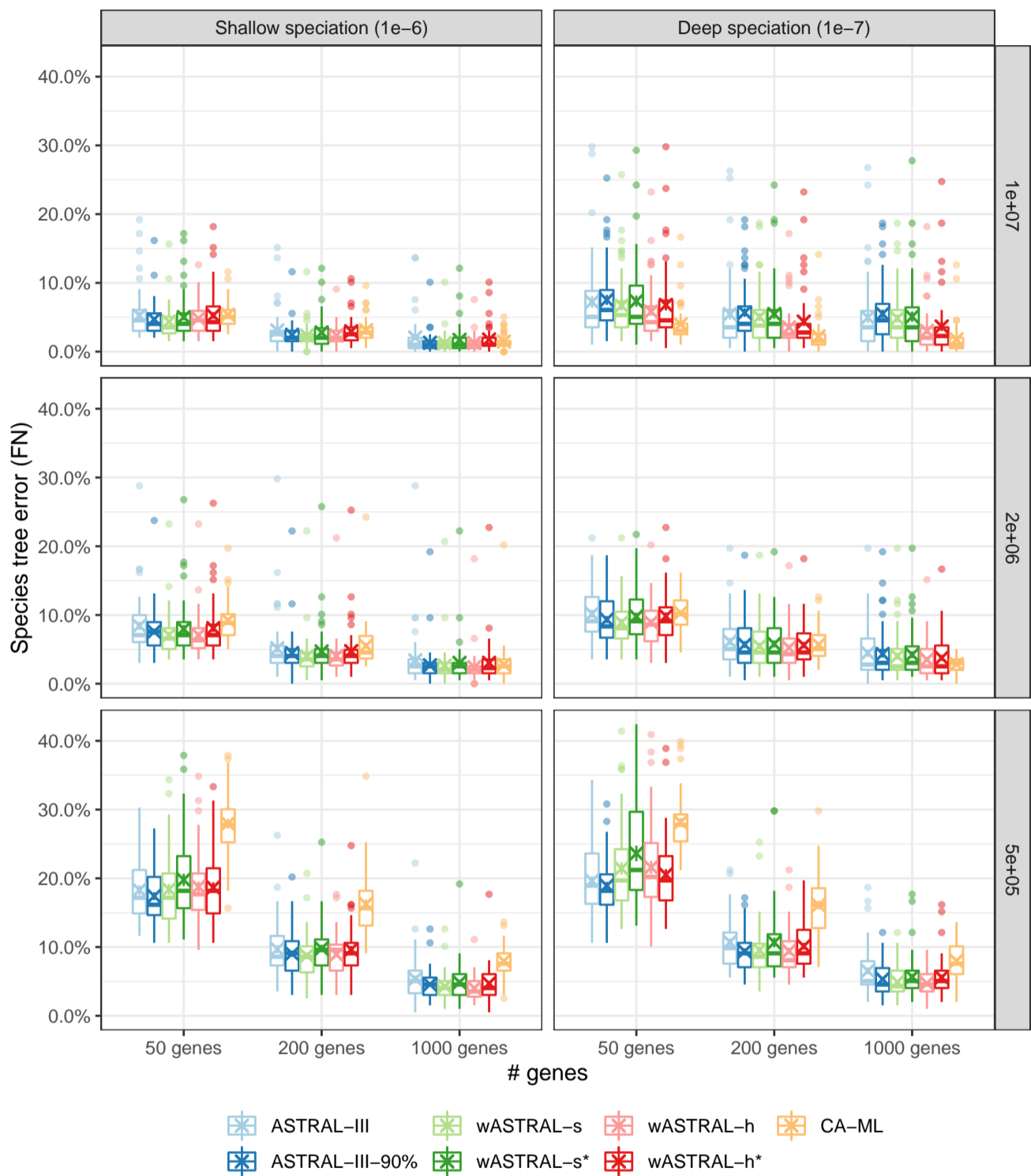


FIG. S6. Species tree error on the S200 dataset with $k = \{50, 200, 1000\}$ and population size (ILS levels). Species tree shape with parameter E1-6 and E1-7 (box columns) and ILS levels (box rows) low ($1e+07$), medium ($2e+06$), and high ($5e+05$) are used. Results with Bayesian supports are labeled wASTRAL-s and wASTRAL-h; results with SH-like support are labeled wASTRAL-s* and wASTRAL-h*.

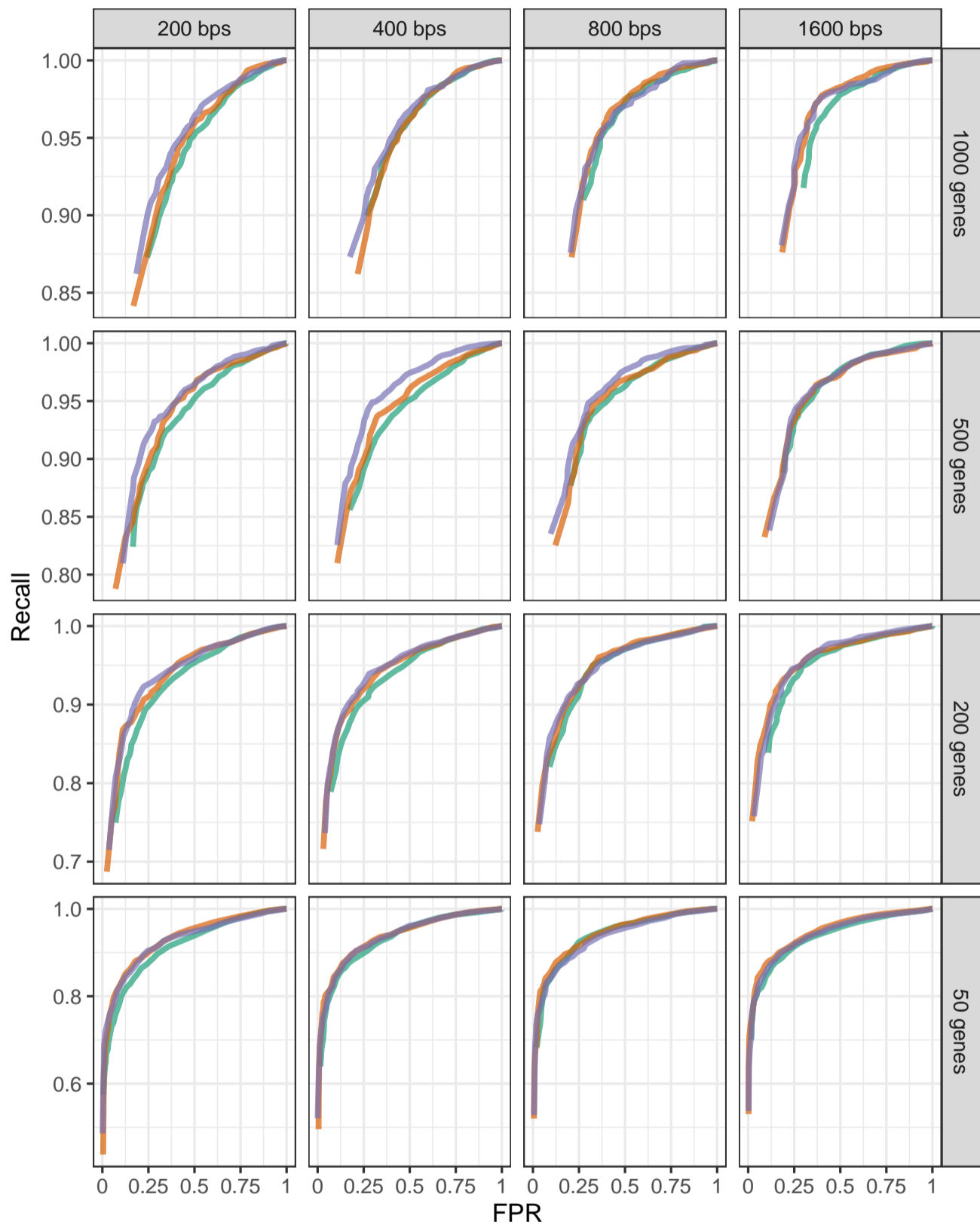


FIG. S7. ROC of S100 dataset with $k = \{50, 200, 500, 1000\}$ and gene sequence length $\{200, 400, 800, 1600\}$ as we change the threshold of support considered. Results with aBayes supports are labelled wASTRAL-s and wASTRAL-h; results with FastTree-2 bootstrap support are labelled wASTRAL-s* and wASTRAL-h*.

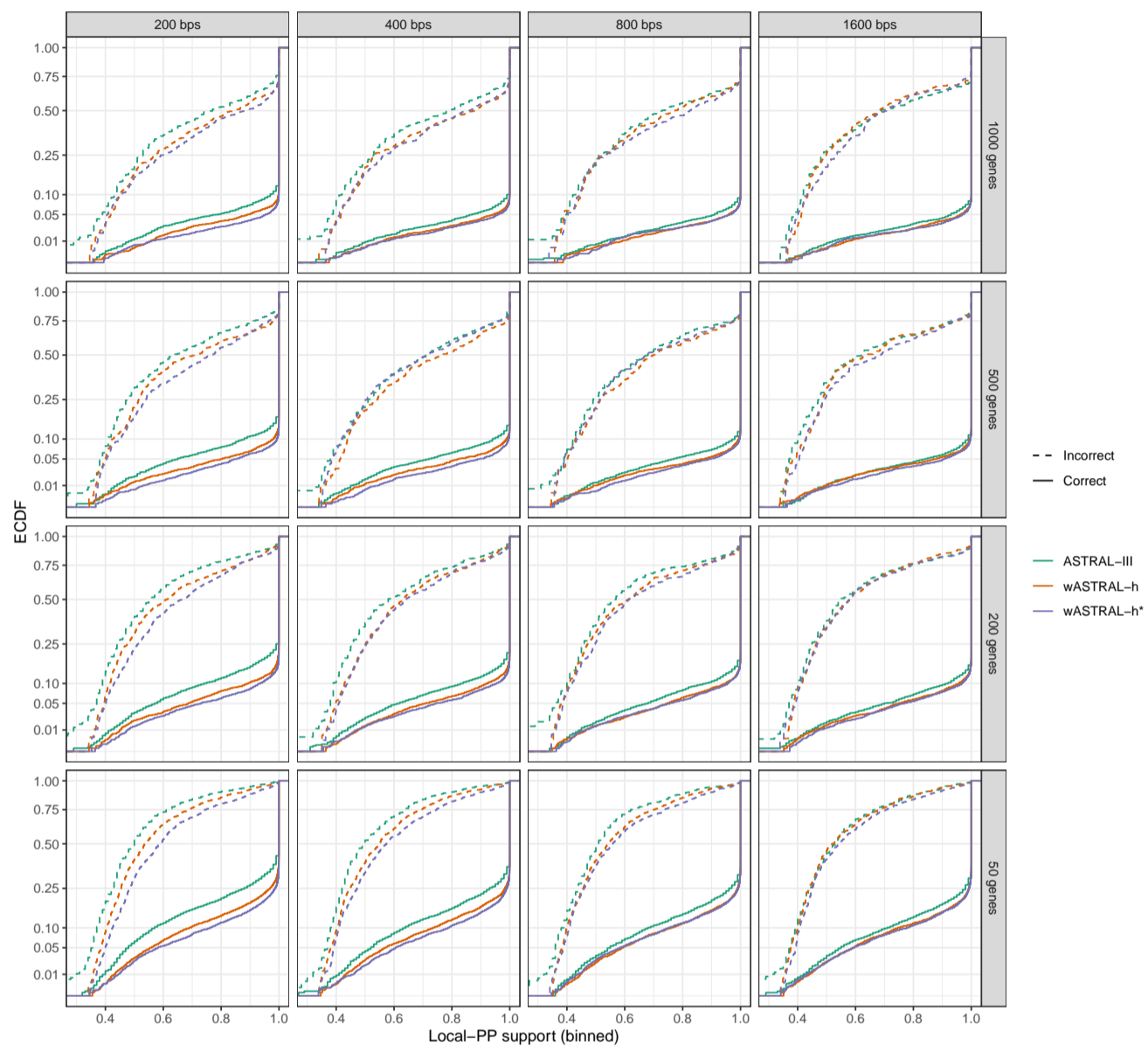


FIG. S8. ECDF of S100 dataset with $k = \{50, 200, 500, 1000\}$ and gene sequence length $\{200, 400, 800, 1600\}$. Results with aBayes supports are labelled wASTRAL-s and wASTRAL-h; results with FastTree-2 bootstrap support are labelled wASTRAL-s* and wASTRAL-h*.

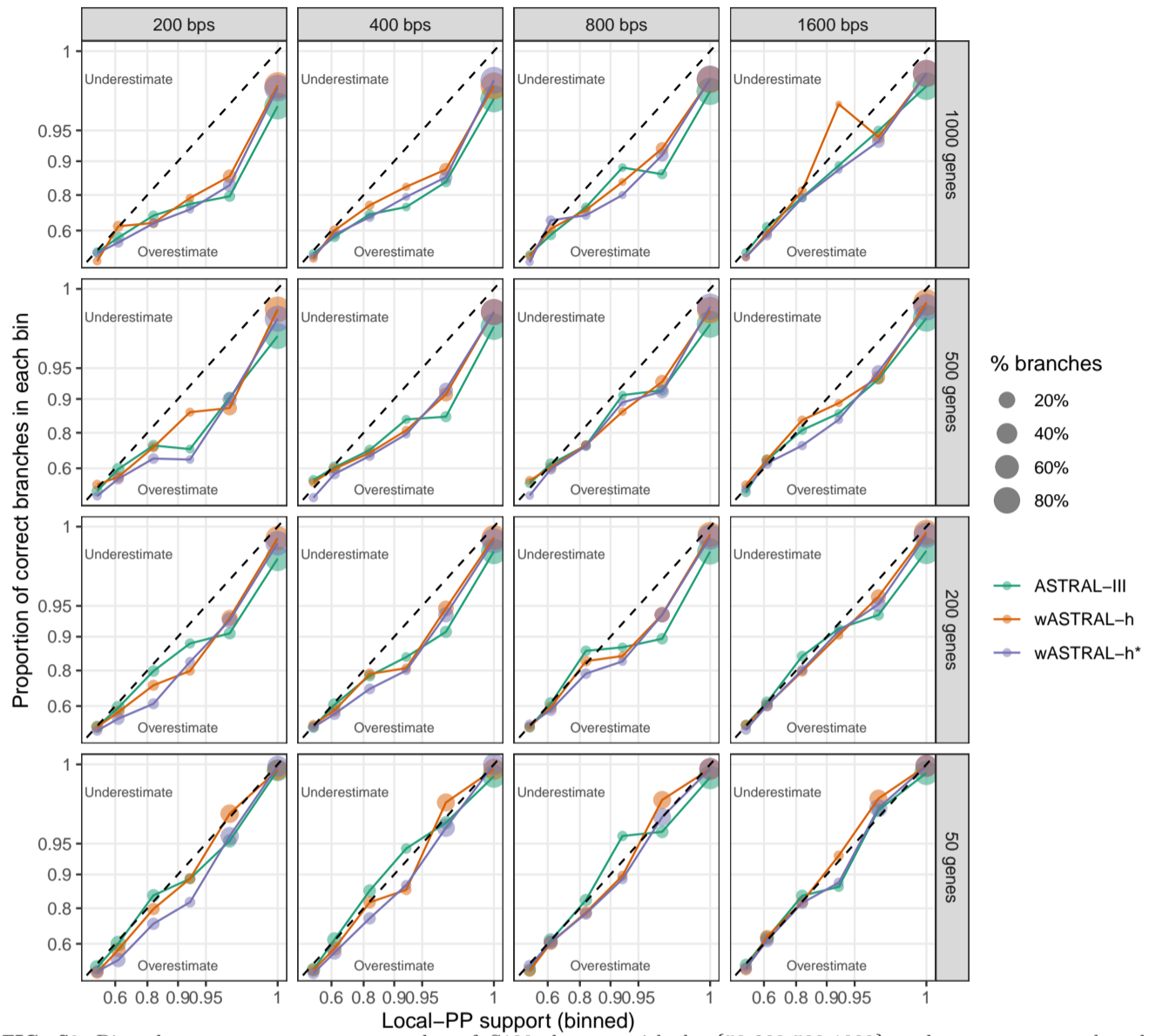


FIG. S9. Binned accuracy-verses-support plot of S100 dataset with $k = \{50, 200, 500, 1000\}$ and gene sequence length $\{200, 400, 800, 1600\}$. Results with aBayes supports are labelled wASTRAL-s and wASTRAL-h; results with FastTree-2 bootstrap support are labelled wASTRAL-s* and wASTRAL-h*.

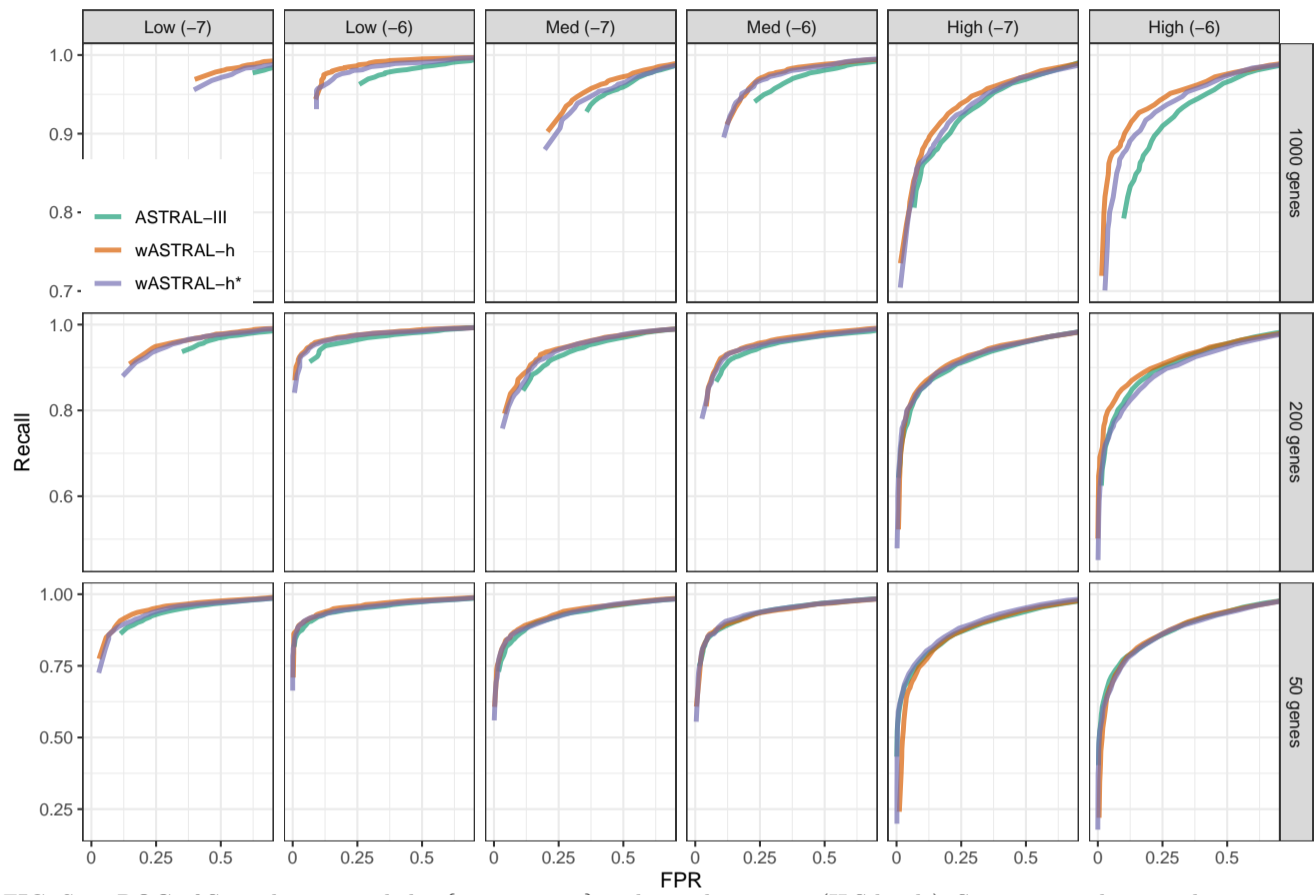


FIG. S10. ROC of S200 dataset with $k = \{50, 200, 1000\}$ and population size (ILS levels). Species tree shape with parameter E1-6 and E1-7 are used. Results with aBayes supports are labeled wASTRAL-h; results with SH-like support are labelled wASTRAL-h*.

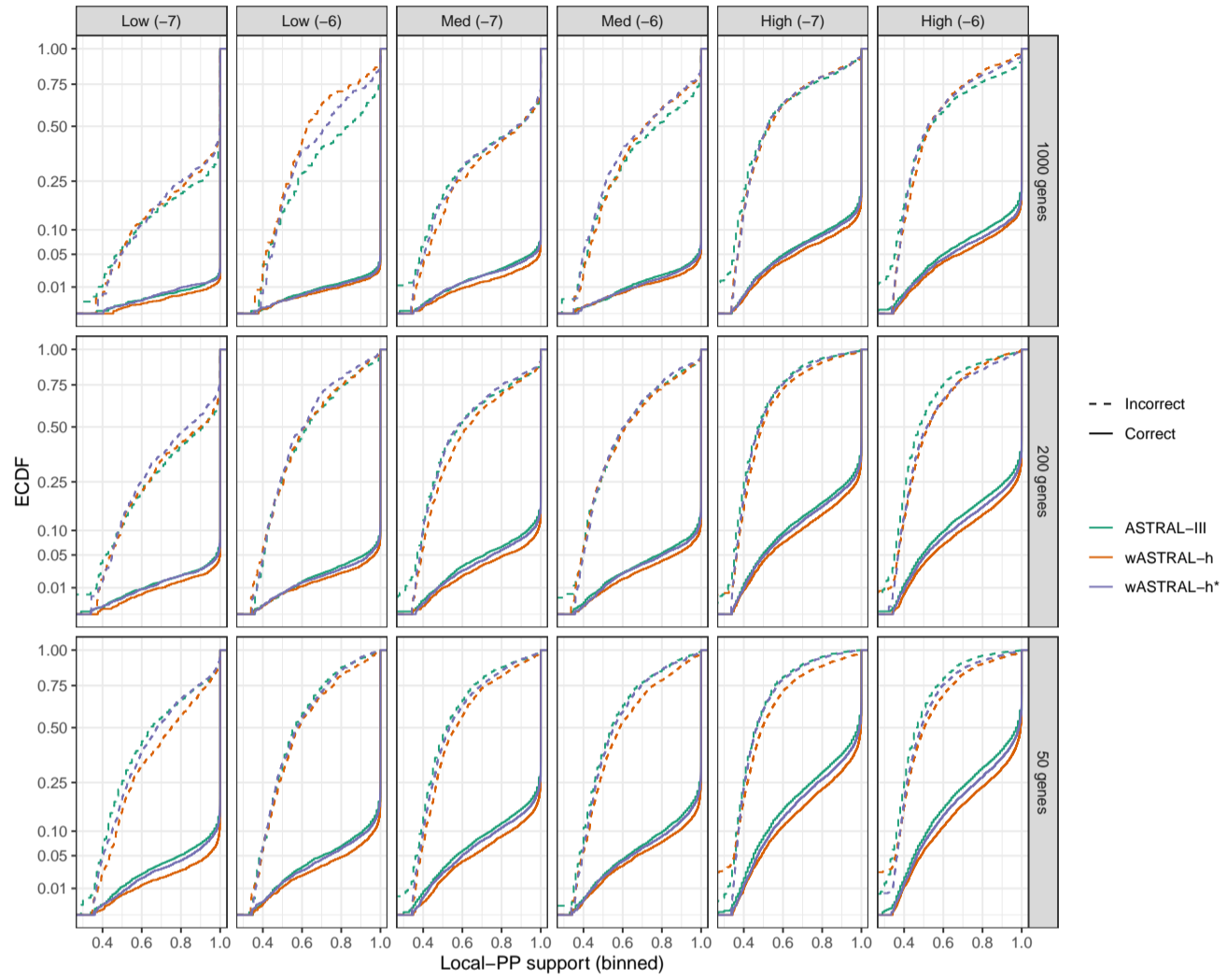


FIG. S11. ECDF of S200 dataset with $k = \{50, 200, 1000\}$ and population size (ILS levels). Species tree shape with parameter E1-6 and E1-7 (box columns) and ILS levels (box rows) low ($1e+07$), medium ($2e+06$), and high ($5e+05$) are used. Results with aBayes supports are labelled wASTRAL-h; results with SH-like support are labelled wASTRAL-h*.

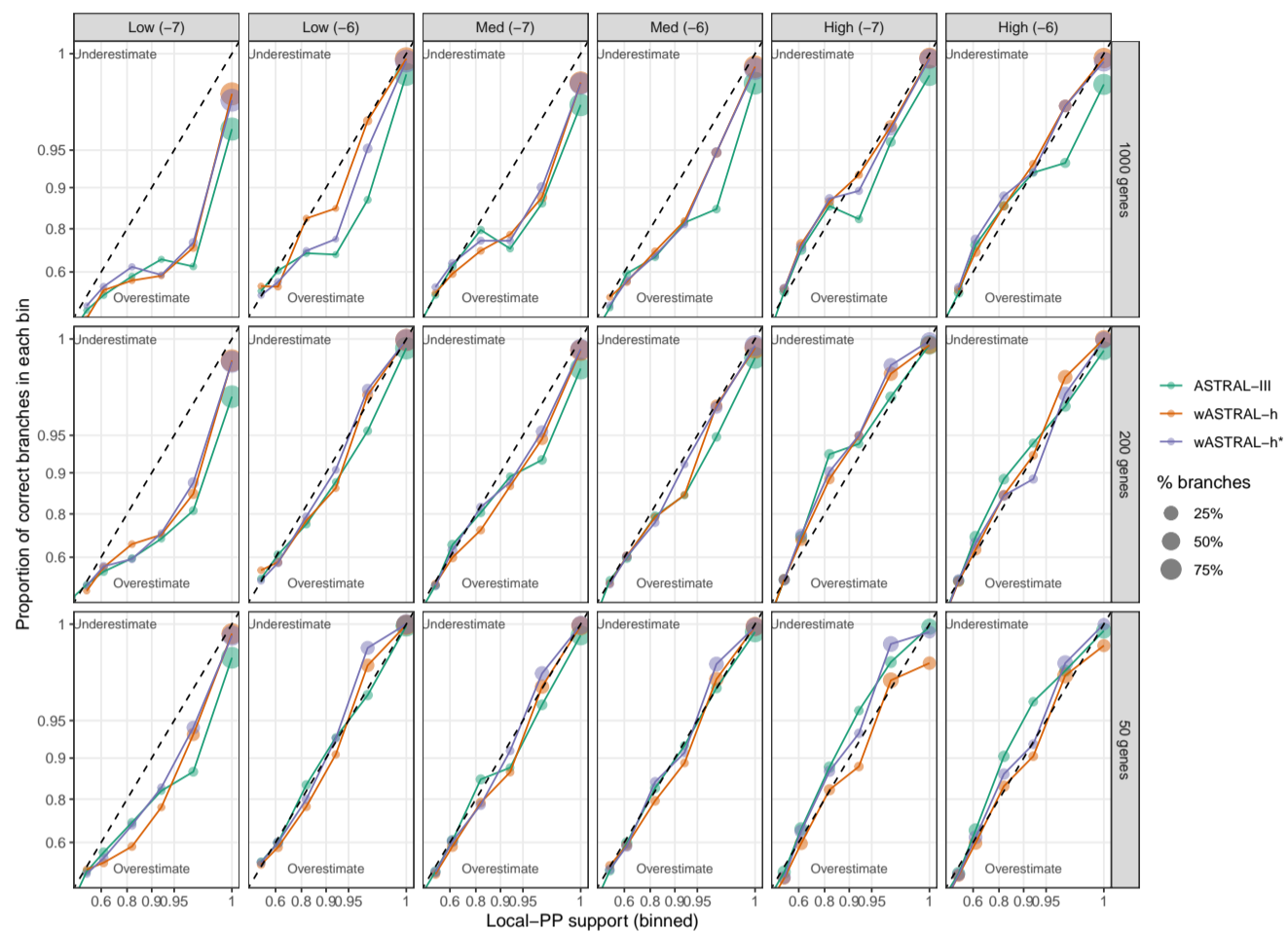


FIG. S12. Binned accuracy-verses-support plot of S200 dataset with $k = \{50, 200, 1000\}$ and population size (ILS levels). Species tree shape with parameter E1-6 and E1-7 are used. Results with aBayes supports are labeled wASTRAL-h; results with SH-like support are labeled wASTRAL-h*.

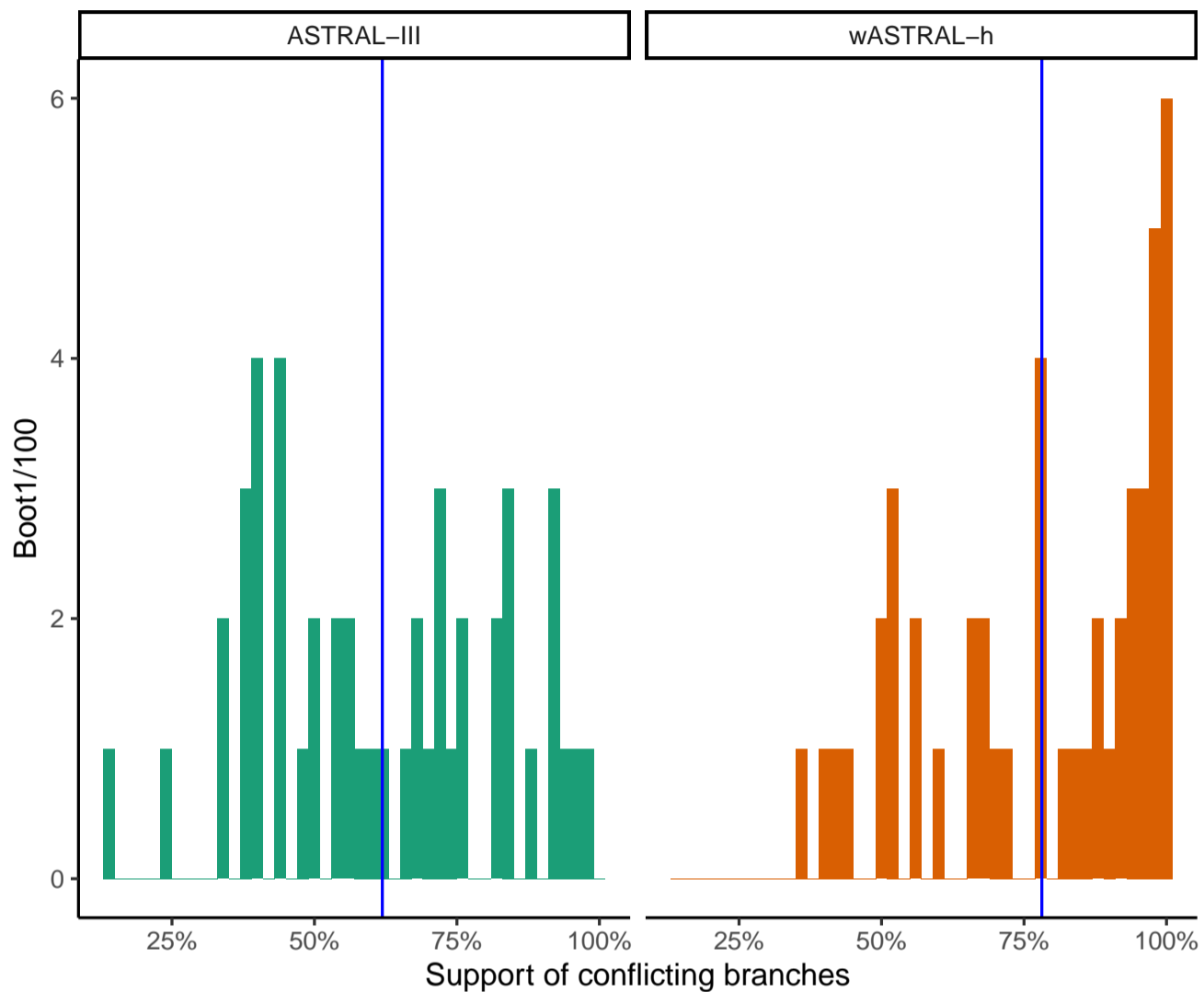


FIG. S13. The distribution of support values of conflicting branches between wASTRAL-h and ASTRAL-III on the 1kp dataset. The ASTRAL-III conflicting branches range between 14% and 99.00% with a mean of 62%. The wASTRAL-h conflicting branches range between 37% and 99.98% with a mean of 78%.

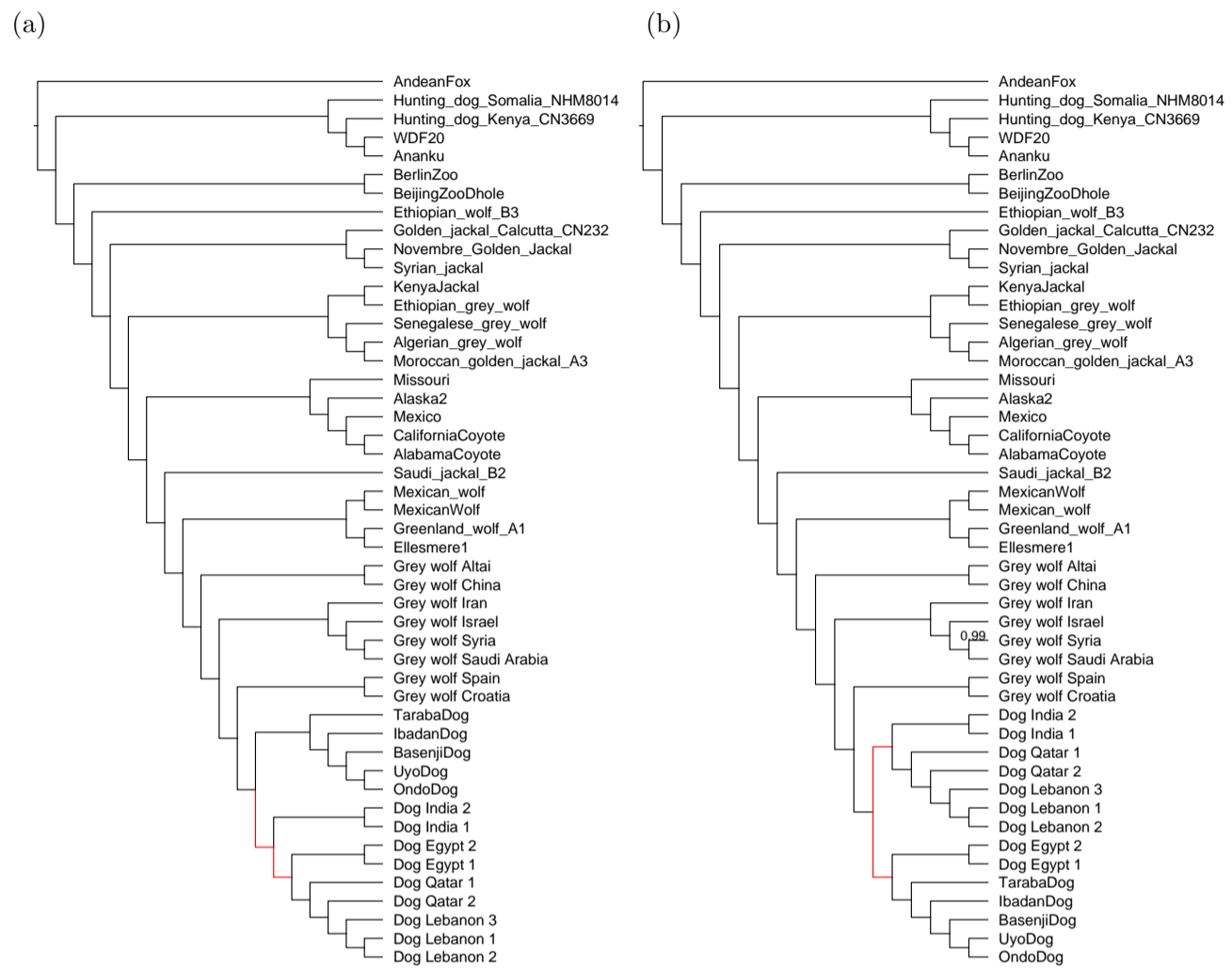


FIG. S14. Inferred species trees (a) from wASTRAL-hybrid with FastTree-2 branch support values as weights using all 459,450 gene trees and (b) from ASTRAL-III using a subset of 100,000 gene trees on canis dataset. Branches support of 100% are omitted.

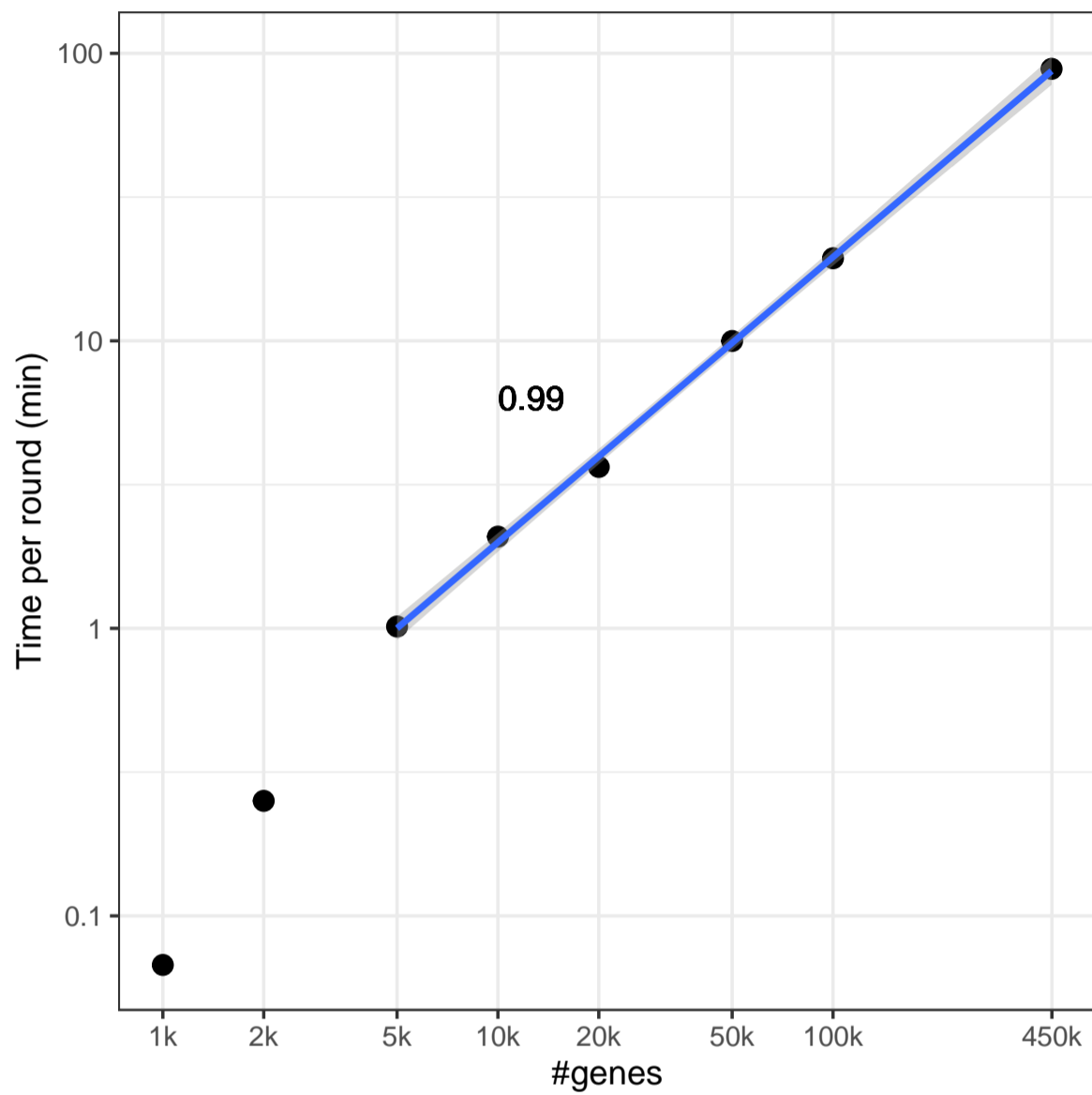
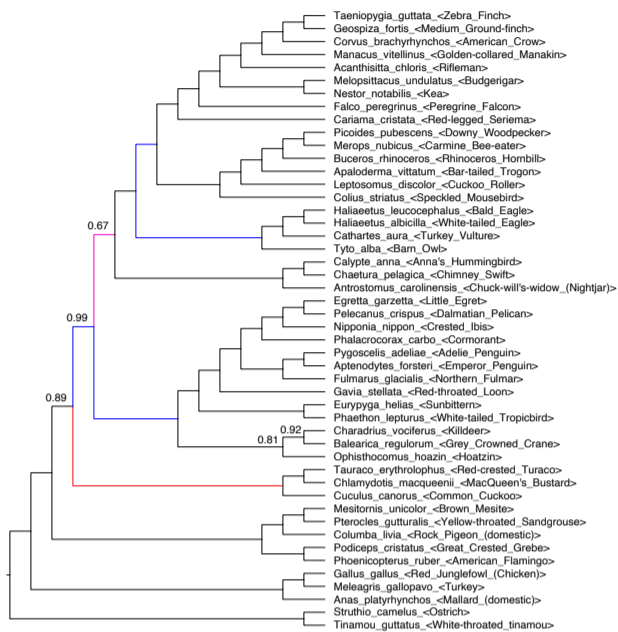


FIG. S15. Normalized time per round of placement by dividing running time by the total number of rounds of placements for wASTRAL-h on the Canis dataset for various k using the new pipeline.

(a)



(b)

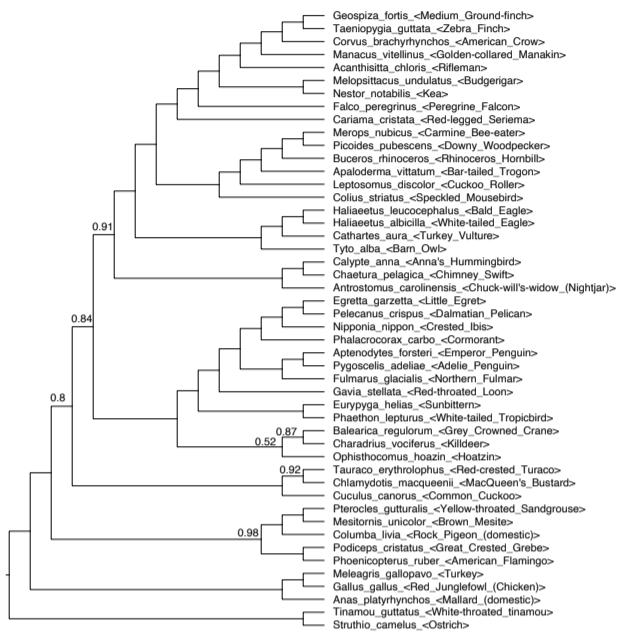
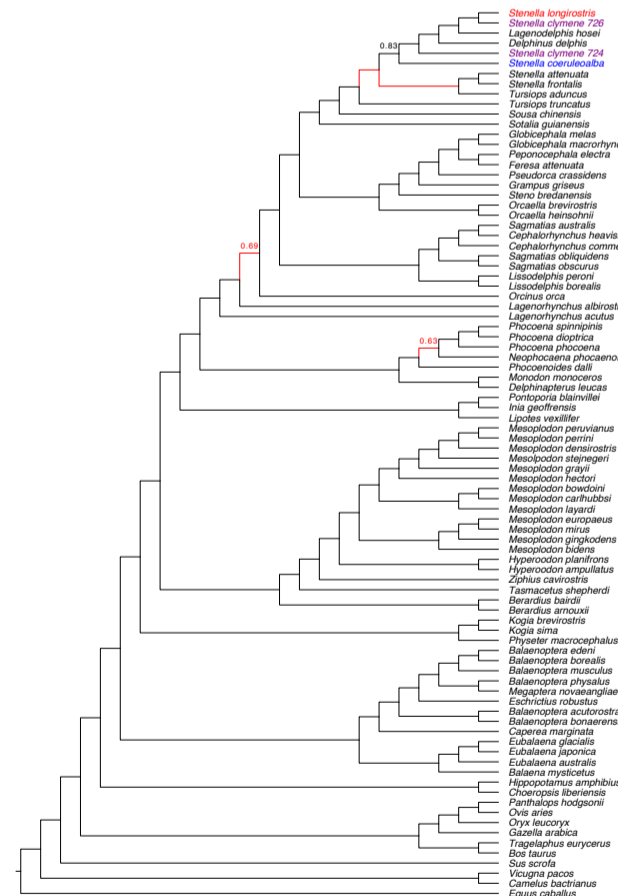


FIG. S16. Inferred species trees from (a) wASTRAL-hybrid with normalized bootstrap support values as weights and (b) ASTRAL-III on gene trees with contracted branches on avian dataset. Branches support of 100% are omitted. Branches that disagree with concatenation (blue), MP-EST binned (red) or both (purple) are identified on the wASTRAL-h tree.

(a)



(b)

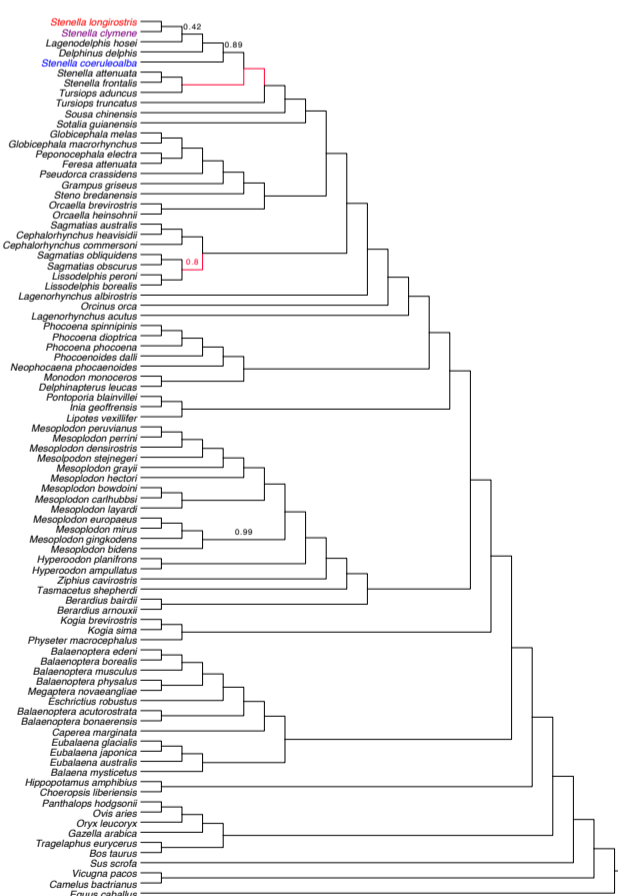


FIG. S17. Inferred species trees from (a) wASTRAL-hybrid with normalized Bayesian support values as weights (with clades of taxa from the same species contracted) and (b) ASTRAL-multi on cetacean dataset. Branches support of 100% are omitted. Branches conflicting with RAxML concatenation are marked red.

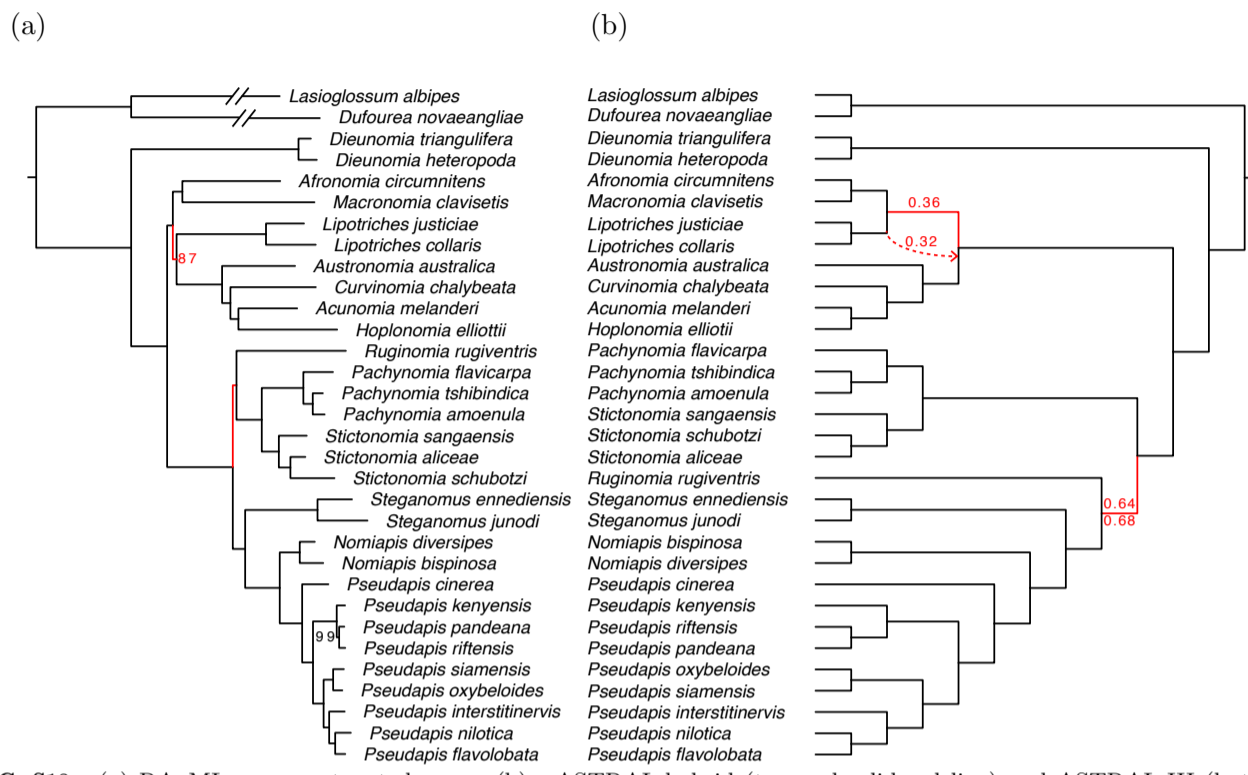
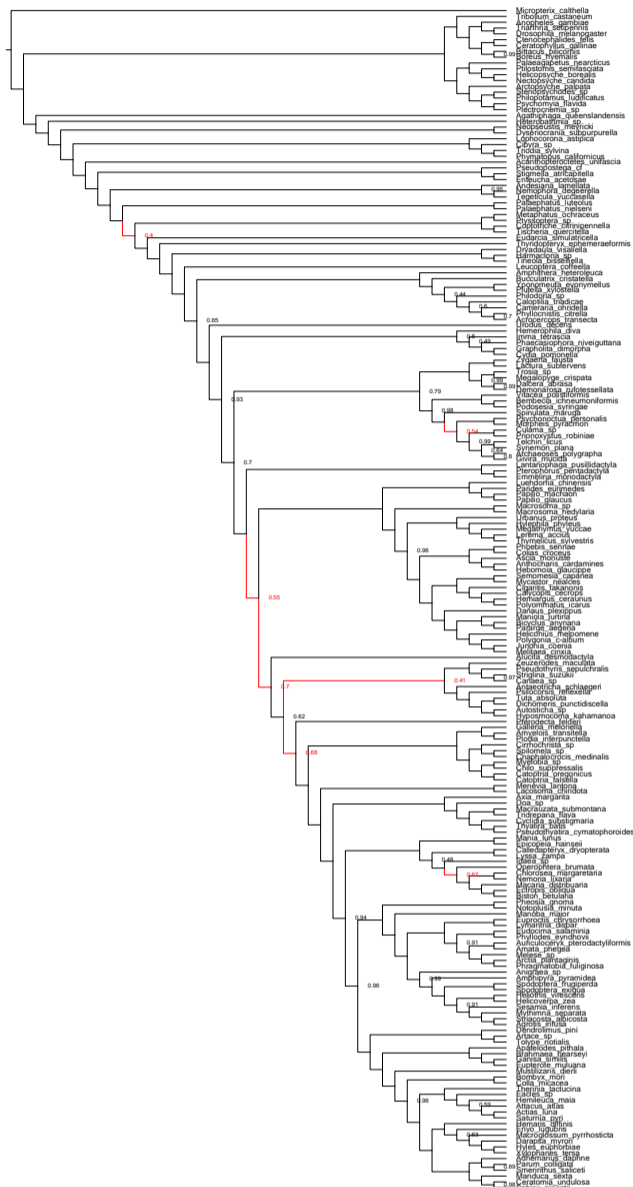


FIG. S18. (a) RAxML on concatenated genes; (b) wASTRAL-hybrid (top and solid red line) and ASTRAL-III (bottom and dashed red line) on Nomiinae dataset.

(a)



(b)

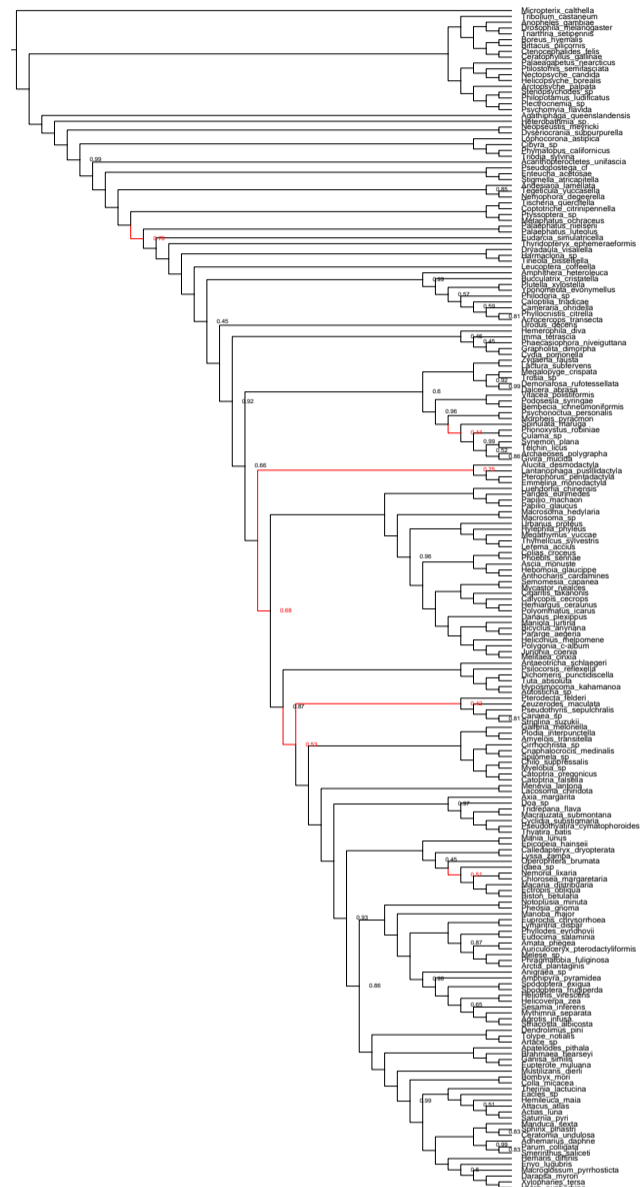


FIG. S19. Inferred species trees from (a) wASTRAL-hybrid with normalized bootstrap support values as weights and (b) ASTRAL-III on Lepidoptera dataset.

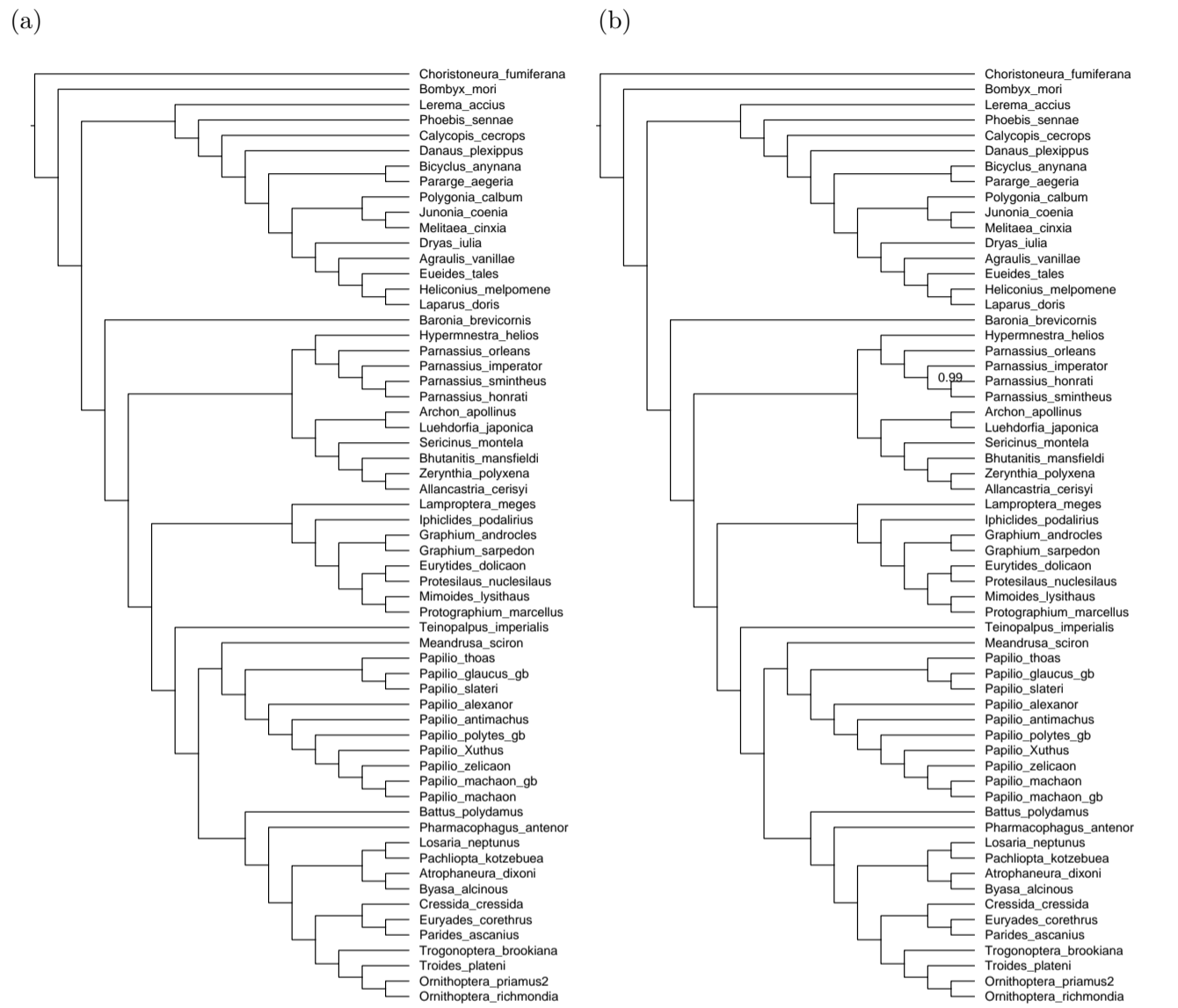


FIG. S20. Inferred species trees from (a) wASTRAL-hybrid with normalized approximate Bayesian support values as weights and (b) ASTRAL-III on Papilionidae dataset.

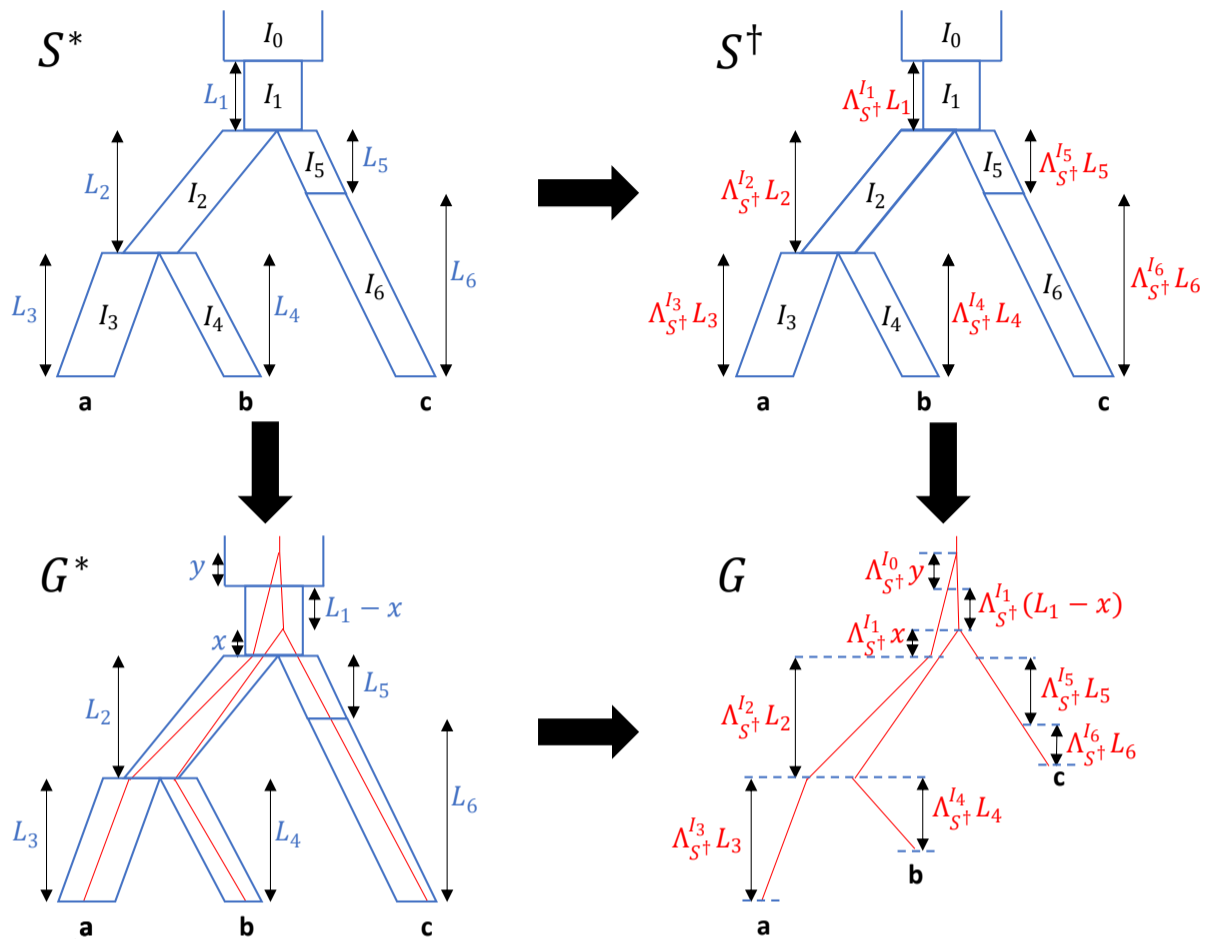


FIG. S21. An illustration of the process of creating a random gene tree with branch lengths in SU. Branches in the true species tree S^* are broken into intervals $I_0 \dots I_6$. The species tree with SU branch lengths S^\dagger is created by multiplying each branch length in S^* with a corresponding multiplier; the multipliers are jointly drawn from some distribution and are drawn independently across gene trees. Gene tree G^* is sampled under MSC process from S^* independent of S^\dagger . However, it inherits the same division of its lineages into segments as S^* at the same locations. The gene tree with SU branch lengths G is created by translating branch lengths of G^* into SU by multiplying the CU length of each of segment I_i by $\Lambda_{S^\dagger}^{I_i}$, the multiplier associated with the segment I_i in S^\dagger and hence G .

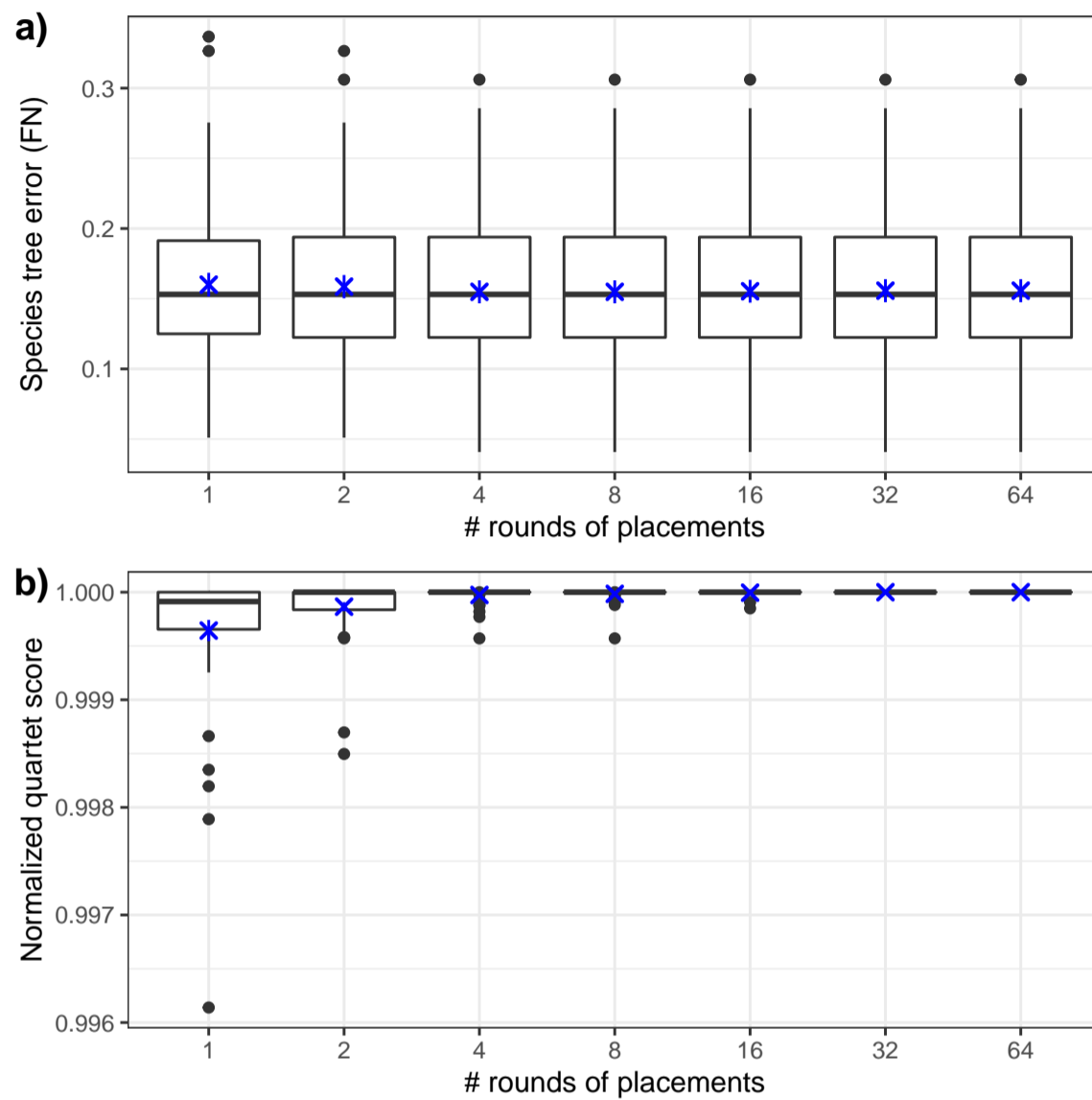


FIG. S22. The species tree estimation error (FN) of wASTRAL-h on S100 dataset as we change the number of rounds of placements in the base algorithm (r). The most difficult case where gene length=200 and $k=50$ is selected. Mean and standard error (50 replicates) are shown in blue.

Supplementary Algorithm

Algorithm S1 Recursive placement algorithm. **PLACE** inserts the species i into an existing species tree S and computes tripartition scores $W(A|B|C, \mathcal{G}) := \sum_{G \in \mathcal{G}} W(A \cap \mathcal{L}_G | B \cap \mathcal{L}_G | C \cap \mathcal{L}_G, G)$ for all tripartitions resulting from adding i onto each branch of S . A global counter Q and a set of per-node counters $w_A, w_B, w_C, w_{\cdot}^+, w_{\cdot}^-, w_{\cdot|\cdot}, w_{\cdot|\cdot}$ are all initialized to 0. **OptimalTreeDP** is defined in Algorithm S2. Each gene tree is rooted on an arbitrary branch e and the support of e is kept for the branch on one side of the root and zero support is given to the branch on the other side of root. \mathcal{L}_v is the set of leaves under v .

```

1: procedure PLACE( $i, S, \mathcal{G}$ ) ▷ Places species  $i$  on tree  $S$  according to  $\mathcal{G}$ 
2:    $W \leftarrow$  empty lookup table ▷ global variables
3:   COLORLEAFSET( $\mathcal{L}_S, C, \emptyset, \mathcal{G}, W$ ) ▷ Color all leaves of  $S$  as  $C$ 
4:   COLORLEAFSET( $\{i\}, B, \emptyset, \mathcal{G}, W$ ) ▷ Color new species  $i$  as  $B$ 
5:   COLORNODE(the root of  $S, i, S, \mathcal{G}, W$ ) ▷ Traverse  $S$  bottom up
6:    $O \leftarrow$  OPTIMALTREEDP( $\mathcal{L}_S \cup \{i\}, \mathcal{L}_S \cup \{i\}, W$ )
7:   return ( $W, O$ , edge of  $S$  onto which  $i$  is added to get  $O$ )
8: procedure COLORLEAFSET( $\mathcal{L}^*, X, T, \mathcal{G}, W$ ) ▷ Condition: Coloring  $\mathcal{L}^*$  as  $X$  should match  $T$ 
9:   for  $G \in \mathcal{G}$  do
10:    for  $j \in \mathcal{L}^* \cap \mathcal{L}_G$  do
11:       $W[T] \leftarrow$  UPDATECOUNTERS(leaf node corresponding to  $j$  in  $g, X$ )
12: procedure COLORNODE( $w, i, S, \mathcal{G}, W$ ) ▷ On start:  $i$  is  $B$ , others are  $C$ ; On exit:  $w$  is  $A$ , others kept
13:   if  $w$  is a leaf then
14:     COLORLEAFSET( $\mathcal{L}_w, A, \mathcal{L}_w | \{i\} | \mathcal{L}_S - \mathcal{L}_w, \mathcal{G}, W$ )
15:   else
16:     ( $u, v$ ) := ( the larger child of  $w$ , the smaller child of  $w$  )
17:     COLORNODE( $v, i, S, \mathcal{G}, W$ ) ▷ Recurse on  $v$ , the smaller child
18:     COLORLEAFSET( $\mathcal{L}_v, C, \emptyset, \mathcal{G}, W$ ) ▷ Undo coloring of  $v$  to enable recursing on  $u$ .
19:     COLORNODE( $u, i, S, \mathcal{G}, W$ ) ▷ Recurse on  $u$ , the large child
20:     COLORLEAFSET( $\mathcal{L}_v, B, \mathcal{L}_u | \{i\} \cup \mathcal{L}_v | \mathcal{L}_S - \mathcal{L}_w, \mathcal{G}, W$ ) ▷ Tripartition of  $w$  when adding  $i$  above  $v$ 
21:     COLORLEAFSET( $\{i\}, A, \{i\} \cup \mathcal{L}_u | \mathcal{L}_v | \mathcal{L}_S - \mathcal{L}_w, \mathcal{G}, W$ ) ▷ Tripartition of  $w$  when adding  $i$  above  $u$ 
22:     COLORLEAFSET( $\{i\}, C, \mathcal{L}_u | \mathcal{L}_v | \{i\} \cup \mathcal{L}_S - \mathcal{L}_w, \mathcal{G}, W$ ) ▷ Tripartition of  $w$  when adding  $i$  above  $w$ 
23:     COLORLEAFSET( $\{i\}, B, \emptyset, \mathcal{G}, W$ )
24:     COLORLEAFSET( $\mathcal{L}_v, A, \mathcal{L}_w | \{i\} | \mathcal{L}_S - \mathcal{L}_w, \mathcal{G}, W$ ) ▷ Tripartition of the new parent of  $i$  and  $w$ 
25: procedure RECURSIVEUPDATE( $w$ )
26:   ( $u, v, e$ ) := ( the left child of  $w$ , the right child of  $w$ , the parent branch of  $w$  )
27:   for  $(X, Y, Z) \in \{(A, B, C), (B, C, A), (C, A, B)\}$  do
28:      $Q \leftarrow Q - w_{XX|YZ}$ 
29:      $w_{XX|YZ} \leftarrow v_X u_{YZ|X} + u_X v_{YZ|X} + u_{XX|Z} v_Y + v_{XX|Z} u_Y + u_{XX|Y} v_Z + v_{XX|Y} u_Z$ 
30:        $+ (u_{YZ}^+ v_{XX}^+ - u_{YZ}^- v_{XX}^-) + (u_{XX}^+ v_{YZ}^+ - u_{XX}^- v_{YZ}^-)$ 
31:      $Q \leftarrow Q + w_{XX|YZ}$ 
32:     if  $w$  is not the root then
33:       ( $w_X, w_Y, w_Z$ )  $\leftarrow ((u_X + v_X)e^{-l(e)}, (u_Y + v_Y)e^{-l(e)}, (u_Z + v_Z)e^{-l(e)})$ 
34:        $w_{XX}^+ \leftarrow u_{XX}^+ + v_{XX}^+ + u_X v_X$ 
35:        $w_{XX}^- \leftarrow (u_{XX}^- + v_{XX}^- + u_X v_X)(1 - s(e))$ 
36:        $w_{YZ}^+ \leftarrow u_{YZ}^+ + v_{YZ}^+ + u_Y v_Z + u_Z v_Y$ 
37:        $w_{YZ}^- \leftarrow (u_{YZ}^- + v_{YZ}^- + u_Y v_Z + u_Z v_Y)(1 - s(e))$ 
38:        $w_{YZ|X} \leftarrow (u_{YZ|X} + v_{YZ|X} + (u_{YZ}^+ - u_{YZ}^-)v_X + u_X(v_{YZ}^+ - v_{YZ}^-))e^{-l(e)}$ 
39:        $w_{XX|Y} \leftarrow (u_{XX|Y} + v_{XX|Y} + (u_{XX}^+ - u_{XX}^-)v_Y + u_Y(v_{XX}^+ - v_{XX}^-))e^{-l(e)}$ 
40:        $w_{XX|Z} \leftarrow (u_{XX|Z} + v_{XX|Z} + (u_{XX}^+ - u_{XX}^-)v_Z + u_Z(v_{XX}^+ - v_{XX}^-))e^{-l(e)}$ 
41:       RECURSIVEUPDATE(the parent of  $w$ )
42: procedure UPDATECOUNTERS( $w, X$ ) ▷  $w$  is a leaf,  $X$  is a color
43:    $e :=$  the parent branch of  $w$ 
44:   ( $w_A, w_B, w_C$ )  $\leftarrow (0, 0, 0)$ 
45:    $w_X \leftarrow e^{-l(e)}$ 
46:   RECURSIVEUPDATE(the parent of  $w$ )
47:   return  $Q$ 

```

Algorithm S2 The Algorithm S2 of $O(n^2kH\log n)$ running time. At start, the function is called as with $\mathcal{L}_S, \mathcal{G}, r$ as input.

```

1: procedure NAIVEPLACEMENT( $T, \mathcal{G}, r$ )
2:    $W^* \leftarrow$  empty lookup table from tripartitions to their weights
3:   for  $i \in \{1, \dots, r\}$  do
4:     shuffle  $T$ 
5:      $S_i \leftarrow$  tree with leaves  $T_1, T_2,$  and  $T_3$ 
6:     for  $j \in \{4, \dots, |T|\}$  do
7:        $W_i, S_i, e \leftarrow$  PLACE( $T_j, S_i, \mathcal{G}$ )
8:     Add all elements of  $W_i$  to  $W^*$ 
9:   return OPTIMALTREEDP( $T, T, W^*$ )
10: procedure OPTIMALTREEDP( $P, \mathcal{L}, W$ )
11:   if DPTree( $P$ ) available then
12:     return DPTree( $P$ )
13:   if  $|P| = 1$  then
14:     DPscore( $P$ )  $\leftarrow$  0
15:     DPTree( $P$ )  $\leftarrow$  Singleton rooted tree with leafset  $P$ 
16:   else
17:      $X \leftarrow -\infty$ 
18:     for  $A \in \{A : W[A|P - A|\mathcal{L} - P] \text{ has been computed}\}$  do
19:        $S_1 \leftarrow$  OPTIMALTREEDP( $A, \mathcal{L}, W$ )
20:        $S_2 \leftarrow$  OPTIMALTREEDP( $P - A, \mathcal{L}, W$ )
21:       if DPscore( $A$ ) + DPscore( $P - A$ ) +  $W[A|P - A|\mathcal{L} - P] > X$  then
22:          $X \leftarrow$  DPscore( $A$ ) + DPscore( $P - A$ ) +  $W[A|P - A|\mathcal{L} - P]$ 
23:         DPTree( $P$ )  $\leftarrow$  merge subtrees  $S_1$  and  $S_2$  at root
24:     DPscore( $P$ )  $\leftarrow$   $X$ 
25:   return DPTree( $P$ )

```

Algorithm S3 The DAC algorithm of $O(n^{1.5+\epsilon}k)$ running time given some assumptions. OptimalTreeDP and NaivePlacement are defined in Algorithm S2, and Place is defined in Algorithm S1. At start, the function is called as with $\mathcal{L}_S, \mathcal{G}, r$ as input.

```

1: procedure TWOSTEPPLACEMENT( $T, \mathcal{G}, r$ )
2:    $W^* \leftarrow$  empty lookup table from tripartitions to their weights
3:   for  $i \in \{1, \dots, r\}$  do
4:      $T_i \leftarrow$  a subsample of  $T$  by removing each element independently with probability  $1 - 1/\sqrt{|T|}$ 
5:      $S_i :=$  NAIVEPLACEMENT( $T_i, \mathcal{G}, \sqrt{|T|}$ )
6:     for  $e \in E_{S_i}$  do
7:        $C_e \leftarrow$  empty list
8:       for  $j \in T - T_i$  do
9:          $W, S_o, e \leftarrow$  PLACE( $j, S_i, \mathcal{G}$ )
10:        add  $T_j$  to  $C_e$ 
11:      $C_\emptyset \leftarrow$  empty list
12:      $S'_i \leftarrow S_i$ 
13:     for  $e \in$  branches of  $S_i$  do
14:        $S_e \leftarrow S_i$ 
15:       for  $j \in C_e$  do
16:          $W, S_o, e' \leftarrow$  PLACE( $j, S_e, \mathcal{G}$ )
17:         if  $e' \in S_i - \{e\}$  then
18:           add  $j$  to  $C_\emptyset$ 
19:         else
20:            $S_e \leftarrow S_o$ 
21:        $S'_i \leftarrow$  The merger of compatible trees  $S_e$  and  $S'_i$ 
22:     for  $j \in C_\emptyset$  do
23:        $W_i, S'_i, e \leftarrow$  PLACE( $j, S'_i, \mathcal{G}$ )
24:     if  $C_\emptyset = \emptyset$  then
25:        $W_i, S'_i, e \leftarrow$  PLACE( $\emptyset, S'_i, \mathcal{G}$ )
26:     Add all elements of  $W_i$  to  $W^*$ 
27:   return OPTIMALTREEDP( $T, T, W^*$ )

```

Proofs

Weighting by support: Proof of Proposition 1 and Theorem 1

For ease of reference, we reproduce Table 2 from the main paper here:

$\mathbb{E}[(\cdot)(\cdot) \alpha_{G,Q}]$	$\delta_G(ab cd)$	$\delta_G(ac bd)$
$\delta_{G^*}(ab cd)$	$\geq \frac{1}{3}(1+2\theta_Q)(\alpha_{G,Q} + \frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q))$	$\leq \frac{1}{3}(1+2\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q))$
$\delta_{G^*}(ac bd)$	$\geq \frac{1}{3}(1-\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q))$	$\leq \frac{1}{3}(1-\theta_Q)(\alpha_{G,Q} + \frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q))$
$\delta_{G^*}(ad bc)$	$\geq \frac{1}{3}(1-\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q))$	$\leq \frac{1}{3}(1-\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q))$
$\mathbb{E}[(\cdot)(\cdot) \alpha_{G,Q}]$	$w_G(ab cd)$	$w_G(ac bd)$
$\delta_{G^*}(ab cd)$	$\geq \frac{1}{3}(1+2\theta_Q)(\alpha_{G,Q} + \frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q))^2$	$\leq \frac{1}{3}(1+2\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q))^2$
$\delta_{G^*}(ac bd)$	$\geq \frac{1}{3}(1-\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q))^2$	$\leq \frac{1}{3}(1-\theta_Q)(\alpha_{G,Q} + \frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q))^2$
$\delta_{G^*}(ad bc)$	$\geq \frac{1}{3}(1-\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q))^2$	$\leq \frac{1}{3}(1-\theta_Q)(\frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q))^2$

Recall that the expected value and variance of $\alpha_{G,Q}$ across genes is denoted by $\bar{\alpha}_Q$ and σ_α^2 .

PROPOSITION 1. For each estimated gene tree G , $\mathbb{E}[\delta_G(ab|cd) - \delta_G(ac|bd)] \geq \theta_Q \bar{\alpha}_Q - \frac{2}{3}(1 - \bar{\alpha}_Q)\beta_Q$ and $\mathbb{E}[w_G(ab|cd) - w_G(ac|bd)] \geq \frac{1}{9}\theta_Q(3 + 2\beta_Q)(\bar{\alpha}_Q^2 + \sigma_\alpha^2) + \frac{2}{9}(3 - \beta_Q)\theta_Q \bar{\alpha}_Q - \frac{4}{9}(1 - \bar{\alpha}_Q)\beta_Q$.

Proof. To prove the Proposition, we start with the following lemma.

LEMMA 1. For each estimated gene tree G with a given $\alpha_{G,Q}$,

$$\mathbb{E}[\delta_G(ab|cd) - \delta_G(ac|bd)|\alpha_{G,Q}] \geq \theta_Q \alpha_{G,Q} - \frac{2}{3}(1 - \alpha_{G,Q})\beta_Q$$

and

$$\mathbb{E}[w_G(ab|cd) - w_G(ac|bd)|\alpha_{G,Q}] \geq \frac{1}{9}(3\alpha_{G,Q} - 2\beta_Q + 2\alpha_{G,Q}\beta_Q + 6)\theta_Q \alpha_{G,Q} - \frac{4}{9}(1 - \alpha_{G,Q})\beta_Q.$$

Proof. From Table 2, we can compute

$$\begin{aligned} & \mathbb{E}[\delta_G(ab|cd) - \delta_G(ac|bd)|\alpha_{G,Q}] \\ &= \mathbb{E}\left[(\delta_G(ab|cd) - \delta_G(ac|bd))(\delta_{G^*}(ab|cd) + \delta_{G^*}(ac|bd) + \delta_{G^*}(ad|bc))\right|\alpha_{G,Q}] \\ &\geq \frac{1}{3}\left((1+2\theta_Q)\alpha_{G,Q} + \frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q)\right) - \frac{1}{3}\left((1-\theta_Q)\alpha_{G,Q} + \frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q)\right) \\ &= \theta_Q \alpha_{G,Q} - \frac{2}{3}(1 - \alpha_{G,Q})\beta_Q; \end{aligned}$$

similarly,

$$\begin{aligned} & \mathbb{E}[w_G(ab|cd) - w_G(ac|bd)|\alpha_{G,Q}] \\ &= \mathbb{E}\left[(w_G(ab|cd) - w_G(ac|bd))(\delta_{G^*}(ab|cd) + \delta_{G^*}(ac|bd) + \delta_{G^*}(ad|bc))\right|\alpha_{G,Q}] \\ &\geq \frac{1}{3}(1+2\theta_Q)\alpha_{G,Q}\left(\alpha_{G,Q} + \frac{2}{3}(1-\alpha_{G,Q})(1-\beta_Q)\right) + \left(\frac{1}{3}(1-\alpha_{G,Q})(1-\beta_Q)\right)^2 \\ &\quad - \frac{1}{3}(1-\theta_Q)\alpha_{G,Q}\left(\alpha_{G,Q} + \frac{2}{3}(1-\alpha_{G,Q})(1+\beta_Q)\right) - \left(\frac{1}{3}(1-\alpha_{G,Q})(1+\beta_Q)\right)^2 \\ &\geq \frac{1}{9}(3\alpha_{G,Q} - 2\beta_Q + 2\alpha_{G,Q}\beta_Q + 6)\theta_Q \alpha_{G,Q} - \frac{4}{9}(1 - \alpha_{G,Q})\beta_Q. \end{aligned}$$

□

From this lemma, we can prove the proposition. First, assume $\alpha_{G,Q}$ is drawn from a discrete distribution. Then,

$$\begin{aligned}\mathbb{E}[\delta_G(ab|cd) - \delta_G(ac|bd)] &= \sum_{\alpha_{G,Q}} \mathbb{E}[\delta_G(ab|cd) - \delta_G(ac|bd) | \alpha_{G,Q}] \mathbb{P}(\alpha_{G,Q}) \\ &\geq \sum_{\alpha_{G,Q}} (\theta_Q \alpha_{G,Q} - \frac{2}{3}(1 - \alpha_{G,Q})\beta_Q) \mathbb{P}(\alpha_{G,Q}) = \theta_Q \bar{\alpha}_Q - \frac{2}{3}(1 - \bar{\alpha}_Q)\beta_Q\end{aligned}$$

and

$$\begin{aligned}\mathbb{E}[w_G(ab|cd) - w_G(ac|bd)] &= \sum_{\alpha_{G,Q}} \mathbb{E}[w_G(ab|cd) - w_G(ac|bd) | \alpha_{G,Q}] \mathbb{P}(\alpha_{G,Q}) \\ &\geq \sum_{\alpha_{G,Q}} \left(\frac{1}{9}(3\alpha_{G,Q} - 2\beta_Q + 2\alpha_{G,Q}\beta_Q + 6)\theta_Q \alpha_{G,Q} - \frac{4}{9}(1 - \alpha_{G,Q})\beta_Q \right) \mathbb{P}(\alpha_{G,Q}) \\ &= \frac{1}{9}\theta_Q(3 + 2\beta_Q)\mathbb{E}[\alpha_{G,Q}^2] + \frac{2}{9}(3 - \beta_Q)\theta_Q \bar{\alpha}_Q - \frac{4}{9}(1 - \bar{\alpha}_Q)\beta_Q \\ &= \frac{1}{9}\theta_Q(3 + 2\beta_Q)(\bar{\alpha}_Q^2 + \sigma_\alpha^2) + \frac{2}{9}(3 - \beta_Q)\theta_Q \bar{\alpha}_Q - \frac{4}{9}(1 - \bar{\alpha}_Q)\beta_Q.\end{aligned}$$

It is straightforward to change these calculations to use integral instead of sum and $\mathbb{P}(\alpha_{G,Q})$ to the PDF in the case that the distribution of $\alpha_{G,Q}$ is continuous. \square

THEOREM 1. *Given estimated gene trees furnished with support generated under MSC+Error+Support model, there exist conditions where (3) guarantee a statistically consistent estimator of S^* but (2) does not, and the reverse is not true.*

Proof. Recall that (1) states

$$W(S, \mathcal{G}) := \sum_{G \in \mathcal{G}} \sum_{Q \in \mathcal{Q}(S)} w_G(S \upharpoonright Q).$$

It means that in order to produce a statistically consistent estimator using (1), the following equation must be satisfied for the true species tree topology S^* and any species tree topology S :

$$\mathbb{E}[W(S^*, \mathcal{G}) - W(S, \mathcal{G})] = |\mathcal{G}| \sum_{Q \in \mathcal{Q}(S)} \mathbb{E}[w_G(S^* \upharpoonright Q) - w_G(S \upharpoonright Q)] \geq 0 \quad (9)$$

Notice that proving for any quartet $Q = \{a, b, c, d\}$ we have $\mathbb{E}[w_G(ab|cd) - w_G(ac|bd)] \geq 0$ and $\mathbb{E}[w_G(ab|cd) - w_G(ad|bc)] \geq 0$ where $S^* \upharpoonright Q = ab|cd$ is sufficient to prove (9); on the other hand, proving for any quartet $Q = \{a, b, c, d\}$ where the internal branch of $S^* \upharpoonright Q$ corresponds to only one branch in S^* , we have $\mathbb{E}[w_G(ab|cd) - w_G(ac|bd)] \geq 0$ and $\mathbb{E}[w_G(ab|cd) - w_G(ad|bc)] \geq 0$ where $S^* \upharpoonright Q = ab|cd$ is necessary to prove (9).

Thus, from Proposition 1, we have guaranteed statistical consistency for wASTRAL for support under

$$D = \bigcap_{Q \in \mathcal{Q}(S)} \{(\theta_Q, \bar{\alpha}_Q, \sigma_\alpha, \beta_Q) \in [0, 1]^4 : \frac{1}{9}\theta_Q(3 + 2\beta_Q)(\bar{\alpha}_Q^2 + \sigma_\alpha^2) + \frac{2}{9}(3 - \beta_Q)\theta_Q \bar{\alpha}_Q - \frac{4}{9}(1 - \bar{\alpha}_Q)\beta_Q \geq 0\}.$$

Similarly, we have guaranteed statistical consistency for unweighted ASTRAL under

$$D' = \bigcap_{Q \in \mathcal{Q}(S)} \{(\theta_Q, \bar{\alpha}_Q, \sigma_\alpha, \beta_Q) \in [0, 1]^4 : \theta_Q \bar{\alpha}_Q - \frac{2}{3}(1 - \bar{\alpha}_Q)\beta_Q \geq 0\}.$$

To prove Theorem 1, we only need to prove that D' is a proper subset of D .

We can prove $D' \subseteq D$, as for any Q , if $(\theta_Q, \bar{\alpha}_Q, \sigma_\alpha, \beta_Q) \in [0, 1]^4$ and $\theta_Q \bar{\alpha}_Q - \frac{2}{3}(1 - \bar{\alpha}_Q)\beta_Q \geq 0$, then

$$\begin{aligned} & \frac{1}{9}\theta_Q(3+2\beta_Q)(\bar{\alpha}_Q^2 + \sigma_\alpha^2) + \frac{2}{9}(3-\beta_Q)\theta_Q\bar{\alpha}_Q - \frac{4}{9}(1-\bar{\alpha}_Q)\beta_Q \\ &= \frac{1}{9}\theta_Q(3+2\beta_Q)\sigma_\alpha^2 + \frac{1}{3}\theta_Q(1-\theta_Q)\bar{\alpha}_Q^2 + \left(\frac{1}{3}\theta_Q\bar{\alpha}_Q + \frac{2}{3}\right)(\theta_Q\bar{\alpha}_Q - \frac{2}{3}(1-\bar{\alpha}_Q)\beta_Q) \geq 0. \end{aligned}$$

We can also prove $D' \neq D$, as if for some Q , $\theta_Q = 0.25, \bar{\alpha}_Q = 0.5, \beta_Q = 0.4$,

$$\theta_Q\bar{\alpha}_Q - \frac{2}{3}(1-\bar{\alpha}_Q)\beta_Q = -\frac{1}{120} < 0$$

and

$$\frac{1}{9}\theta_Q(3+2\beta_Q)(\bar{\alpha}_Q^2 + \sigma_\alpha^2) + \frac{2}{9}(3-\beta_Q)\theta_Q\bar{\alpha}_Q - \frac{4}{9}(1-\bar{\alpha}_Q)\beta_Q = \frac{7}{720} + \frac{19}{180}\sigma_\alpha^2 > 0.$$

Thus D' is a proper subset of D and we conclude the proof.

□

Weighting by length: Proof of Propositions 2 and 3 and Theorem 2

Before providing the proofs, we remind the reader of one property of the coalescent model. According to the coalescent model, at any point along a branch of the species tree with i gene tree lineages, the time (i.e., distance) x to the next coalescent event, reducing the number of lineages to $i-1$, is exponentially distributed with the rate $\binom{i}{2}$, resulting in probability density function (PDF):

$$\frac{i(i-1)}{2} e^{-\frac{i(i-1)}{2}x}, \quad (10)$$

and the two lineages that coalesce are independent of x .

PROPOSITION 2. For a true quartet species tree S^* with topology $ab|cd$ and input gene trees \mathcal{G} generated under the naive model with any multiplier λ , let f be the distance between anchors of S^* . As $f \rightarrow 0$, given $k = \Theta(f^{-2})$ gene trees, we have $\text{Var}[X_G] = \Theta_f(1)$ and

$$\frac{\mathbb{E}[X_G]}{\sqrt{\text{Var}[X_G]}} = \frac{1+4\lambda}{1+2\lambda} \sqrt{\frac{3}{2}} f + O(f^2).$$

Proof. We analyze balanced and unbalanced trees separately.

Case 1: Unbalanced trees (i.e., the root of S^* has a terminal branch as a child). W.o.l.g., we assume the root branch is located on branch leading to d .

Let p, q , and r be the MRCA nodes of (a, b) , (a, c) , and (a, d) on rooted species tree S^* , respectively. Let p' and r' be the points of coalescence of leaves a, b and leaves c, d on the rooted gene tree G , respectively. Let x, y_0 , and z be the CU difference in heights of points (p, p') , (q, r) , and (r, r') , respectively. Note that f is the length of (p, q) . Let $L := l_{S^*}(a, p) + l_{S^*}(b, p) + l_{S^*}(c, r) + l_{S^*}(d, r)$. Notice that $l_G(a, p) + l_G(b, p) + l_G(c, r) + l_G(d, r) = \lambda L$ and $l_G(a, b) + l_G(c, d) = \lambda(2x + 2z + L)$.

Let $f_X(x)$ be the probability density that x is the CU difference in heights of (p, p') and p' is the lowest point of coalescence. Notice that by (10):

$$f_X(x) = \begin{cases} e^{-x} & 0 \leq x \leq f \\ \frac{1}{\binom{2}{3}} \left(e^{-f} \binom{2}{3} e^{-\binom{2}{3}(x-f)} \right) = e^{-3x+2f} & f \leq x \leq f + y_0 \\ \frac{1}{\binom{2}{4}} \left(e^{-f} e^{-\binom{2}{3}y_0} \binom{2}{4} e^{-\binom{2}{4}(x-f-y_0)} \right) = e^{-6x+5f+3y_0} & f + y_0 \leq x \end{cases}$$

Let $f_{Z|X}(z; x)$ be the probability density that z is the CU difference in heights of (r, r') , conditioned on that x is the CU difference in heights of (p, p') and p' is the lowest point of coalescence. Notice that:

$$f_{Z|X}(z; x) = \begin{cases} e^{-z} & 0 \leq x \leq f + y_0 \text{ and } 0 \leq z \\ e^{-(z-(x-f-y_0))} = e^{-z+x-f-y_0} & 0 \leq x - f - y_0 \leq z \end{cases}$$

We specify three coalescence scenarios by indicator functions $\delta_1, \delta_2, \delta_3$: *i*) δ_1 indicates $0 \leq x < f$; *ii*) δ_2 indicates $f \leq x < f + y_0$; *iii*) δ_3 indicates $f + y_0 \leq x$.

Note that

$$\begin{aligned} \mathbb{E}[w_G(ab|cd)] &= \mathbb{E}[(\delta_1 + \delta_2 + \delta_3)w_G(ab|cd)] \\ \mathbb{E}[w_G^2(ab|cd)] &= \mathbb{E}[(\delta_1 + \delta_2 + \delta_3)w_G^2(ab|cd)]. \end{aligned}$$

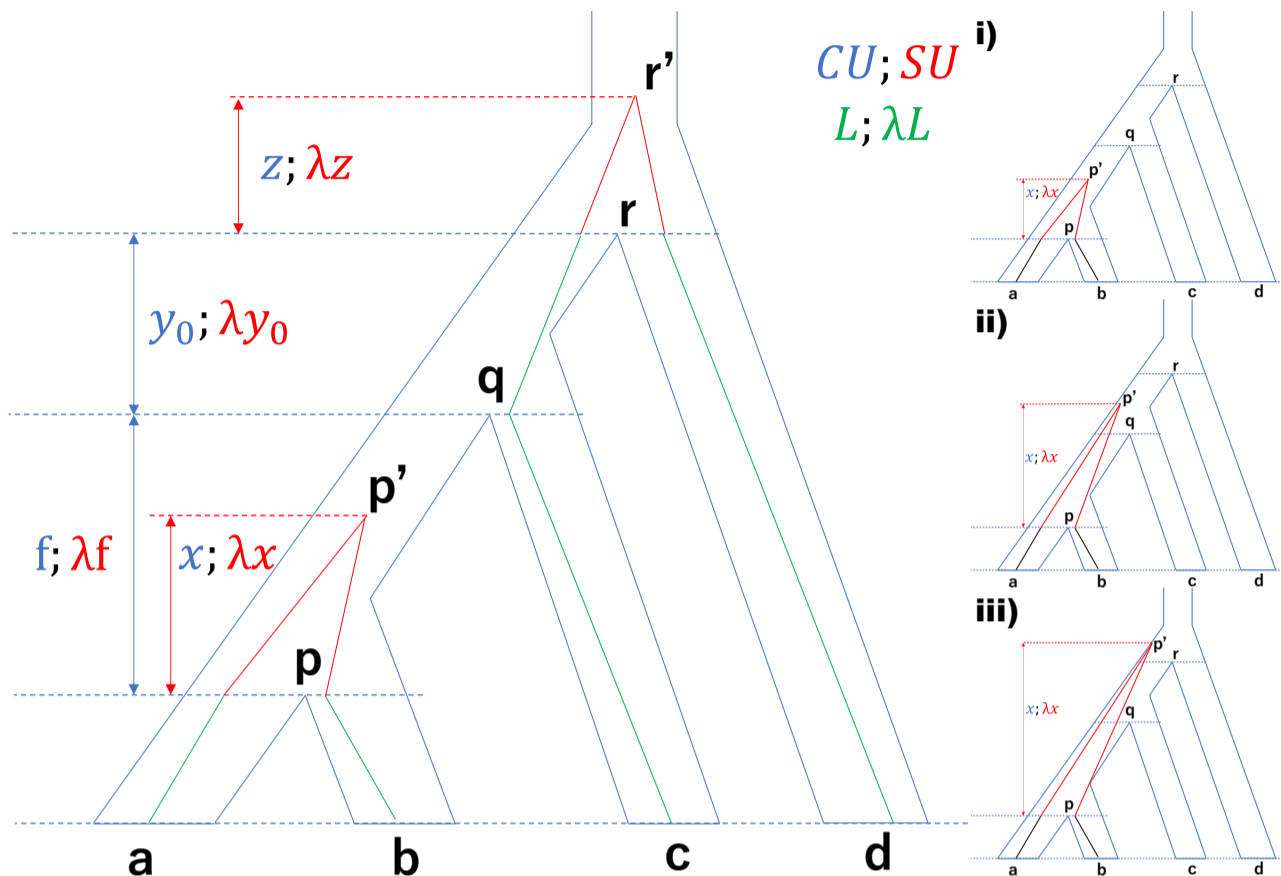


FIG. S23. Illustration of the unbalanced case. Lengths in CU/SU units are denoted in blue/red. Branches in green have a total length $L/\lambda L$ in CU/SU units. The right-hand side shows the position of p' in relation to q and r in various cases.

Similarly, since only scenarios 2 and 3 have deep coalescence events that may lead to gene tree disagreement with the species tree, and by the symmetry of all three topologies under scenarios 2 and 3,

$$\begin{aligned}\mathbb{E}[w_G(ac|bd)] &= \mathbb{E}[(\delta_2 + \delta_3)w_G(ab|cd)] \\ \mathbb{E}[w_G^2(ac|bd)] &= \mathbb{E}[(\delta_2 + \delta_3)w_G^2(ab|cd)].\end{aligned}$$

Thus,

$$\mathbb{E}[X_G] = \mathbb{E}[w_G(ab|cd)] - \mathbb{E}[w_G(ac|bd)] = \mathbb{E}[\delta_1 w_G(ab|cd)], \quad (11)$$

and since $w_G(ab|cd)w_G(ac|bd) = 0$,

$$\begin{aligned}\text{Var}[X_G] &= \mathbb{E}[X_G^2] - \mathbb{E}^2[X_G] = \mathbb{E}[w_G^2(ab|cd) + w_G^2(ac|bd)] - \mathbb{E}^2[X_G] \\ &= \mathbb{E}[(\delta_1 + 2\delta_2 + 2\delta_3)w_G^2(ab|cd)] - \mathbb{E}^2[X_G].\end{aligned} \quad (12)$$

We next compute both elements of (11) as well as some elements of (12) (others will not be necessary).

- δ_1 : When G has topology $ab|cd$, p' must be the lowest point of coalescence. Thus,

$$\begin{aligned}\mathbb{E}[\delta_1 w_G(ab|cd)] &= \int_0^f \int_0^{+\infty} e^{-\lambda(2x+2z+L)} f_X(x) f_{Z|X}(z;x) dz dx \\ &= \int_0^f \int_0^{+\infty} e^{-\lambda(2x+2z+L)} e^{-x} e^{-z} dz dx \\ &= \frac{e^{-\lambda L} (1 - e^{-(1+2\lambda)f})}{(1+2\lambda)^2};\end{aligned}$$

$$\mathbb{E}[\delta_1 w_G^2(ab|cd)] \leq \mathbb{E}[\delta_1 w_G(ab|cd)] = O(f).$$

- δ_2 : When G has topology $ab|cd$, p' must be the lowest point of coalescence. Thus,

$$\begin{aligned} \mathbb{E}[\delta_2 w_G^2(ab|cd)] &= \int_f^{f+y_0} \int_0^{+\infty} e^{-\lambda(4x+4z+2L)} f_X(x) f_{Z|Y}(z;y) dz dx \\ &= \int_f^{f+y_0} \int_0^{+\infty} e^{-\lambda(4x+4z+2L)} e^{-3x+2f} e^{-z} dz dx \\ &= \frac{1 - e^{-(3+4\lambda)y_0}}{(1+4\lambda)(3+4\lambda)} e^{-(1+4\lambda)f - 2\lambda L}. \end{aligned}$$

- δ_3 : When G has the topology $ab|cd$, either p' or q' must be the lowest point of coalescence, and by symmetry, the two cases must have the same PDFs. Thus,

$$\begin{aligned} \mathbb{E}[\delta_3 w_G^2(ab|cd)] &= \int_{f+y_0}^{+\infty} \int_{x-f-y_0}^{+\infty} e^{-\lambda(4x+4z+2L)} 2f_X(x) f_{Z|X}(z;x) dz dx \\ &= \int_{f+y_0}^{+\infty} \int_{x-f-y_0}^{+\infty} e^{-\lambda(4x+4z+2L)} 2e^{-6x+5f+3y_0} e^{-z+x-f-y_0} dz dx \\ &= \int_{f+y_0}^{+\infty} e^{-4\lambda(x+x-f-y_0)-2\lambda L} 2e^{-6x+5f+3y_0} \frac{1}{1+4\lambda} dx \\ &= \frac{1}{(3+4\lambda)(1+4\lambda)} e^{-(1+4\lambda)f - (3+4\lambda)y_0 - 2\lambda L}. \end{aligned}$$

Replacing in (11), we get

$$\mathbb{E}[X_G] = \mathbb{E}[\delta_1 w_G(ab|cd)] = \frac{e^{-\lambda L} (1 - e^{-(1+2\lambda)f})}{(1+2\lambda)^2} = \frac{e^{-\lambda L}}{1+2\lambda} f + O(f^2);$$

and replacing in (12), we get

$$\begin{aligned} \text{Var}[X_G] &= \mathbb{E}[(\delta_1 + 2\delta_2 + 2\delta_3) w_G^2(ab|cd)] - \mathbb{E}^2[X_G] = \mathbb{E}[2(\delta_2 + \delta_3) w_G^2(ab|cd)] + O(f) \\ &= \frac{2e^{-(1+4\lambda)f - 2\lambda L}}{(3+4\lambda)(1+4\lambda)} + O(f) = \frac{2e^{-2\lambda L}}{(3+4\lambda)(1+4\lambda)} + O(f), \end{aligned}$$

from which our assumption of $\text{Var}[X_G] = \Omega(1)$ follows.

Case 2: Balanced tree.

Let p, q , and r be the MRCA nodes of (a, b) , (c, d) , and (a, d) on rooted species tree S^* , respectively. Let p' and q' be the points of coalescence of leaves a, b and leaves c, d on the rooted gene tree G , respectively. Let x, x_0, y , and y_0 be the CU difference in heights of points (p, p') , (p, r) , (q, q') , and (q, r) , respectively. Note that $f = x + y$ is CU length of path (p, q) . Let $L := l_{S^*}(a, p) + l_{S^*}(b, p) + l_{S^*}(c, q) + l_{S^*}(d, q)$. Notice that $l_G(a, p) + l_G(b, p) + l_G(c, q) + l_G(d, q) = \lambda L$ and $l_G(a, b) + l_G(c, d) = \lambda(2x + 2y + L)$.

We specify three coalescence scenarios by indicator functions $\delta_1, \delta_2, \delta_3$: *i*) δ_1 indicates $0 \leq x < x_0$; *ii*) δ_2 indicates $x_0 \leq x, 0 \leq y < y_0$; *iii*) δ_3 indicates $x_0 \leq x, y_0 \leq y$.

Note that

$$\begin{aligned} \mathbb{E}[w_G(ab|cd)] &= \mathbb{E}[(\delta_1 + \delta_2 + \delta_3) w_G(ab|cd)] \\ \mathbb{E}[w_G^2(ab|cd)] &= \mathbb{E}[(\delta_1 + \delta_2 + \delta_3) w_G^2(ab|cd)]. \end{aligned}$$

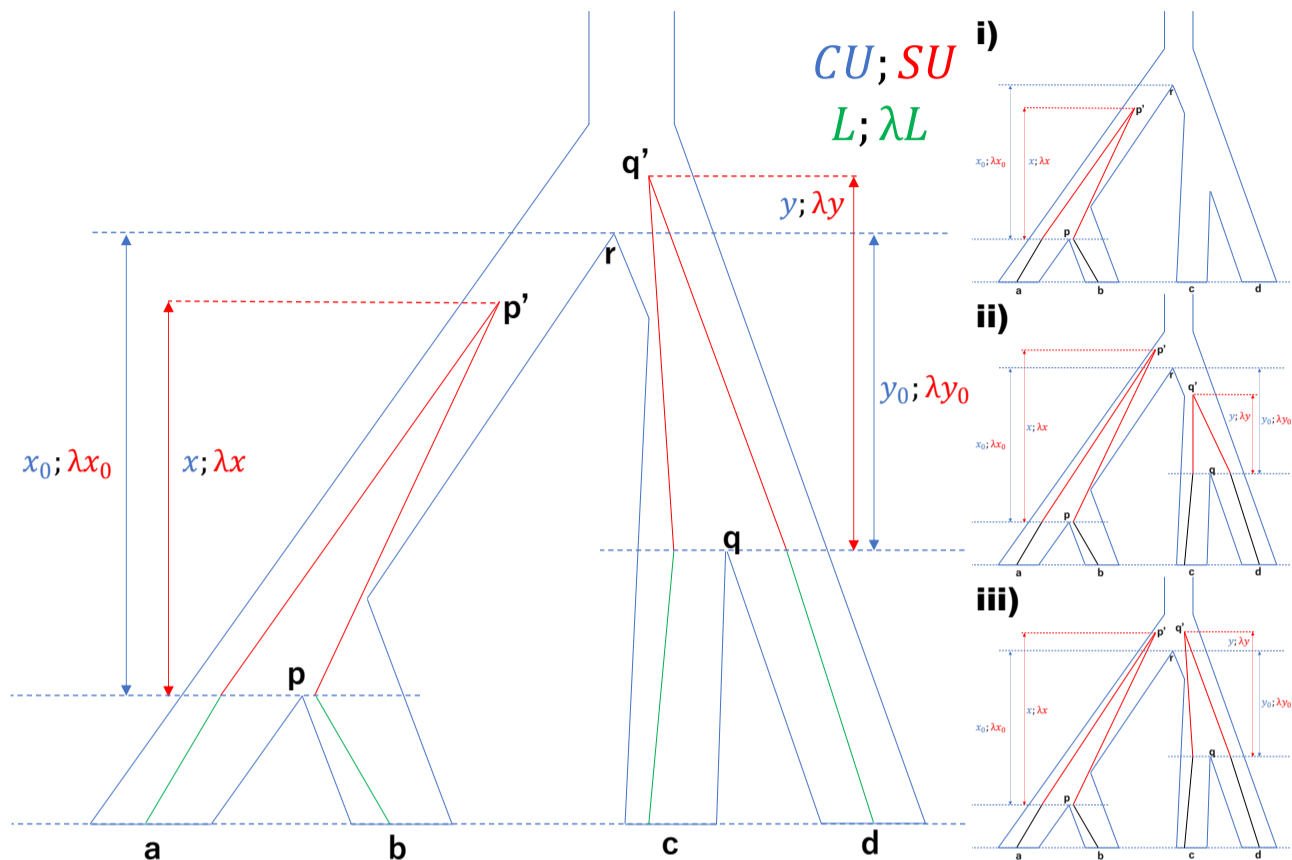


FIG. S24. Illustration of the unbalanced case. Lengths in CU/SU units are denoted in blue/red. Branches in green have a total length $L/\lambda L$ in CU/SU units. The right-hand side shows the position of p' and q' in relation to r in various cases.

Similarly, since only scenarios 3 have deep coalescence events that may lead to gene tree disagreement with the species tree, and by the symmetry of all three topologies under scenarios 3,

$$\begin{aligned}\mathbb{E}[w_G(ac|bd)] &= \mathbb{E}[\delta_3 w_G(ab|cd)] \\ \mathbb{E}[w_G^2(ac|bd)] &= \mathbb{E}[\delta_3 w_G^2(ab|cd)].\end{aligned}$$

Thus,

$$\mathbb{E}[X_G] = \mathbb{E}[w_G(ab|cd)] - \mathbb{E}[w_G(ac|bd)] = \mathbb{E}[(\delta_1 + \delta_2)w_G(ab|cd)]; \quad (13)$$

and since $w_G(ab|cd)w_G(ac|bd) = 0$,

$$\begin{aligned}\text{Var}[X_G] &= \mathbb{E}[X_G^2] - \mathbb{E}^2[X_G] = \mathbb{E}[w_G^2(ab|cd) + w_G^2(ac|bd)] - \mathbb{E}^2[X_G] \\ &= \mathbb{E}[(\delta_1 + \delta_2 + 2\delta_3)w_G^2(ab|cd)] - \mathbb{E}^2[X_G].\end{aligned} \quad (14)$$

- δ_1 : Here,

$$\begin{aligned}\mathbb{E}[\delta_1 w_G(ab|cd)] &= \int_0^{x_0} \int_0^{+\infty} e^{-\lambda(2x+2y+L)} e^{-x} e^{-y} dy dx \\ &= \frac{e^{-\lambda L}(1 - e^{-(1+2\lambda)x_0})}{(1+2\lambda)^2} = \frac{e^{-\lambda L} x_0}{1+2\lambda} + O(x_0^2) = \frac{e^{-\lambda L} x_0}{1+2\lambda} + O(f^2);\end{aligned}$$

and

$$\mathbb{E}[\delta_1 w_G^2(ab|cd)] \leq \mathbb{E}[\delta_1 w_G(ab|cd)] = O(f).$$

- δ_2 : Here,

$$\begin{aligned}\mathbb{E}[\delta_2 w_G(ab|cd)] &= \int_{x_0}^{+\infty} \int_0^{y_0} e^{-\lambda(2x+2y+L)} e^{-x} e^{-y} dy dx \\ &= \frac{e^{-\lambda L} (1 - e^{-(1+2\lambda)y_0}) e^{-(1+2\lambda)x_0}}{(1+2\lambda)^2} = \frac{e^{-\lambda L} y_0}{1+2\lambda} + O(f^2);\end{aligned}$$

and

$$\mathbb{E}[\delta_2 w_G^2(ab|cd)] \leq \mathbb{E}[\delta_2 w_G(ab|cd)] = O(f).$$

- δ_3 : Similar to the unbalanced case, when G has the topology $ab|cd$, either p' or q' must be the lowest point of coalescence, and by symmetry, the two cases must have the same PDFs. Thus,

$$\begin{aligned}\mathbb{E}[\delta_3 w_G^2(ab|cd)] &= \int_{x_0}^{+\infty} \int_{x-x_0+y_0}^{+\infty} e^{-\lambda(4x+4y+2L)} 2e^{-x_0} e^{-y_0} e^{-6x+6x_0} e^{-y+x-x_0+y_0} dy dx \\ &= \int_{x_0}^{+\infty} e^{-4\lambda(x+x-x_0+y_0)-2\lambda L} 2e^{-x_0} e^{-y_0} e^{-6x+6x_0} \frac{1}{1+4\lambda} dx \\ &= \frac{1}{(3+4\lambda)(1+4\lambda)} e^{-(1+4\lambda)(x_0+y_0)-2\lambda L} = \frac{1}{(3+4\lambda)(1+4\lambda)} e^{-(1+4\lambda)f-2\lambda L}.\end{aligned}$$

Replacing in (13), we get

$$\mathbb{E}[X_G] = \mathbb{E}[(\delta_1 + \delta_2)w_G(ab|cd)] = \frac{e^{-\lambda L}(x_0 + y_0)}{1+2\lambda} + O(f^2) = \frac{e^{-\lambda L}f}{1+2\lambda} + O(f^2);$$

and replacing in (14), we get

$$\begin{aligned}\text{Var}[X_G] &= \mathbb{E}[(\delta_1 + \delta_2 + 2\delta_3)w_G^2(ab|cd)] - \mathbb{E}^2[X_G] \\ &= \mathbb{E}[2\delta_3 w_G^2(ab|cd)] + O(f) \\ &= \frac{2e^{-(1+4\lambda)f-2\lambda L}}{(3+4\lambda)(1+4\lambda)} + O(f) = \frac{2e^{-2\lambda L}}{(3+4\lambda)(1+4\lambda)} + O(f),\end{aligned}$$

from which our assumption of $\text{Var}[X_G] = \Theta_f(1)$ follows.

Thus, in both balanced and unbalanced cases,

$$\frac{\mathbb{E}[X_G]}{\sqrt{\text{Var}[X_G]}} = \frac{\frac{e^{-\lambda L}f}{1+2\lambda} + O(f^2)}{\sqrt{\frac{2e^{-2\lambda L}}{(1+4\lambda)(3+4\lambda)} + O(f)}} = \sqrt{1 + \frac{4\lambda + 4\lambda^2}{3(1+2\lambda)^2}} \sqrt{\frac{3}{2}} f + O(f^2)$$

□

PROPOSITION 3. For a true quartet species tree S^* with topology $ab|cd$ and input gene trees \mathcal{G} generated under the variable rate model, let f be the distance between anchors of S^* and L be the total length of all other branches. Assume that for every branch segment I , the variance of its multiplier is bounded above: $\text{Var}(\Lambda_{S^*}^I) \leq \varepsilon^2$ where $\varepsilon^2 = \frac{e^{-\lambda L}}{(16+32\lambda)+(6+32\lambda+32\lambda^2)L} \left(\frac{20(\lambda+\lambda^2)}{9(1+2\lambda)^2}\right)^3$. As $f \rightarrow 0$, given $k = \Theta(f^{-2})$ gene trees, we have $\text{Var}[X_G] = \Theta_f(1)$ and

$$\frac{\mathbb{E}[X_G]}{\sqrt{\text{Var}[X_G]}} \geq \sqrt{\frac{3}{2}} \left(1 - \frac{4\lambda^2}{(1+4\lambda)^2}\right)^{-\frac{1}{2}} f + O(f^2).$$

Proof. We follow the same logic in proof of Proposition 2.

Case 1: Unbalanced trees. Let $P(x)$ be functions to random variables denoting SU difference in heights of points (p, p') where p' is x CU distance above p ; let $R(z)$ be functions to random variables denoting SU difference in heights of points (r, r') where r' is z CU distance above r . Note that $P(f+y_0)+R(z)=P(f+y_0+z)$ where $P(f+y_0)$ denote the SU length of (p, r) . Let random variable $\Lambda := (l_{S^\dagger}(a, p) + l_{S^\dagger}(b, p) + l_{S^\dagger}(c, r) + l_{S^\dagger}(d, r))$ be the total SU terminal branch lengths and the constant value L be the CU distance corresponding to Λ .

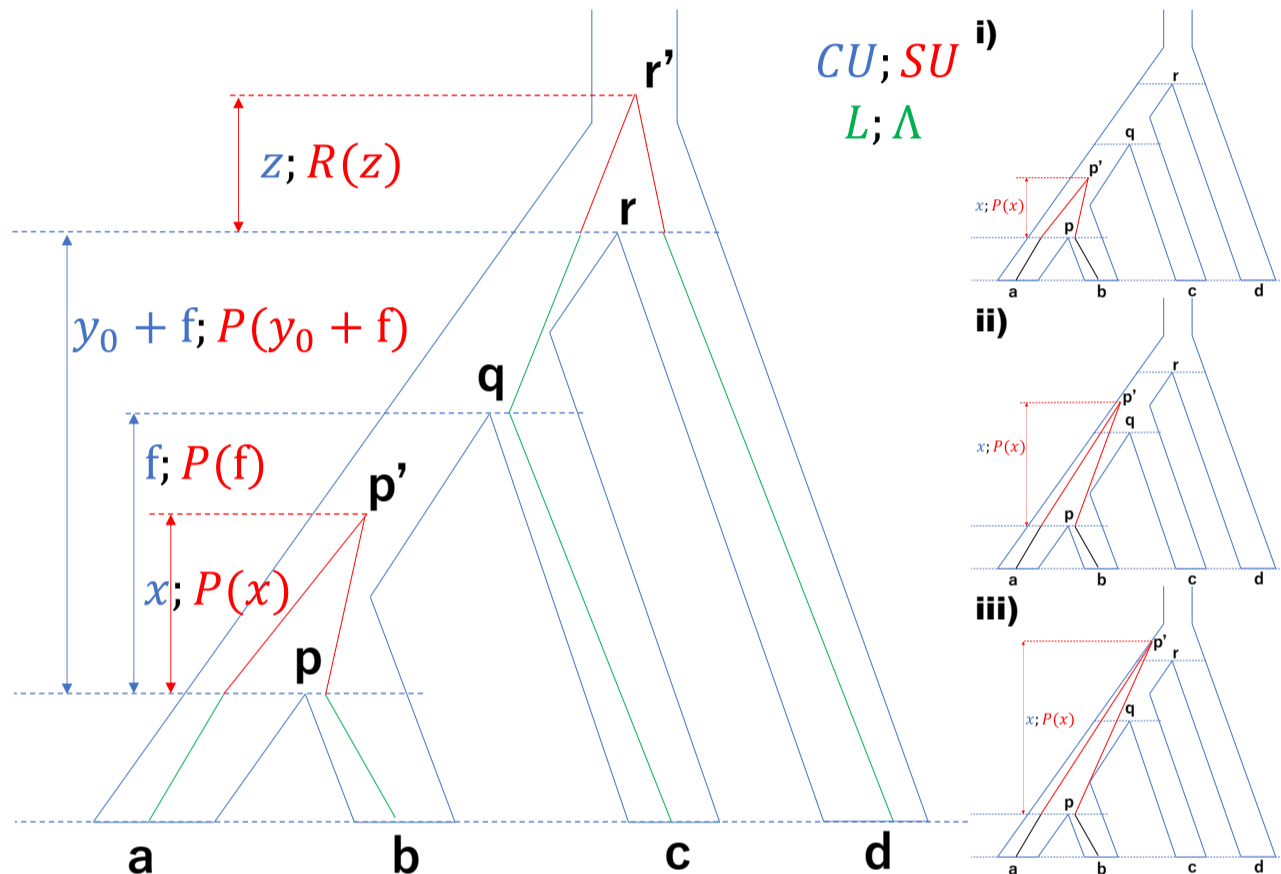


FIG. S25. Illustration of the unbalanced case. Lengths in CU/SU units are denoted in blue/red. Branches in green have a total length L/Λ in CU/SU units. The right-hand side shows the position of p' in relation to q and r in various cases.

- δ_1 : When G has topology $ab|cd$, p' must be the lowest point of coalescence. Thus,

$$\begin{aligned}
& \mathbb{E}[\delta_1 w_G(ab|cd)] \\
&= \mathbb{E}\left[\int_0^f \int_0^{+\infty} e^{-2P(x)-2R(z)-\Lambda} f_X(x) f_{Z|X}(z;x) dz dx\right] \\
&= \mathbb{E}\left[\int_0^f \int_0^{+\infty} e^{-2P(x)-2R(z)-\Lambda} e^{-x} e^{-z} dz dx\right] \\
&= \mathbb{E}\left[\int_0^f \int_0^{+\infty} e^{-2P(x)-2R(z)-\Lambda-x-z} dz dx\right];
\end{aligned}$$

and

$$\mathbb{E}[\delta_1 w_G^2(ab|cd)] \leq \mathbb{E}[\delta_1 w_G(ab|cd)] = O(f).$$

- δ_2 : When G has topology $ab|cd$, p' must be the lowest point of coalescence. Thus,

$$\begin{aligned}
& \mathbb{E}[\delta_2 w_G^2(ab|cd)] \\
&= \mathbb{E}\left[\int_f^{f+y_0} \int_0^{+\infty} e^{-4P(x)-4R(z)-2\Lambda} f_X(x) f_{Z|Y}(z;y) dz dx\right] \\
&= \mathbb{E}\left[\int_f^{f+y_0} \int_0^{+\infty} e^{-4P(x)-4R(z)-2\Lambda} e^{-3x+2f} e^{-z} dz dx\right] \\
&= \int_f^{f+y_0} \int_0^{+\infty} \mathbb{E}[e^{-4P(x)-4R(z)-2\Lambda}] e^{-3x-z+2f} dz dx.
\end{aligned}$$

- δ_3 : When G has the topology $ab|cd$, either p' or q' must be the lowest point of coalescence, and by symmetry, the two cases must have the same PDFs. Thus,

$$\begin{aligned}
& \mathbb{E}[\delta_3 w_G^2(ab|cd)] \\
&= \mathbb{E}\left[\int_{f+y_0}^{+\infty} \int_{x-f-y_0}^{+\infty} e^{-4P(x)-4R(z)-2\Lambda} 2f_X(x) f_{Z|X}(z;x) dz dx\right] \\
&= \mathbb{E}\left[\int_{f+y_0}^{+\infty} \int_{x-f-y_0}^{+\infty} e^{-4P(x)-4R(z)-2\Lambda} 2e^{-6x+5f+3y_0} e^{-z+x-f-y_0} dz dx\right] \\
&= \int_{f+y_0}^{+\infty} \int_{x-f-y_0}^{+\infty} \mathbb{E}[e^{-4P(x)-4R(z)-2\Lambda}] 2e^{-5x-z+4f+2y_0} dz dx.
\end{aligned}$$

Replacing in (11), by Jensen's inequality, we get

$$\begin{aligned}
\mathbb{E}[X_G] &= \mathbb{E}[\delta_1 w_G(ab|cd)] = \mathbb{E}\left[\int_0^f \int_0^{+\infty} e^{-2P(x)-2R(z)-\Lambda-x-z} dz dx\right] \\
&\geq \int_0^f \int_0^{+\infty} e^{-2P(x)-2R(z)-\Lambda-x-z} dz dx \\
&= \int_0^f \int_0^{+\infty} e^{-2\lambda x-2\lambda z-\lambda L-x-z} dz dx \\
&= \frac{e^{-\lambda L}(1-e^{-(1+2\lambda)f})}{(1+2\lambda)^2} = \frac{e^{-\lambda L}}{1+2\lambda} f + O(f^2).
\end{aligned}$$

And replacing in (12), we get

$$\begin{aligned}
\text{Var}[X_G] &= \mathbb{E}[(\delta_1 + 2\delta_2 + 2\delta_3)w_G^2(ab|cd)] - \mathbb{E}^2[X_G] = \mathbb{E}[2(\delta_2 + \delta_3)w_G^2(ab|cd)] + O(f) \\
&= \int_f^{f+y_0} \int_0^{+\infty} \mathbb{E}[e^{-4P(x)-4R(z)-2\Lambda}] 2e^{-3x-z+2f} dz dx \\
&\quad + \int_{f+y_0}^{+\infty} \int_{x-f-y_0}^{+\infty} \mathbb{E}[e^{-4P(x)-4R(z)-2\Lambda}] 4e^{-5x-z+2f+2y_0} dz dx + O(f),
\end{aligned}$$

from which our assumption of $\text{Var}[X_{G^*}] = \Theta_f(1)$ follows.

Let $F_P(u;x)$, $F_R(v;z)$, and $F_\Lambda(w)$ be the CDF of $P(x)$, $R(z)$, and Λ respectively; let $F_{PRA}(u,v,w;x,z)$ and $F_{PRA}(u,v,w;x,z)$ be the joint CDF and the joint PDF. Let $F_P^{-1}(t;x)$, $F_R^{-1}(t;z)$, and $F_\Lambda^{-1}(t)$ be the inverse function of CDF of $P(x)$, $R(z)$, and Λ respectively.

Then,

$$\begin{aligned}
& \mathbb{E} \left[e^{-2(2P(x)+2R(z)+\Lambda)} \right] \\
&= \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} e^{-2(2u+2v+w)} F_{PRA}(u, v, w; x, z) dw dv du \\
&= \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} e^{-2(2u+2v+w)} \frac{\partial^3 F_{PRA}}{\partial u \partial v \partial w} dw dv du \\
&= \int_0^{+\infty} \int_0^{+\infty} \left(e^{-2(2u+2v+w)} \frac{\partial^2 F_{PRA}}{\partial u \partial v} \Big|_{w=0}^{+\infty} \right. \\
&\quad \left. - \int_0^{+\infty} (-2) e^{-2(2u+2v+w)} \frac{\partial^2 F_{PRA}}{\partial u \partial v} dw \right) dv du \\
&= \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} 2e^{-2(2u+2v+w)} \frac{\partial^2 F_{PRA}}{\partial u \partial v} dv du dw \\
&= \int_0^{+\infty} \int_0^{+\infty} \left(2e^{-2(2u+2v+w)} \frac{\partial F_{PRA}}{\partial u} \Big|_{v=0}^{+\infty} - \int_0^{+\infty} (-8) e^{-2(2u+2v+w)} \frac{\partial F_{PRA}}{\partial u} dv \right) du dw \\
&= \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} 8e^{-2(2u+2v+w)} \frac{\partial F_{PRA}}{\partial u} du dw dv \\
&= \int_0^{+\infty} \int_0^{+\infty} \left(8e^{-2(2u+2v+w)} F_{PRA}(u, v, w; x, z) \Big|_{u=0}^{+\infty} \right. \\
&\quad \left. - \int_0^{+\infty} (-32) e^{-2(2u+2v+w)} F_{PRA}(u, v, w; x, z) du \right) dw dv \\
&= \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} 32e^{-2(2u+2v+w)} F_{PRA}(u, v, w; x, z) dw dv du \\
&\leq \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} 32e^{-2(2u+2v+w)} \min\{F_P(u; x), F_R(v; z), F_\Lambda(w)\} dw dv du \\
&= \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} 32e^{-2(2u+2v+w)} \left(\int_0^1 1_{t \leq F_P(u; x)} 1_{t \leq F_R(v; z)} 1_{t \leq F_\Lambda(w)} dt \right) dw dv du \\
&= \int_0^1 \int_0^{+\infty} \int_0^{+\infty} \int_0^{+\infty} 32e^{-2(2u+2v+w)} 1_{u \geq F_P^{-1}(t; x)} 1_{v \geq F_R^{-1}(t; z)} 1_{w \geq F_\Lambda^{-1}(t)} dw dv du dt \\
&= \int_0^1 \int_{F_P^{-1}(t; x)}^{+\infty} \int_{F_R^{-1}(t; z)}^{+\infty} \int_{F_\Lambda^{-1}(t)}^{+\infty} 32e^{-2(2u+2v+w)} dw dv du dt \\
&= \int_0^1 e^{-2(2F_P^{-1}(t; x) + 2F_R^{-1}(t; z) + F_\Lambda^{-1}(t))} dt.
\end{aligned}$$

Thus, for any $0 < t_0 < 1$,

$$\begin{aligned}
& \mathbb{E} \left[e^{-2(2P(x)+2R(z)+\Lambda)} \right] \\
&\leq \int_0^1 e^{-2(2F_P^{-1}(t; x) + 2F_R^{-1}(t; z) + F_\Lambda^{-1}(t))} dt \\
&\leq \int_0^{t_0} \overbrace{e^{-2(2F_P^{-1}(0; x) + 2F_R^{-1}(0; z) + F_\Lambda^{-1}(0))}}^1 dt + \int_{t_0}^1 e^{-2(2F_P^{-1}(t_0; x) + 2F_R^{-1}(t_0; z) + F_\Lambda^{-1}(t_0))} dt \\
&\leq t_0 + e^{-2(2F_P^{-1}(t_0; x) + 2F_R^{-1}(t_0; z) + F_\Lambda^{-1}(t_0))}.
\end{aligned}$$

By Chebyshev's inequality (using $t_0^{-\frac{1}{2}}$ as the constant), $F_P^{-1}(t_0; x) \geq (\lambda - \frac{\varepsilon}{\sqrt{t_0}})x$, $F_R^{-1}(t_0; z) \geq (\lambda - \frac{\varepsilon}{\sqrt{t_0}})y$, and $F_\Lambda^{-1}(t_0) \geq (\lambda - \frac{\varepsilon}{\sqrt{t_0}})L$. Thus,

$$\mathbb{E} \left[e^{-2(2P(x)+2R(z)+\Lambda)} \right] \leq t_0 + e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4x+4z+2L)}.$$

Thus,

$$\begin{aligned}
\text{Var}[X_{G^*}] &\leq \int_f^{f+y_0} \int_0^{+\infty} \left(t_0 + e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4x+4z+2L)} \right) 2e^{-3x-z+2f} dz dx \\
&\quad + \int_{f+y_0}^{+\infty} \int_{x-f-y_0}^{+\infty} \left(t_0 + e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4x+4z+2L)} \right) 4e^{-5x-z+2f+2y_0} dz dx + O(f) \\
&= \int_f^{f+y_0} \left(2t_0 e^{-3x+2f} + \frac{2}{1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}}} e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4x+2L)-3x+2f} \right) dx \\
&\quad + \int_{f+y_0}^{+\infty} \left(4t_0 e^{-6x+3f+3y_0} + \frac{4}{1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}}} e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(8x-4f-4y_0+2L)-6x+3f+3y_0} \right) dx + O(f) \\
&= \frac{2}{3} t_0 (e^{-f} - e^{-f-3y_0}) + \frac{2}{(1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})(3+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})} \left(e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4f+2L)-f} - e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4f+4y_0+2L)-f-3y_0} \right) \\
&\quad + \frac{4}{6} t_0 e^{-3f-3y_0} + \frac{4}{(1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})(6+8\lambda - \frac{8\varepsilon}{\sqrt{t_0}})} e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4f+4y_0+2L)-3f-3y_0} + O(f) \\
&= \frac{2}{3} t_0 + \frac{2e^{-2L(\lambda - \frac{\varepsilon}{\sqrt{t_0}})}}{(1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})(3+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})} + O(f).
\end{aligned}$$

Case 2: Balanced tree. Let $P(x)$ be functions to random variables denoting SU difference in heights of points (p, p') where p' is x CU distance above p ; let $Q(y)$ be functions to random variables denoting SU difference in heights of points (q, q') where q' is y CU distance above q . Note that $P(x_0+z) - P(x_0) = Q(y_0+z) - Q(y_0)$ where $P(x_0)$ and $Q(y_0)$ denote the SU length of (p, r) and (q, r) , respectively. Let random variable $\Lambda := (l_{S^+}(a, p) + l_{S^+}(b, p) + l_{S^+}(c, q) + l_{S^+}(d, q))$ be the total SU terminal branch lengths and the constant value L be the CU distance corresponding to Λ .

- δ_1 : Here,

$$\mathbb{E}[\delta_1 w_G(ab|cd)] = \mathbb{E} \left[\int_0^{x_0} \int_0^{+\infty} e^{-2P(x)-2Q(y)-\Lambda} e^{-x} e^{-y} dy dx \right];$$

and

$$\mathbb{E}[\delta_1 w_G^2(ab|cd)] \leq \mathbb{E}[\delta_1 w_G(ab|cd)] = O(f).$$

- δ_2 : Here,

$$\mathbb{E}[\delta_2 w_G(ab|cd)] = \mathbb{E} \left[\int_{x_0}^{+\infty} \int_0^{y_0} e^{-2P(x)-2Q(y)-\Lambda} e^{-x} e^{-y} dy dx \right];$$

and

$$\mathbb{E}[\delta_2 w_G^2(ab|cd)] \leq \mathbb{E}[\delta_2 w_G(ab|cd)] = O(f).$$

- δ_3 : Similar to the unbalanced case, when G has the topology $ab|cd$, either p' or q' must be the lowest point of coalescence, and by symmetry, the two cases must have the same PDFs. Thus,

$$\begin{aligned}
\mathbb{E}[\delta_3 w_G^2(ab|cd)] &= \mathbb{E} \left[\int_{x_0}^{+\infty} \int_{x-x_0+y_0}^{+\infty} e^{-4P(x)-4Q(y)-2\Lambda} 2e^{-x_0} e^{-y_0} e^{-6x+6x_0} e^{-y+x-x_0+y_0} dy dx \right] \\
&= \int_{x_0}^{+\infty} \int_{x-x_0+y_0}^{+\infty} \mathbb{E} \left[e^{-4P(x)-4Q(y)-2\Lambda} \right] 2e^{-5x-y+4x_0} dy dx.
\end{aligned}$$

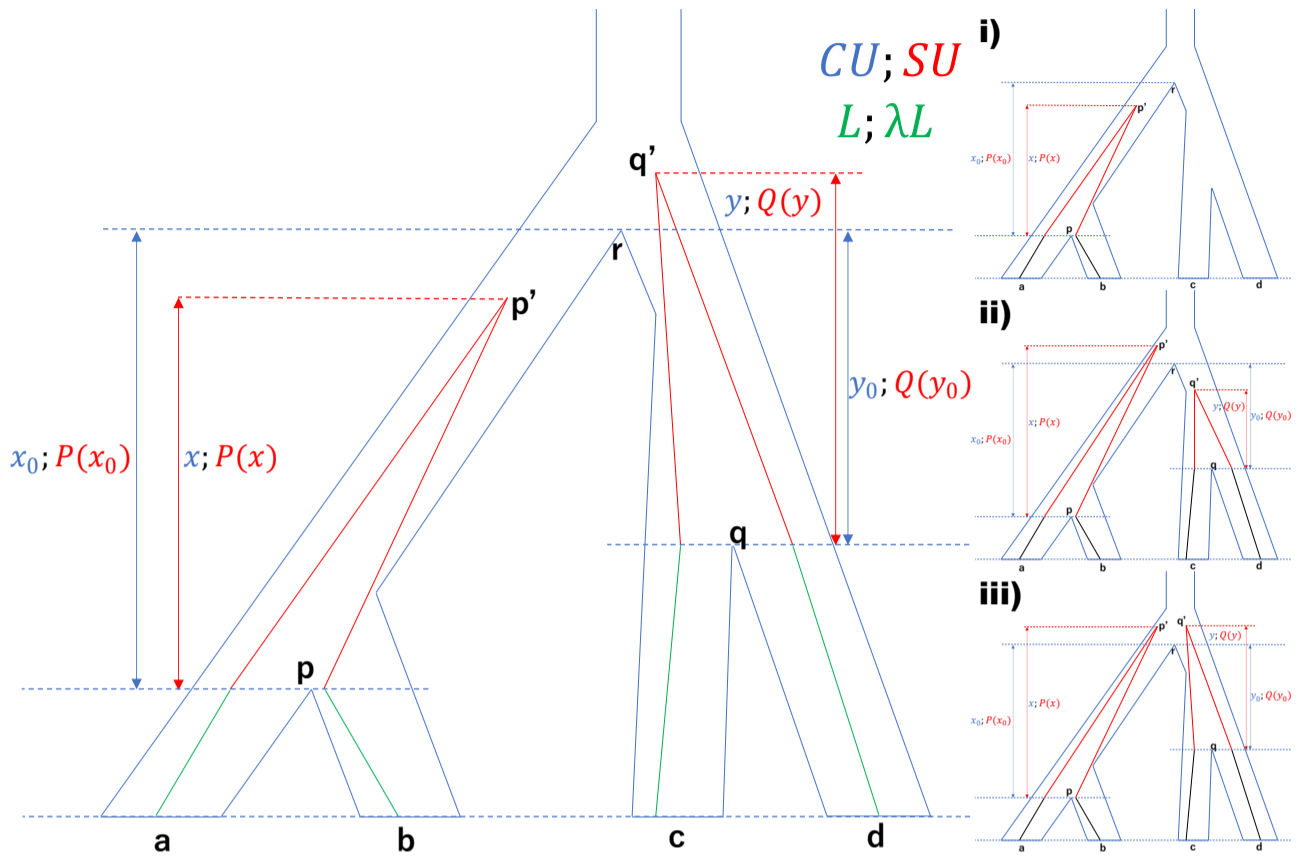


FIG. S26. Illustration of the unbalanced case. Lengths in CU/SU units are denoted in blue/red. Branches in green have a total length L/λ in CU/SU units. The right-hand side shows the position of p' and q' in relation to r in various cases.

Replacing in (13), we get

$$\begin{aligned}
\mathbb{E}[X_G] &= \mathbb{E}[(\delta_1 + \delta_2)w_G(ab|cd)] \\
&= \mathbb{E}\left[\int_0^{x_0} \int_0^{+\infty} e^{-2P(x)-2Q(y)-\Lambda} e^{-x} e^{-y} dy dx + \int_{x_0}^{+\infty} \int_0^{y_0} e^{-2P(x)-2Q(y)-\Lambda} e^{-x} e^{-y} dy dx\right] \\
&\geq \int_0^{x_0} \int_0^{+\infty} e^{-2\lambda x-2\lambda y-\lambda L} e^{-x} e^{-y} dy dx + \int_{x_0}^{+\infty} \int_0^{y_0} e^{-2\lambda x-2\lambda y-\lambda L} e^{-x} e^{-y} dy dx \\
&= \frac{(x_0 + y_0)e^{-\lambda L}}{1+2\lambda} + O(f^2) = \frac{fe^{-\lambda L}}{1+2\lambda} + O(f^2);
\end{aligned}$$

and replacing in (14), for any $0 < t_0 < 1$,

$$\begin{aligned}
\text{Var}[X_G] &= \mathbb{E}[(\delta_1 + \delta_2 + 2\delta_3)w_G^2(ab|cd)] - \mathbb{E}^2[X_G] \\
&= \mathbb{E}[2\delta_3 w_G^2(ab|cd)] + O(f) \\
&= \int_{x_0}^{+\infty} \int_{x-x_0+y_0}^{+\infty} \mathbb{E}\left[e^{-4P(x)-4Q(y)-2\Lambda}\right] 4e^{-5x-y+4x_0} dy dx \\
&\leq \int_{x_0}^{+\infty} \int_{x-x_0+y_0}^{+\infty} \left(t_0 + e^{(-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4x+4y+2L)}\right) 4e^{-5x-y+4x_0} dy dx + O(f) \\
&= \int_{x_0}^{+\infty} \left(4e^{-6x-y_0+5x_0} t_0 + \frac{4}{1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}}} e^{-6x-y_0+5x_0 + (-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(8x-4x_0+4y_0+2L)}\right) dx + O(f) \\
&= \frac{4}{6} e^{-x_0-y_0} t_0 + \frac{4}{(1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})(6+8\lambda - \frac{8\varepsilon}{\sqrt{t_0}})} e^{-x_0-y_0 + (-\lambda + \frac{\varepsilon}{\sqrt{t_0}})(4x_0+4y_0+2L)} + O(f) \\
&= \frac{2}{3} t_0 + \frac{2e^{-2L(\lambda - \frac{\varepsilon}{\sqrt{t_0}})}}{(1+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})(3+4\lambda - \frac{4\varepsilon}{\sqrt{t_0}})} + O(f),
\end{aligned}$$

from which our assumption of $\text{Var}[X_G] = \Theta_f(1)$ follows.

Thus, for both balanced and unbalanced trees, the variance is bounded the by same expression, and thus in both cases,

$$\begin{aligned}
\text{Var}[X_{G^*}] &\leq \frac{2}{3} t_0 + 2 \frac{e^{-2\lambda L}}{(1+4\lambda)(3+4\lambda)} \frac{1}{\left(1 - \frac{4\varepsilon}{(1+4\lambda)\sqrt{t_0}}\right) \left(1 - \frac{4\varepsilon}{(3+4\lambda)\sqrt{t_0}}\right) e^{-\frac{2\varepsilon L}{\sqrt{t_0}}}} + O(f) \\
&\leq \frac{2}{3} t_0 + 2 \frac{e^{-2\lambda L}}{\left(1 - \frac{4\varepsilon}{(1+4\lambda)\sqrt{t_0}}\right) \left(1 - \frac{4\varepsilon}{(3+4\lambda)\sqrt{t_0}}\right) \left(1 - \frac{2\varepsilon L}{\sqrt{t_0}}\right)} + O(f) \\
&\leq \frac{2}{3} t_0 + 2 \frac{e^{-2\lambda L}}{\left(1 - \frac{4\varepsilon}{(1+4\lambda)\sqrt{t_0}} - \frac{4\varepsilon}{(3+4\lambda)\sqrt{t_0}} - \frac{2\varepsilon L}{\sqrt{t_0}}\right)} + O(f) \\
&= \frac{2}{3} t_0 + \frac{2e^{-2\lambda L}}{(3+16\lambda+16\lambda^2) - \frac{\varepsilon}{\sqrt{t_0}}((16+32\lambda) + (6+32\lambda+32\lambda^2)L)} + O(f).
\end{aligned}$$

Now, let $C := (16+32\lambda) + (6+32\lambda+32\lambda^2)L$, $t_0 = \left(\frac{C^{\frac{1}{3}} \varepsilon^{\frac{1}{3}}}{(3+16\lambda+16\lambda^2)e^{\frac{2}{3}\lambda L}}\right)^2$, we get

$$\begin{aligned}
\text{Var}[X_{G^*}] &\leq \frac{2e^{-2\lambda L}}{3(3+16\lambda+16\lambda^2)^2} \left((\varepsilon e^{\lambda L} C)^{\frac{2}{3}} + \frac{9+48\lambda+48\lambda^2}{1 - (\varepsilon e^{\lambda L} C)^{\frac{2}{3}}} \right) + O(f) \\
&= \frac{2e^{-2\lambda L}}{3(3+16\lambda+16\lambda^2)} \left(\frac{(\varepsilon e^{\lambda L} C)^{\frac{2}{3}}}{3+16\lambda+16\lambda^2} + 3 + \frac{3(\varepsilon e^{\lambda L} C)^{\frac{2}{3}}}{1 - (\varepsilon e^{\lambda L} C)^{\frac{2}{3}}} \right) + O(f).
\end{aligned}$$

Now, recalling that $\varepsilon = \frac{e^{-\lambda L}}{C} \left(\frac{20(\lambda+\lambda^2)}{9(1+2\lambda)^2}\right)^{\frac{3}{2}}$,

$$\begin{aligned}
\text{Var}[X_{G^*}] &\leq \frac{2}{3(3+16\lambda+16\lambda^2)(1+2\lambda)^2} \\
&\quad \left(\frac{\frac{20}{9}(\lambda+\lambda^2)}{3+16\lambda+16\lambda^2} + 3 + \frac{3(\frac{20}{9})(\lambda+\lambda^2)}{1-\frac{\frac{20}{9}(\lambda+\lambda^2)}{1+4\lambda+4\lambda^2}} \right) + O(f) \\
&\leq \frac{2}{3(3+16\lambda+16\lambda^2)(1+2\lambda)^2} \left(\frac{20}{27}\lambda + 3 + \frac{\frac{20}{3}(\lambda+\lambda^2)}{1-\frac{5}{9}} \right) + O(f) \\
&= \frac{2}{3(3+16\lambda+16\lambda^2)(1+2\lambda)^2} \left(\frac{20}{27}\lambda + 3 + 15(\lambda+\lambda^2) \right) + O(f) \\
&< \frac{2}{3(1+2\lambda)^2} \left(\frac{3+16\lambda+15\lambda^2}{3+16\lambda+16\lambda^2} \right) + O(f).
\end{aligned}$$

□

THEOREM 2. *Under the conditions of Proposition 2 or Proposition 3,*

$$P\left(\sum_{G \in \mathcal{G}} w_G(ab|cd) \leq \sum_{G \in \mathcal{G}} w_G(ac|bd)\right) \leq P\left(\sum_{G \in \mathcal{G}} \delta_G(ab|cd) \leq \sum_{G \in \mathcal{G}} \delta_G(ac|bd)\right).$$

Proof. We start with proving this theorem under the conditions of Proposition 2. Recall $X_G := w_G(ab|cd) - w_G(ac|bd)$ and $Y_G := \delta_G(ab|cd) - \delta_G(ac|bd)$, and let $\bar{X}_{\mathcal{G}} = \frac{1}{k} \sum_{G \in \mathcal{G}} X_G$ and $\bar{Y}_{\mathcal{G}} = \frac{1}{k} \sum_{G \in \mathcal{G}} Y_G$. Recall also that under Proposition 2, proved below, under conditions of Theorem 2, we have $\text{Var}[\bar{X}_{\mathcal{G}}] = \Omega(1)$ and

$$\frac{\mathbb{E}[X_G]}{\sqrt{\text{Var}[X_G]}} = -\sqrt{\frac{3+16\lambda+16\lambda^2}{3+16\lambda+15\lambda^2}} \sqrt{\frac{3}{2}} f + O(f^2). \quad (15)$$

Similarly, we can compute the ratio of mean and variance for Y (corresponding to unweighted ASTRAL):

$$\mathbb{E}[Y_G] := \mathbb{E}[\delta_G(ab|cd) - \delta_G(ac|bd)] = 1 - e^{-f} = f + O(f^2)$$

$$\text{Var}[Y_G] := \text{Var}[\delta_G(ab|cd) - \delta_G(ac|bd)] = \frac{5}{3} e^{-f} - e^{-2f} = \frac{2}{3} + O(f)$$

and thus,

$$\frac{\mathbb{E}[Y_G]}{\sqrt{\text{Var}[Y_G]}} = \sqrt{\frac{3}{2}} f + O(f^2). \quad (16)$$

Given Proposition 2, we can use Berry–Esseen theorem to derive

$$\begin{aligned}
P(\bar{X}_{\mathcal{G}} \leq 0) &= P\left(\frac{\sqrt{k}}{\sqrt{\text{Var}[X_G]}} (\bar{X}_{\mathcal{G}} - \mathbb{E}[X_G]) \leq -\frac{\sqrt{k}}{\sqrt{\text{Var}[X_G]}} \mathbb{E}[X_G]\right) = \\
&\quad \Phi\left(-\sqrt{k} \frac{\mathbb{E}[X_G]}{\sqrt{\text{Var}[X_G]}}\right) + O\left(\frac{1}{\sqrt{k}}\right),
\end{aligned}$$

where Φ denotes CDF of the standard Normal distribution. Since $k = \Theta(f^{-2})$,

$$P(\bar{X}_{\mathcal{G}} \leq 0) = \Phi\left(-\sqrt{k} \frac{\mathbb{E}[X_G]}{\sqrt{\text{Var}[X_G]}}\right) + O(f) \quad (17)$$

and

$$P(\bar{Y}_{\mathcal{G}} \leq 0) = \Phi\left(-\sqrt{k} \frac{\mathbb{E}[Y_G]}{\sqrt{\text{Var}[Y_G]}}\right) + O(f), \quad (18)$$

Combining equations (17) and (18) with (15) and (16), we get

$$P\left(\sum_{G \in \mathcal{G}} w_G(ab|cd) \leq \sum_{G \in \mathcal{G}} w_G(ac|bd)\right) = \Phi\left(-\sqrt{\frac{3+16\lambda+16\lambda^2}{3+16\lambda+15\lambda^2}} \sqrt{\frac{3}{2}} f\sqrt{k}\right) + O(f)$$

and

$$P\left(\sum_{G \in \mathcal{G}} \delta_G(ab|cd) \leq \sum_{G \in \mathcal{G}} \delta_G(ac|bd)\right) = \Phi\left(-\sqrt{\frac{3}{2}} f\sqrt{k}\right) + O(f).$$

As $f \rightarrow 0$, the interval $(-\sqrt{1 + \frac{4\lambda+4\lambda^2}{3(1+2\lambda)^2}} \sqrt{\frac{3}{2}} f\sqrt{k}, -\sqrt{\frac{3}{2}} f\sqrt{k})$ does not shrink because $\Theta(f\sqrt{k}) = \Theta(1)$. Thus, we have

$$\Phi\left(-\sqrt{\frac{3}{2}} f\sqrt{k}\right) - \Phi\left(-\sqrt{1 + \frac{4\lambda+4\lambda^2}{3(1+2\lambda)^2}} \sqrt{\frac{3}{2}} f\sqrt{k}\right) = \Theta(1)$$

ensuring that

$$P\left(\sum_{G \in \mathcal{G}} w_G(ab|cd) \leq \sum_{G \in \mathcal{G}} w_G(ac|bd)\right) \leq P\left(\sum_{G \in \mathcal{G}} \delta_G(ab|cd) \leq \sum_{G \in \mathcal{G}} \delta_G(ac|bd)\right).$$

The proof under Proposition 3 is similar. Recall that under Proposition 3, $\text{Var}[X_{G^*}] = \Theta_f(1)$ and

$$\frac{\mathbb{E}[X_{G^*}]}{\sqrt{\text{Var}[X_{G^*}]}} \geq \sqrt{\frac{3}{2}} \left(1 - \frac{4\lambda^2}{(1+4\lambda)^2}\right)^{-\frac{1}{2}} f + O(f^2). \quad (19)$$

Given this result, the rest of the proof is similar to the proof under the conditions of Proposition 2, culminating in

$$P\left(\sum_{G^* \in \mathcal{G}} w_{G^*}(ab|cd) \leq \sum_{G^* \in \mathcal{G}} w_{G^*}(ac|bd)\right) \leq \Phi\left(-\left(1 - \frac{4\lambda^2}{(1+4\lambda)^2}\right)^{-\frac{1}{2}} \sqrt{\frac{3}{2}} f\sqrt{k}\right) + O(f).$$

□

Placement-based Algorithm

In this section, for a node v in tree G , we let \mathcal{L}_v denote the set of leaves under v .

Proof of Theorem 3

THEOREM 3. *Let S be a species tree, i be a species not in \mathcal{L}_S , \mathcal{S} be the set of possible species tree topologies by placing i onto S , and S' be the output of Algorithm S1. Then, $W(S', \mathcal{G}) = \max_{\hat{S} \in \mathcal{S}} W(\hat{S}, \mathcal{G})$.*

Proof. We start with two propositions, proved below.

PROPOSITION 5. *After each call to $\text{ColorLeafSet}(\mathcal{L}^*, X, T, \mathcal{G}, W)$ with a $T \neq \emptyset$, $W[T] = \sum_{G \in \mathcal{G}} W(T, G)$.*

PROPOSITION 6. *Before calling OptimalTreeDP in line 6 of Algorithm S1, lookup table W contains all tripartitions corresponding to internal nodes of all tree topologies in \mathcal{S} .*

By Proposition 6, all tripartitions corresponding to internal nodes of all tree topologies in \mathcal{S} pre-computed. Then, OptimalTreeDP uses a dynamic programming algorithm similar to the one formulated by Mirarab and Warnow 2015 to compute $\arg\max_{\hat{S} \in \mathcal{S}} W(\hat{S}, \mathcal{G})$. \square

PROPOSITION 5. *After each call to $\text{ColorLeafSet}(\mathcal{L}^*, X, T, \mathcal{G}, W)$ with a $T \neq \emptyset$, $W[T] = \sum_{G \in \mathcal{G}} W(T, G)$.*

Proof. For a gene tree node w and a color X , let \mathcal{L}_w^X denote the set of leaves in \mathcal{L}_w colored by X . For an internal node w , let u, v be the children of w , p be the parent of w (if w is not the root), and e denote the branch (w, p) . For a leaf i and internal node w , let $\mathcal{P}_{i,w}$ denote path between i and w and $s(\mathcal{P}) = 1 - \prod_{\hat{e} \in \mathcal{P}} (1 - s(\hat{e}))$. For leaves i, j , let $m(i, j)$ denote MRCA of i and j . Referring back to Table S1, we first establish the connection between recursive formulas of the algorithm and counter definitions.

- When $u_X = \sum_{i \in \mathcal{L}_u^X} e^{-l(\mathcal{P}_{i,w})}$, $v_X = \sum_{i \in \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,w})}$,

$$w_X := \left((u_X + v_X) e^{-l(e)} \right) = \sum_{i \in \mathcal{L}_w^X} e^{-l(\mathcal{P}_{i,w})} e^{-l(e)} = \sum_{i \in \mathcal{L}_w^X} e^{-l(\mathcal{P}_{i,p})}.$$

- When $u_{XX}^+ = \sum_{\{i,j\} \subseteq \mathcal{L}_u^X} e^{-l(\mathcal{P}_{i,j})}$, $v_{XX}^+ = \sum_{\{i,j\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,j})}$,

$$\begin{aligned} w_{XX}^+ &:= u_{XX}^+ + v_{XX}^+ + u_X v_X = \sum_{\{i,j\} \subseteq \mathcal{L}_u^X} e^{-l(\mathcal{P}_{i,j})} + \sum_{\{i,j\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,j})} + \sum_{i \in \mathcal{L}_u^X} e^{-l(\mathcal{P}_{i,w})} \sum_{j \in \mathcal{L}_v^X} e^{-l(\mathcal{P}_{j,w})} \\ &= \sum_{\{i,j\} \subseteq \mathcal{L}_u^X} e^{-l(\mathcal{P}_{i,j})} + \sum_{\{i,j\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,j})} + \sum_{i \in \mathcal{L}_u^X} \sum_{j \in \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,j})} = \sum_{\{i,j\} \subseteq \mathcal{L}_w^X} e^{-l(\mathcal{P}_{i,j})}. \end{aligned}$$

- For $X \neq Y$, when $u_{XY}^+ = \sum_{(i,j) \in \mathcal{L}_u^X \times \mathcal{L}_u^Y} e^{-l(\mathcal{P}_{i,j})}$, $v_{XY}^+ = \sum_{(i,j) \in \mathcal{L}_v^X \times \mathcal{L}_v^Y} e^{-l(\mathcal{P}_{i,j})}$,

$$\begin{aligned} w_{XY}^+ &:= u_{XY}^+ + v_{XY}^+ + u_X v_Y + u_Y v_X \\ &= \sum_{(i,j) \in \mathcal{L}_u^X \times \mathcal{L}_u^Y} e^{-l(\mathcal{P}_{i,j})} + \sum_{(i,j) \in \mathcal{L}_v^X \times \mathcal{L}_v^Y} e^{-l(\mathcal{P}_{i,j})} + \sum_{(i,j) \in \mathcal{L}_u^X \times \mathcal{L}_v^Y} e^{-l(\mathcal{P}_{i,j})} + \sum_{(i,j) \in \mathcal{L}_v^X \times \mathcal{L}_u^Y} e^{-l(\mathcal{P}_{i,j})} \\ &= \sum_{\{i,j\} \subseteq \mathcal{L}_w^X \times \mathcal{L}_w^Y} e^{-l(\mathcal{P}_{i,j})}. \end{aligned}$$

- When $u_{\bar{X}X} = \sum_{\{i,j\} \subseteq \mathcal{L}_u^X} e^{-l(\mathcal{P}_{i,j})} \prod_{\hat{e} \in \mathcal{P}_{m(i,j),w}} (1-s(\hat{e}))$, $v_{\bar{X}X} = \sum_{\{i,j\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,j})} \prod_{\hat{e} \in \mathcal{P}_{m(i,j),w}} (1-s(\hat{e}))$,

$$\begin{aligned} w_{\bar{X}X} &:= (u_{\bar{X}X} + v_{\bar{X}X} + u_X v_X) (1-s(e)) \\ &= \sum_{\{i,j\} \subseteq \mathcal{L}_u^X} e^{-l(\mathcal{P}_{i,j})} \prod_{\hat{e} \in \mathcal{P}_{m(i,j),p}} (1-s(\hat{e})) + \sum_{\{i,j\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,j})} \prod_{\hat{e} \in \mathcal{P}_{m(i,j),p}} (1-s(\hat{e})) \\ &\quad + \sum_{(i,j) \in \mathcal{L}_u^X \times \mathcal{L}_v^X} e^{-l(\mathcal{P}_{i,j})} (1-s(e)) = \sum_{\{i,j\} \subseteq \mathcal{L}_w^X} e^{-l(\mathcal{P}_{i,j})} \prod_{\hat{e} \in \mathcal{P}_{m(i,j),p}} (1-s(\hat{e})). \end{aligned}$$

- When $u_{\bar{X}Y} = \sum_{(i,j) \in \mathcal{L}_u^X \times \mathcal{L}_u^Y} e^{-l(\mathcal{P}_{i,j})} (1-s(\mathcal{P}_{m(i,j),w}))$, $v_{\bar{X}Y} = \sum_{(i,j) \in \mathcal{L}_u^X \times \mathcal{L}_v^Y} e^{-l(\mathcal{P}_{i,j})} (1-s(\mathcal{P}_{m(i,j),w}))$, and $X \neq Y$, similarly,

$$w_{\bar{X}Y} := (u_{\bar{X}Y} + v_{\bar{X}Y} + u_X v_Y + u_Y v_X) (1-s(e)) = \sum_{(i,j) \in \mathcal{L}_w^X \times \mathcal{L}_w^Y} e^{-l(\mathcal{P}_{i,j})} (1-s(\mathcal{P}_{m(i,j),p})).$$

- For $X \neq Y$, when $u_{XX|Y} = \sum_{\{i,j\} \subseteq \mathcal{L}_u^X} \sum_{k \in \{k' \in \mathcal{L}_v^Y : \mathcal{L}_{m(i,j)} \subsetneq \mathcal{L}_{m(i,k')}\}} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),m(i,k)})$, $v_{XX|Y} = \sum_{\{i,j\} \subseteq \mathcal{L}_v^X} \sum_{k \in \{k' \in \mathcal{L}_v^Y : \mathcal{L}_{m(i,j)} \subsetneq \mathcal{L}_{m(i,k')}\}} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),m(i,k)})$,

$$w_{XX|Y} := (u_{XX|Y} + v_{XX|Y} + (u_{XX}^+ - u_{XX}^-)v_Y + u_Y(v_{XX}^+ - v_{XX}^-)) e^{-l(e)}.$$

Notice that $(u_{XX}^+ - u_{XX}^-)v_Y = \sum_{\{i,j\} \subseteq \mathcal{L}_u^X} \sum_{k \in \mathcal{L}_v^Y} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),w})$ and $u_Y(v_{XX}^+ - v_{XX}^-) = \sum_{\{i,j\} \subseteq \mathcal{L}_v^X} \sum_{k \in \mathcal{L}_u^Y} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),w})$. Thus,

$$\begin{aligned} w_{XX|Y} &= \sum_{\{i,j\} \subseteq \mathcal{L}_w^X} \sum_{k \in \{k' \in \mathcal{L}_w^Y : \mathcal{L}_{m(i,j)} \subsetneq \mathcal{L}_{m(i,k')}\}} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),m(i,k)}) e^{-l(e)} \\ &= \sum_{\{i,j\} \subseteq \mathcal{L}_w^X} \sum_{k \in \{k' \in \mathcal{L}_w^Y : \mathcal{L}_{m(i,j)} \subsetneq \mathcal{L}_{m(i,k')}\}} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,p})} s(\mathcal{P}_{m(i,j),m(i,k)}). \end{aligned}$$

- Similarly, when $u_{XY|Z} = \sum_{(i,j) \in \mathcal{L}_u^X \times \mathcal{L}_u^Y} \sum_{k \in \{k' \in \mathcal{L}_u^Z : \mathcal{L}_{m(i,j)} \subsetneq \mathcal{L}_{m(i,k')}\}} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),m(i,k)})$, $v_{XY|Z} = \sum_{(i,j) \in \mathcal{L}_v^X \times \mathcal{L}_v^Y} \sum_{k \in \{k' \in \mathcal{L}_v^Z : \mathcal{L}_{m(i,j)} \subsetneq \mathcal{L}_{m(i,k')}\}} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),m(i,k)})$, for distinct X, Y, Z ,

$$w_{XY|Z} = \sum_{(i,j) \in \mathcal{L}_w^X \times \mathcal{L}_w^Y} \sum_{k \in \{k' \in \mathcal{L}_w^Z : \mathcal{L}_{m(i,j)} \subsetneq \mathcal{L}_{m(i,k')}\}} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,p})} s(\mathcal{P}_{m(i,j),m(i,k)}).$$

- For distinct X, Y, Z ,

$$\begin{aligned} w_{XX|YZ} &:= v_X u_{YZ|X} + u_X v_{YZ|X} + u_{XX|Z} v_Y + v_{XX|Z} u_Y + u_{XX|Y} v_Z + v_{XX|Y} u_Z \\ &\quad + (u_{YZ}^+ v_{XX}^+ - u_{YZ}^- v_{XX}^-) + (u_{XX}^+ v_{YZ}^+ - u_{XX}^- v_{YZ}^-). \end{aligned}$$

Notice that,

$$\begin{aligned} v_X u_{YZ|X} &= \sum_{(h,i,j,k) \in \mathcal{L}_v^X \times \mathcal{L}_u^Y \times \mathcal{L}_u^Z \times \mathcal{L}_u^X} \delta_G(hk|ij) e^{-l(\mathcal{P}_{h,w})} e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,w})} s(\mathcal{P}_{m(i,j),m(i,k)}) \\ &= \sum_{(h,i,j,k) \in \mathcal{L}_v^X \times \mathcal{L}_u^Y \times \mathcal{L}_u^Z \times \mathcal{L}_u^X} \delta_G(hk|ij) e^{-l(\mathcal{P}_{i,j})-l(\mathcal{P}_{k,h})} s(\mathcal{P}_{m(i,j),m(i,k)}) \\ &= \sum_{(h,i,j,k) \in \mathcal{L}_v^X \times \mathcal{L}_u^Y \times \mathcal{L}_u^Z \times \mathcal{L}_u^X} w_G(hk|ij). \end{aligned}$$

Similarly,

$$\begin{aligned}
u_X v_{YZ|X} &= \sum_{\substack{h \in \mathcal{L}_u^X \\ i \in \mathcal{L}_v^Y \\ j \in \mathcal{L}_v^Z \\ k \in \mathcal{L}_v^X}} w_G(hk|ij), u_{XX|Z} v_Y = \sum_{\substack{\{h,i\} \subseteq \mathcal{L}_u^X \\ j \in \mathcal{L}_u^Z \\ k \in \mathcal{L}_v^Y}} w_G(hi|jk), v_{XX|Z} u_Y = \sum_{\substack{\{h,i\} \subseteq \mathcal{L}_v^X \\ j \in \mathcal{L}_v^Z \\ k \in \mathcal{L}_u^Y}} w_G(hi|jk), \\
u_{XX|Y} v_Z &= \sum_{\substack{\{h,i\} \subseteq \mathcal{L}_u^X \\ j \in \mathcal{L}_u^Y \\ k \in \mathcal{L}_v^Z}} w_G(hi|jk), v_{XX|Y} u_Z = \sum_{\substack{\{h,i\} \subseteq \mathcal{L}_v^X \\ j \in \mathcal{L}_v^Y \\ k \in \mathcal{L}_u^Z}} w_G(hi|jk).
\end{aligned}$$

Also,

$$\begin{aligned}
u_{YZ}^+ v_{XX}^+ - u_{YZ}^- v_{XX}^- &= \sum_{(h,i) \in \mathcal{L}_u^Y \times \mathcal{L}_u^Z} \sum_{\{j,k\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{h,i}) - l(\mathcal{P}_{j,k})} \\
&\quad - \sum_{(h,i) \in \mathcal{L}_u^Y \times \mathcal{L}_u^Z} \sum_{\{j,k\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{h,i}) - l(\mathcal{P}_{j,k})} \prod_{\hat{e} \in \mathcal{P}_{m(h,i),w}} (1 - s(\hat{e})) \prod_{\hat{e} \in \mathcal{P}_{m(j,k),w}} (1 - s(\hat{e})) \\
&= \sum_{(h,i) \in \mathcal{L}_u^Y \times \mathcal{L}_u^Z} \sum_{\{j,k\} \subseteq \mathcal{L}_v^X} e^{-l(\mathcal{P}_{h,i}) - l(\mathcal{P}_{j,k})} \left(1 - \prod_{\hat{e} \in \mathcal{P}_{m(h,i),m(j,k)}} (1 - s(\hat{e})) \right) \\
&= \sum_{(h,i) \in \mathcal{L}_u^Y \times \mathcal{L}_u^Z} \sum_{\{j,k\} \subseteq \mathcal{L}_v^X} w_G(hi|jk).
\end{aligned}$$

Similarly,

$$u_{XX}^+ v_{YZ}^+ - u_{XX}^- v_{YZ}^- = \sum_{\{h,i\} \subseteq \mathcal{L}_u^X} \sum_{(j,k) \in \mathcal{L}_v^Y \times \mathcal{L}_v^Z} w_G(hi|jk).$$

Notice that above cases count exactly once all quartets $hi|jk$ for all leaf nodes h, i colored X, j colored Y, k colored Z such that MRCA of h, i, j, k is w ; namely,

$$w_{XX|YZ} = \sum_{\{h,i\} \subseteq \mathcal{L}_w^X} \sum_{j \in \mathcal{L}_w^Y} \sum_{k \in \{k' : k' \in \mathcal{L}_w^Z, \text{MRCA}(h, i, j, k') = w\}} w_G(hi|jk).$$

- We define $I(G)$ to be the set of internal nodes of gene tree G and \mathcal{L}_G^X be the set of leaves of gene tree G with color X . It is trivial to verify that at the

$$Q = \sum_{G \in \mathcal{G}} \sum_{w \in I(G)} w_{AA|BC} + \sum_{G \in \mathcal{G}} \sum_{w \in I(G)} w_{BB|CA} + \sum_{G \in \mathcal{G}} \sum_{w \in I(G)} w_{CC|AB}.$$

At the end of procedure **UpdateCounters**, $\sum_{w \in I(G)} w_{XX|YZ} = \sum_{\{h,i\} \subseteq \mathcal{L}_G^X} \sum_{(j,k) \in \mathcal{L}_G^Y \times \mathcal{L}_G^Z} w_G(hi|jk)$.

Thus, Q returned by **UpdateCounters** satisfies:

$$Q = \sum_{G \in \mathcal{G}} \left(\sum_{\substack{\{h,i\} \subseteq \mathcal{L}_G^A \\ (j,k) \in \mathcal{L}_G^B \times \mathcal{L}_G^C}} w_G(hi|jk) + \sum_{\substack{\{h,i\} \subseteq \mathcal{L}_G^B \\ (j,k) \in \mathcal{L}_G^C \times \mathcal{L}_G^A}} w_G(hi|jk) + \sum_{\substack{\{h,i\} \subseteq \mathcal{L}_G^C \\ (j,k) \in \mathcal{L}_G^A \times \mathcal{L}_G^B}} w_G(hi|jk) \right).$$

For tripartition $T = A|B|C$, note that by assumption, before the call, all the gene tree leaves are colored such that recoloring \mathcal{L}^* by X would produce a coloring that matches T . Thus, at the end of the call to **ColorLeafSet**, for each gene tree G , we have $A \cap \mathcal{L}_G = \mathcal{L}_G^A$, $B \cap \mathcal{L}_G = \mathcal{L}_G^B$, and $C \cap \mathcal{L}_G = \mathcal{L}_G^C$.

Then, the value returned by **UpdateCounters** satisfies:

$$Q = \sum_{G \in \mathcal{G}} W(A|B|C, G). \quad (20)$$

It can be easily verified that after each call to `ColorLeafSet`($\mathcal{L}^*, X, T, \mathcal{G}, W$), the species tree tripartition T matches the coloring of all gene trees as required by conditions of (20), concluding $W[T] = Q = \sum_{G \in \mathcal{G}} W(T, G)$. \square

PROPOSITION 6. Before calling `OptimalTreeDP` in line 6 of Algorithm S1, lookup table W contains all tripartitions corresponding to internal nodes of all tree topologies in \mathcal{S} .

Proof. Each $\hat{S} \in \mathcal{S}$ places i above a different node w of S creating a new node corresponding to tripartition $\mathcal{L}_w|\{i\}|\mathcal{L}_S - \mathcal{L}_w$ covered in line 24. Besides new nodes, each existing internal node w of S will correspond to a different tripartition after placing i onto S depending on the relative location of w and i . Let u, v denote the larger and the smaller child of w . Node w corresponds to $\mathcal{L}_u|\{i\} \cup \mathcal{L}_v|\mathcal{L}_S - \mathcal{L}_w$ if i is under u , corresponds to $\{i\} \cup \mathcal{L}_u|\mathcal{L}_v|\mathcal{L}_S - \mathcal{L}_w$ if i is under v , and corresponds to $\mathcal{L}_u|\mathcal{L}_v|\{i\} \cup \mathcal{L}_S - \mathcal{L}_w$ if i is above w . All three cases for each node w is covered in lines 20–22. \square

Proof of Theorem 4

THEOREM 4. If there exists a species tree topology S^* satisfying that for each quartet subtree $ab|cd$,

$$\sum_{G \in \mathcal{G}} w(ab|cd) > \max \left(\sum_{G \in \mathcal{G}} w(ac|bd), \sum_{G \in \mathcal{G}} w(ad|bc) \right), \quad (6)$$

then the output of Algorithm S2 will be S^* .

Proof. We start with a Corollary 1 of Theorem 3

COROLLARY 1. Assuming (6), if S is compatible with the true tree S^* , then S' is compatible with S^* .

By induction, W_i in line 8 of Algorithm S2 should contain all tripartitions of S^* , as at that time $S_i = S^*$ by Corollary 1. Consequentially, the output of Algorithm S2 must also be S^* . \square

Proof of Proposition 4

PROPOSITION 4. The time complexity of Algorithm S2 is $O(kHn^2 \log n)$.

Proof. We begin with a proposition and a corollary.

PROPOSITION 7. Procedure `ColorNode` on any species tree node w takes $O(kH|\mathcal{L}_w| \log |\mathcal{L}_w|)$ time.

PROOF (SKETCH) OF PROPOSITION 7. We can prove this proposition by induction. For an internal node w with larger child u and smaller child v , if for some constant $C \geq \frac{6}{\log 2}$, `ColorNode` on u calls `UpdateCounters` at most $Ck|\mathcal{L}_u|(\log |\mathcal{L}_u| + 1)$ times and `ColorNode` on u calls `UpdateCounters` at most $Ck|\mathcal{L}_v|(\log |\mathcal{L}_v| + 1)$ times, then `ColorNode` on w calls `UpdateCounters` at most

$$\begin{aligned} & Ck|\mathcal{L}_u|(\log |\mathcal{L}_u| + 1) + Ck|\mathcal{L}_v|(\log |\mathcal{L}_v| + 1) + 3k(|\mathcal{L}_v| + 1) \\ & \leq Ck|\mathcal{L}_u|(\log |\mathcal{L}_w| + 1) + Ck|\mathcal{L}_v|(\log \frac{|\mathcal{L}_w|}{2} + 1) + 6k|\mathcal{L}_v| \\ & \leq Ck|\mathcal{L}_u|(\log |\mathcal{L}_w| + 1) + Ck|\mathcal{L}_v|(\log |\mathcal{L}_w| + 1) - Ck|\mathcal{L}_v|\log 2 + 6k|\mathcal{L}_v| \\ & \leq Ck|\mathcal{L}_w|(\log |\mathcal{L}_w| + 1) + (6 - C\log 2)k|\mathcal{L}_v| \\ & \leq Ck|\mathcal{L}_w|(\log |\mathcal{L}_w| + 1) \text{ times.} \end{aligned}$$

It is easy to verify that each `UpdateCounters` takes $O(H_G)$ time where H_G is the height of the gene tree, and thus `ColorNode` on node w takes $O(kH|\mathcal{L}_w|\log|\mathcal{L}_w|)$ time. \square

COROLLARY 2 (Corollary of Proposition 7). *For any tree topology S with n species, the `Place` procedure on S takes $O(kHn\log n)$ time.*

`NaivePlacement` of taxon set T makes $r(|T|-3)$ calls to `Place`, each of which takes $O(kH|T|\log|T|)$ time. Thus, `NaivePlacement` takes $O(rkH|T|^2\log|T|)$ time and when $T=\mathcal{L}_S$ and $r=O(1)$, $O(rkH|T|^2\log|T|)=O(n^2kH\log n)$. \square

Proofs of Theorems 6 and Theorem 5

THEOREM 6. *Under the conditions of Theorem 4, the DAC Algorithm S3 will output S^* .*

Proof. By Theorem 4, S_i in line 5 of Algorithm S3 are compatible with S^* . With Corollary 1, by induction, each S_e in line 21 of Algorithm S3 is compatible with S^* . Consequentially, W_i in line 26 contain all tripartitions of S^* , as at that time $S'_i=S^*$, and the output of Algorithm S3 must also be S^* . \square

THEOREM 5. *When the inequality condition in Theorem 4 is satisfied, then the time complexity of the DAC algorithm is $O(n^{1.5+\epsilon}kH)$ with arbitrarily high probability.*

PROOF (SKETCH). From the inequality (6), we can trivially deduct that S^* is the species tree topology that maximizes the weighted quartet score, and each S_i in line 5 of Algorithm S3 is compatible to S^* . Also, each C_e in line 15 of Algorithm S3 equals the set of species under the edges coming off of the internal nodes on the path of S^* corresponding to e .

We now introduce a proposition

PROPOSITION 8. *With high probability, $\max_{e \in E_{S_i}} |C_e| \leq 2\sqrt{n}\log n + O(\sqrt{n})$.*

Proof. For each pair of nodes u, v of S^* , let $C_{u,v} := \{x : x \in \mathcal{L}_S, u \text{ is not on } \mathcal{P}_{x,v} \text{ and } v \text{ is not on } \mathcal{P}_{x,u}\}$. It is easy to verify that for every e of S_i , $C_e = C_{u,v}$ for some nodes u, v of S^* . For every u and v that are sufficiently apart so that $C_{u,v}$ has $2\sqrt{n}\log n + \omega(\sqrt{n})$ elements and a random T_i in line 4 of Algorithm S3,

$$P(C_{u,v} \cap T_i = \emptyset) = \left(1 - \frac{1}{\sqrt{n}}\right)^{|C_{u,v}|} \leq e^{-\frac{1}{\sqrt{n}}|C_{u,v}|} = \frac{1}{n^2} e^{-\omega(1)} = o\left(\frac{1}{n^2}\right).$$

By union bound, the probability that there exists a pair of nodes u, v of S^* such that $|C_{u,v}| \geq 2\sqrt{n}\log n + \omega(\sqrt{n})$ and $C_{u,v} \cap T_i = \emptyset$ is $o(1)$. Since, by definition, $C_e \cap T_i = \emptyset$ for every C_e , with high probability, there exists no C_e having $2\sqrt{n}\log n + \omega(\sqrt{n})$ elements. \square

Since $|T_i| \sim \text{Binomial}(n, \frac{1}{\sqrt{n}})$, with high probability $|T_i| = O(\sqrt{n})$ and calling `NaivePlacement` on line 5 takes $O(n^{1.5}kH\log n)$ time. It is easy to confirm that $C_\emptyset = \emptyset$ and every call to `Place` takes as input a species tree topology of $O(\sqrt{n}\log n)$ species with high probability. Thus, with high probability, each call to `Place` takes $O(\sqrt{n}kH\log^2 n \log \log n)$ time and all $O(n)$ calls to `Place` takes $O(n^{1.5}kH\log^2 n \log \log n)$ time. Therefore, the time complexity of the DAC algorithm is $O(n^{1.5}kH\log^2 n \log \log n) = O(n^{1.5+\epsilon}kH)$ with high probability. \square